

《科学研究費補助金(総合研究A)研究集会報告集》

代 数 的 組 合 せ 論
報 告 集

1989年11月28日～30日

於 甲 南 大 学

ま え が き

この報告集は、1989年11月28日～30日の3日間、甲南大学理学部に於て行なわれた研究集会、「代数的組合せ論」の講演記録である。

集会はJ. Seberry、R. G. Stanton両教授を迎え、多数の参加者を得て、盛会であった。会場の御世話、会の運営には、甲南大学伊藤昇教授、大阪教育大学大山豪、伊藤達郎両教授、鈴木寛氏にたいへん御世話になった。これら、集会に御協力戴いた方々に深く感謝致します。

なお、研究集会の諸経費、この報告集の作成経費は、文部省科学研究費総合研究A（代表者 京都大学土方弘明教授）に依った。

1989年12月

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Mutually Balanced Nested Designs

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Abstract

A mutually balanced nested design of strength t is introduced. It is shown that such a nested design is equivalent to a balanced array of strength t . Some recursive constructions are given for a mutually balanced nested design of strength 2. Some constructions are also given for a balanced incomplete array of strength 2 which are based on a mutually balanced nested design.

1. Introduction

Let S be a set $\{1, 2, \dots, s\}$ of s symbols and let X be the set of all t -dimensional vectors with elements from S . A *balanced array* of strength t with s symbols is a $v \times b$ array A whose elements are from S satisfying the following conditions:

(i) in any t -rowed subarray A_0 of A , the number of columns of A_0 which are equal to x is $\mu(x)$ for any $x \in X$,

(ii) for any permutation matrix P of order t and for any $x \in X$, $\mu(Px) = \mu(x)$.

Such an array is denoted by $BA_\mu(t, s; v)$. If $\mu(x) = \mu(y)$ for every $x, y \in X$, then the array is called an *orthogonal array* of strength t with s symbols.

A balanced array was first introduced and studied by Chakravarti [1,2] in connection with some class of statistical designs. Many authors (e.g. [4-7]) have researched such an array. Kuriki and Fuji-Hara [4] defined an (r, λ) -design with mutually balanced nested subdesigns which is equivalent to a balanced array of strength two. For strength $t(\geq 2)$, we introduce a nested design of strength t satisfying some conditions in Section 3 and show that such a nested design is equivalent

to a balanced array of strength t . We give some recursive constructions of a nested design of strength 2. Gill [3] generalized a balanced array to a balanced incomplete array and gave a construction of a balanced incomplete array of strength 2 there. As an application of results obtained in Section 3, we give some constructions of a balanced incomplete array of strength 2 which are based on a nested design in Section 4.

2. Balanced arrays

Let $W_t^s = \{(d_1, d_2, \dots, d_s); d_i \geq 0, \sum_{i=1}^s d_i = t\}$. For $\mathbf{d} = (d_1, d_2, \dots, d_s) \in W_t^s$, let $X_{\mathbf{d}}$ be the set of all t -dimensional vectors which contain each symbol $i \in S$ in d_i positions. If A is a balanced array of strength t with s symbols, then $\mu(\mathbf{x}) = \mu(\mathbf{y})$ holds for $\mathbf{x}, \mathbf{y} \in X_{\mathbf{d}}$. Obviously, if $\mathbf{x} \in X_{\mathbf{d}}$, then $X_{\mathbf{d}} = \{P\mathbf{x}; \text{ for every permutation matrix } P\}$. Therefore, we can rewrite the conditions (i) and (ii) given in Section 1 for a balanced array as follows:

$C(t, \mathbf{d})$: in any t -rowed subarray A_0 of A , the number of columns of A_0 which are equal to $\mathbf{x} \in X_{\mathbf{d}}$ is $\nu_t(\mathbf{d})$.

If $C(t, \mathbf{d})$ are satisfied for every $\mathbf{d} \in W_t^s$, then A is a balanced array of strength t with s symbols. Throughout this paper, we use the conditions $C(t, \mathbf{d})$ for a balanced array instead of usual conditions given in Section 1. Note that $C(0, \mathbf{0})$ is always satisfied and $\nu_0(\mathbf{0})$ is the number of columns of A .

Lemma 2.1. *In an array with elements from S , if s conditions in the $s + 1$ conditions $C(t - 1, \mathbf{d}), C(t, \mathbf{d} + \mathbf{e}_1), C(t, \mathbf{d} + \mathbf{e}_2), \dots, C(t, \mathbf{d} + \mathbf{e}_s)$ are satisfied, then the remaining condition is also satisfied, where $\mathbf{d} \in W_{t-1}^s$ and \mathbf{e}_i denotes an s -dimensional vector with unity in the i th position and zero in all other positions.*

Proof. Let A be an array with elements from S . Consider a $(t - 1)$ -rowed subarray A_0 of A and a t -rowed subarray A_1 of A containing A_0 as the first $t - 1$ rows. For $\mathbf{x} \in X_{\mathbf{d}}, \mathbf{d} \in W_{t-1}^s$, the number of columns of the subarray A_1 which

are equal to $[x|i]$ is denoted by $\alpha([x|i])$ for each $i \in \mathcal{S}$, where $[a|b]$ denotes the juxtaposition of two vectors a and b . In the subarray A_0 , the number of columns which are equal to x is denoted by $\beta(x)$. Since the vectors $[x|1], [x|2], \dots, [x|s]$ are all distinct,

$$\beta(x) = \alpha([x|1]) + \alpha([x|2]) + \dots + \alpha([x|s])$$

holds.

Now we assume that the following s conditions are satisfied:

$$C(t-1, d), C(t, d + e_1), \dots, C(t, d + e_{s-1}),$$

then $\beta(x) = \nu_{t-1}(d)$, $\alpha([x|1]) = \nu_t(d + e_1)$, \dots , $\alpha([x|s-1]) = \nu_t(d + e_{s-1})$ and the numbers are independent of choice of A_0 and A_1 . Hence, $\alpha([x|s])$ is determined uniquely and the number is independent of choice of A_0 and A_1 . Therefore, the remaining condition $C(t, d + e_s)$ is also satisfied.

For other cases of the assumption, we can prove the statement of this lemma by the similar technique. □

As an immediate consequence of Lemma 2.1, we have:

Lemma 2.2. *In an array with elements from \mathcal{S} , if $C(t, f)$ are satisfied for every $f \in W_t^s$, then $C(t-1, d)$ are also satisfied for every $d \in W_{t-1}^s$.*

Lemma 2.2 implies that if A is a balanced array of strength t with s symbols, then A is also a balanced array of strength $t-1$ with s symbols. Furthermore, by the proof of Lemma 2.1, in a balanced array of strength t with s symbols,

$$\nu_k(d) = \sum_{i=1}^s \nu_{k+1}(d + e_i)$$

hold for every $d \in W_k^s$, $k < t$.

Now we clarify the conditions to be a balanced array which used in succeeding sections.

Lemma 2.3. *In a balanced array A of strength $t-1$ with s symbols, if $C(t, \mathbf{d})$ are satisfied for every $\mathbf{d} = (d_1, d_2, \dots, d_s) \in W_t^s$ such that $d_s = 0$, then A is also a balanced array of strength t with s symbols.*

Proof. We will show that $C(t, \mathbf{g})$ are satisfied for every $\mathbf{g} = (g_1, g_2, \dots, g_s) \in W_t^s$ by induction on g_s .

Since A is a balanced array of strength $t-1$ with s symbols, $C(t-1, \mathbf{f})$ are satisfied for every $\mathbf{f} = (f_1, f_2, \dots, f_s) \in W_{t-1}^s$. For every $\mathbf{f} \in W_{t-1}^s$ such that $f_s = 0$, each of vectors $\mathbf{f} + \mathbf{e}_1, \dots, \mathbf{f} + \mathbf{e}_{s-1}$ does not contain the symbol s . Hence, by the conditions of this lemma, $C(t, \mathbf{f} + \mathbf{e}_1), \dots, C(t, \mathbf{f} + \mathbf{e}_{s-1})$ are satisfied. Therefore, by Lemma 2.1, $C(t, \mathbf{f} + \mathbf{e}_s)$ is also satisfied. This implies that $C(t, \mathbf{g})$ are satisfied for every $\mathbf{g} \in W_t^s$ such that $g_s = 1$.

The remaining part of induction is straightforward, so omitted here. □

By use of Lemma 2.3, noting that $C(0, \mathbf{0})$ is always satisfied, we have the following:

Theorem 2.4. *In an array A with elements from S , if $C(k, \mathbf{d})$ are satisfied for every $\mathbf{d} = (d_1, d_2, \dots, d_s) \in W_k^s$ such that $d_s = 0$, $1 \leq k \leq t$, then A is a balanced array of strength t with s symbols.*

3. Mutually Balanced Nested Designs

Let V be a v -set (called *points* or *varieties*) and \mathcal{B} be a collection of subsets of V (called *blocks*). Then the pair (V, \mathcal{B}) is called a *design*. A *t -wise balanced design* $S_\lambda(t, K; v)$ is a design (V, \mathcal{B}) satisfying the following condition:

for any t -subset T of V , the number of blocks containing T is λ which is independent of the t -subset T chosen.

If, for any u -subset U ($u \leq t$) of V , the number of blocks containing U is constant (say, λ_u) which is independent of the u -subset U chosen, then the design is called

a regular t -wise balanced design $R_\lambda(t, K; v)$, where K denotes the set of block sizes of \mathcal{B} and $\lambda = (\lambda_1, \lambda_2, \dots, \lambda_t)$. In the case $t = 2$, the designs are called a *pairwise balanced design* and a *regular pairwise balanced design* (or an (r, λ) -design, where $r = \lambda_1$ and $\lambda = \lambda_2$), respectively.

Let $\mathcal{B} = \{B_1, B_2, \dots, B_b\}$ and $\mathcal{B}' = \{B'_1, B'_2, \dots, B'_b\}$. If $B'_j \subseteq B_j$ for $1 \leq j \leq b$, then (V, \mathcal{B}') is called a *subdesign* of a design (V, \mathcal{B}) . Note that B'_j can be the empty set. Suppose that there exist s subdesigns $(V, \mathcal{B}^{(i)})$, $i = 1, 2, \dots, s$, of (V, \mathcal{B}) , such that $\bigcup_{i=1}^s B_j^{(i)} = B_j$ and $B_j^{(i)} \cap B_j^{(i')} = \phi$ if $i \neq i'$, for $1 \leq j \leq b$, where $\mathcal{B}^{(i)} = \{B_1^{(i)}, B_2^{(i)}, \dots, B_b^{(i)}\}$. For convenience, we may write a block B of \mathcal{B} as $B = \{B^{(1)}; B^{(2)}; \dots; B^{(s)}\}$, $B^{(i)} \in \mathcal{B}^{(i)}$. If the subdesigns satisfy the following conditions:

- (i) each subdesign $(V, \mathcal{B}^{(i)})$ is an (r_i, λ_i) -design,
- (ii) for any distinct points $x, y \in V$, the number of blocks $B = \{B^{(1)}; B^{(2)}; \dots; B^{(s)}\} \in \mathcal{B}$ which contain x in $B^{(i)}$ and y in $B^{(h)}$ is exactly λ_{ih} ,

then we call it an (r, λ) -design with *mutually balanced nested subdesigns*. Kuriki and Fuji-Hara [4] defined the design and proved that an (r, λ) -design with mutually balanced nested s subdesigns is equivalent to a balanced array of strength 2 with $s + 1$ symbols.

Here we generalize an (r, λ) -design with mutually balanced nested subdesigns to strength t . For $\mathbf{d} = (d_1, d_2, \dots, d_s) \in W_k^s$, let S_1, S_2, \dots, S_s be mutually disjoint subsets of V such that $|S_i| = d_i$. Then we consider the following condition:

$L(k, \mathbf{d})$: the number of blocks $B = \{B^{(1)}; B^{(2)}; \dots; B^{(s)}\} \in \mathcal{B}$ such that each $S_i \subset B^{(i)}$ is exactly $\eta_k(\mathbf{d})$ which is independent of choice of S_1, S_2, \dots, S_s .

If $L(k, \mathbf{d})$ are satisfied for every $\mathbf{d} \in W_k^s$, $1 \leq k \leq t$, then (V, \mathcal{B}) is called a *mutually balanced nested design* (MBND) of strength t . Such a design is denoted by $M_\eta(t, s; v)$, $\eta = \{\eta_1, \eta_2, \dots, \eta_t\}$, where each η_k is an index function from W_k^s to nonnegative integers and $v = |V|$. Note that $L(0, \mathbf{O})$ is always satisfied and $\eta_0(\mathbf{O})$ is the number of blocks. If $L(2, 2e_i)$ is satisfied, then the i th subdesign is

a pairwise balanced design. If $L(1, \mathbf{e}_i)$ and $L(2, 2\mathbf{e}_i)$ are satisfied, then the i th subdesign is an (r, λ) -design such that $r = \eta_1(\mathbf{e}_i)$ and $\lambda = \eta_2(2\mathbf{e}_i)$.

The conditions of the definition of a MBND do not mention about the original big design (V, \mathcal{B}) .

Lemma 3.1. *A mutually balanced nested design $M_\eta(t, s; v)$ (V, \mathcal{B}) is also a regular t -wise balanced design with parameters*

$$\lambda_k = \sum_{\mathbf{d} \in W_k^s} \binom{k}{d_1, d_2, \dots, d_s} \eta_k(\mathbf{d})$$

for $1 \leq k \leq t$, where $\mathbf{d} = (d_1, d_2, \dots, d_s)$ and $\binom{k}{d_1, d_2, \dots, d_s}$ denotes the multinomial coefficient.

Proof. Consider a k -subset U of V for $1 \leq k \leq t$. The number of partition of U into U_1, U_2, \dots, U_s such that $|U_i| = d_i$ is $\binom{k}{d_1, d_2, \dots, d_s}$. For each partition, the number of blocks $B = \{B^{(1)}, B^{(2)}, \dots, B^{(s)}\} \in \mathcal{B}$ such that each $U_i \subset B^{(i)}$ is $\eta_k(\mathbf{d})$. Therefore, the number of blocks of \mathcal{B} containing U is

$$\sum_{\mathbf{d} \in W_k^s} \binom{k}{d_1, d_2, \dots, d_s} \eta_k(\mathbf{d}),$$

which is independent of the k -subset U chosen. Hence (V, \mathcal{B}) is a regular t -wise balanced design. \square

Now we show that a MBND of strength t with s subdesigns is equivalent to a balanced array of strength t with $s + 1$ symbols.

Theorem 3.2. *There exists a mutually balanced nested design $M_\eta(t, s; v)$ if and only if there exists a balanced array $BA_\mu(t, s + 1; v)$ such that*

$$\mu(\mathbf{x}) = \sum_{h=0}^{t-k} (-1)^h \binom{t-k}{h} \sum_{i_1=1}^s \cdots \sum_{i_h=1}^s \eta_{k+h}(\mathbf{d} + \mathbf{e}_{i_1} + \cdots + \mathbf{e}_{i_h}) \quad (3-1)$$

for a t -dimensional vector \mathbf{x} containing each symbol $i \in \mathcal{S}$ in d_i positions and the symbol $s + 1$ in the remaining positions, where $\mathbf{d} = (d_1, d_2, \dots, d_s) \in W_k^s$, $0 \leq k \leq t$, and $\binom{t-k}{h}$ denotes the binomial coefficient.

Proof. Suppose that (V, \mathcal{B}) is a $M_\eta(t, s; v)$. Let $\mathcal{B} = \{B_1, B_2, \dots, B_b\}$, $b = \eta_0(\mathbf{0})$. Each block is partitioned into s subblocks, i.e., $B_j = \{B_j^{(1)}; B_j^{(2)}; \dots; B_j^{(s)}\}$, for $1 \leq j \leq b$. From the blocks, we define a $v \times b$ array $\mathbf{A} = [a_{xj}]$ as

$$a_{xj} = \begin{cases} i, & \text{if a point } x \text{ of } V \text{ occurs } B_j^{(i)}, \\ s+1, & \text{otherwise.} \end{cases}$$

For $\mathbf{d} = (d_1, d_2, \dots, d_s) \in W_k^s$, $1 \leq k \leq t$, let S_1, S_2, \dots, S_s be mutually disjoint subsets of V such that $|S_i| = d_i$. Since (V, \mathcal{B}) is a $M_\eta(t, s; v)$, the number of blocks $B = \{B^{(1)}; B^{(2)}; \dots; B^{(s)}\} \in \mathcal{B}$ such that each $S_i \subset B^{(i)}$ is $\eta_k(\mathbf{d})$ which is independent of choice of S_1, S_2, \dots, S_s . Hence, in the array \mathbf{A} , $C(k, \mathbf{d}')$ is satisfied for $\mathbf{d}' = (d'_1, d'_2, \dots, d'_{s+1}) \in W_k^{s+1}$ such that $d'_i = d_i, i = 1, 2, \dots, s$, and $d'_{s+1} = 0$. Therefore, by Theorem 2.4, \mathbf{A} is a balanced array of strength t with $s+1$ symbols.

Conversely, suppose that \mathbf{A} is a $BA_\mu(t, s+1; v)$. Correspond points of a v -set V to rows of \mathbf{A} and blocks of a collection \mathcal{B} to columns of \mathbf{A} . Each block of \mathcal{B} consists of points of V corresponding to entries $1, 2, \dots, s$ of \mathbf{A} . For each block $B_j \in \mathcal{B}$, $1 \leq j \leq b$, we partition B_j into s subblocks $B_j^{(1)}, B_j^{(2)}, \dots, B_j^{(s)}$ such that $B_j^{(i)}$ consists of points with the entry i . For $\mathbf{d}' = (d'_1, d'_2, \dots, d'_{s+1}) \in W_k^{s+1}$ such that $d'_{s+1} = 0, 1 \leq k \leq t$, let \mathbf{y} be a k -dimensional vector containing each symbol $i \in \mathcal{S}$ in d'_i positions. Since \mathbf{A} is a $BA_\mu(t, s+1; v)$, in any k -rowed subarray \mathbf{A}_0 of \mathbf{A} , the number of columns of \mathbf{A}_0 which are equal to \mathbf{y} is $\nu_k(\mathbf{d}')$. From the definition of the design (V, \mathcal{B}) , $L(k, \mathbf{d})$ is satisfied for $\mathbf{d} = (d_1, d_2, \dots, d_s) \in W_k^s$ such that $d_i = d'_i, i = 1, 2, \dots, s$. Therefore, (V, \mathcal{B}) is a $M_\eta(t, s; v)$.

Finally, we show (3.1) by induction on k . In the case $k = t$, \mathbf{x} does not contain the symbol $s+1$. Hence $\mu(\mathbf{x}) = \eta_t(\mathbf{d})$ holds.

Assuming that (3.1) holds for $k = u+1, u+2, \dots, t$, we will show that (3.1) holds for $k = u$. Since $\mu(\mathbf{x}) = \mu(\mathbf{P}\mathbf{x})$ for any permutation matrix \mathbf{P} of order t , we may assume that, without loss of generality, \mathbf{x} contains the symbol $s+1$ in the last $t-u$ positions, i.e., $\mathbf{x} = [\mathbf{x}^* | s+1 \dots s+1]$, where \mathbf{x}^* is a u -dimensional vector

containing each symbol $i \in \mathcal{S}$ in d_i positions for $\mathbf{d} \in W_u^s$. Then, we have

$$\begin{aligned} \mu(\mathbf{x}) &= \mu([\mathbf{x}^* | s+1 \dots s+1]) \\ &= \eta_u(\mathbf{d}) - \sum_{p=1}^{t-u} \binom{t-u}{p} \sum_{i_1=1}^s \cdots \sum_{i_p=1}^s \mu([\mathbf{x}^* | i_1 \dots i_p s+1 \dots s+1]). \end{aligned} \quad (3-2)$$

Applying the assumption to (3.2), we have

$$\begin{aligned} \mu(\mathbf{x}) &= \eta_u(\mathbf{d}) - \sum_{p=1}^{t-u} \binom{t-u}{p} \sum_{i_1=1}^s \cdots \sum_{i_p=1}^s \sum_{l=0}^{t-u-p} (-1)^l \binom{t-u-p}{l} \\ &\quad \sum_{i_{p+1}=1}^s \cdots \sum_{i_{p+l}=1}^s \eta_{u+p+l}(\mathbf{d} + \mathbf{e}_{i_1} + \cdots + \mathbf{e}_{i_p} + \cdots + \mathbf{e}_{i_{p+1}} + \cdots + \mathbf{e}_{i_{p+l}}) \\ &= \eta_u(\mathbf{d}) - \sum_{p=1}^{t-u} \binom{t-u}{p} \sum_{h=p}^{t-u} (-1)^{h-p} \\ &\quad \binom{t-u-p}{h-p} \sum_{i_1=1}^s \cdots \sum_{i_h=1}^s \eta_{u+h}(\mathbf{d} + \mathbf{e}_{i_1} + \cdots + \mathbf{e}_{i_h}) \\ &= \eta_u(\mathbf{d}) - \sum_{h=1}^{t-u} (-1)^h \binom{t-u}{h} \sum_{p=1}^h (-1)^p \binom{h}{p} \sum_{i_1=1}^s \cdots \sum_{i_h=1}^s \eta_{u+h}(\mathbf{d} + \mathbf{e}_{i_1} + \cdots + \mathbf{e}_{i_h}) \\ &= \eta_u(\mathbf{d}) + \sum_{h=1}^{t-u} (-1)^h \binom{t-u}{h} \sum_{i_1=1}^s \cdots \sum_{i_h=1}^s \eta_{u+h}(\mathbf{d} + \mathbf{e}_{i_1} + \cdots + \mathbf{e}_{i_h}) \\ &= \sum_{h=0}^{t-u} (-1)^h \binom{t-u}{h} \sum_{i_1=1}^s \cdots \sum_{i_h=1}^s \eta_{u+h}(\mathbf{d} + \mathbf{e}_{i_1} + \cdots + \mathbf{e}_{i_h}). \end{aligned}$$

Therefore, (3.1) holds for $k = u$. □

In the case $t = 2$, we give two recursive constructions of a $M_{\mathfrak{q}}(2, s; v)$. Constructions are shown for only $t = 2$. However, it is not difficult to generalize the constructions to any integer t .

Theorem 3.3. *If there exists a $M_{\mathfrak{q}}(2, s; v)$ with the set of block sizes K and if there exist (r', λ') -designs with v' varieties for all v' in K , then there exists a $M_{\mathfrak{q}}(2, s; v)$ such that*

$$\eta'_1(\mathbf{e}_i) = r'\eta_1(\mathbf{e}_i), \quad i = 1, 2, \dots, s,$$

and

$$\eta'_2(\mathbf{e}_i + \mathbf{e}_j) = \lambda'\eta_2(\mathbf{e}_i + \mathbf{e}_j), \quad i, j = 1, 2, \dots, s.$$

Proof. Let (V, \mathcal{B}) be a $M_{\mathfrak{q}}(2, s; v)$ with block sizes K and let $B = \{B^{(1)}; B^{(2)}; \dots; B^{(s)}\} \in \mathcal{B}$ such that $|B| = v'$. From the assumption, there exists an (r', λ') -design (V', \mathcal{E}) with $v' \in K$ varieties. Relabeling V' by varieties of B , we construct a design (B, \mathcal{E}_B) . Then we partition each block of \mathcal{E}_B by the following way:

$$E = \{E^{(1)}; E^{(2)}; \dots; E^{(s)}\} \in \mathcal{E}_B \quad \text{if and only if} \quad E^{(i)} \subseteq B^{(i)}.$$

Applying this method for each block B of \mathcal{B} , we construct a new design (V, \mathcal{B}') . We will show that it is also a MBND. Consider each new subdesign $(V, \mathcal{B}^{(i)'})$, where $\mathcal{B}^{(i)'}$ denotes a collection of i th subblocks of B' . For any point x of V , if $B^{(i)}$ contains x , then x is contained in r' blocks of $\mathcal{E}_B^{(i)}$. Furthermore, for any pair of distinct points x and y of V , if $B^{(i)}$ contains $\{x, y\}$, then the pair is contained in λ' blocks of $\mathcal{E}_B^{(i)}$, where $\mathcal{E}_B^{(i)}$ denotes a collection of i th subblocks of \mathcal{E}_B . Hence $(V, \mathcal{B}^{(i)'})$ is an $(r'\eta_1(\mathbf{e}_i), \lambda'\eta_2(2\mathbf{e}_i))$ -design.

Next consider the condition $L(2, \mathbf{e}_i + \mathbf{e}_j)$ about the new design (V, \mathcal{B}') . For distinct points x and y of V such that $x \in B^{(i)}$ and $y \in B^{(j)}$, the number of blocks of \mathcal{E}_B such that $x \in E^{(i)}$ and $y \in E^{(j)}$ is λ' . So $\eta'_2(\mathbf{e}_i + \mathbf{e}_j) = \lambda'\eta_2(\mathbf{e}_i + \mathbf{e}_j)$ holds. Therefore, (V, \mathcal{B}') is a MBND. \square

Corollary 3.4. *If there exists a $M_{\mathfrak{q}}(2, s; v)$ with the set of block sizes K and if there exist $(r, 1)$ -designs with v' varieties for some v' in K , then there exists a new $M_{\mathfrak{q}}(2, s; v)$ with more blocks and the different set of block sizes.*

Proof. By the same way as Theorem 3.3, we can embed small $(r', 1)$ -designs with v' varieties if $v' \in K$. It is easy to see that the condition $L(2, e_i + e_j)$ about the new design is satisfied. However $L(1, e_i)$ is usually not satisfied. To satisfy the condition, we add some new blocks such that the i th subblock consists only one point and the remaining subblocks are empty, $\{\phi; \dots; \phi; x; \phi; \dots; \phi\}$ for a suitable point x of V , until $L(1, e_i)$ is satisfied. \square

4. Balanced Incomplete Array

A *balanced incomplete array* of strength t with s symbols and d blanks is a $v \times b$ array A whose elements are from \mathcal{S} satisfying the following conditions:

- (i) every row has exactly d blanks,
- (ii) in any t -rowed subarray A_0 of A , the number of columns of A_0 which are equal to \mathbf{x} is $\mu(\mathbf{x})$ for any $\mathbf{x} \in \mathbf{X}$,
- (iii) for any permutation matrix P of order t and for any $\mathbf{x} \in \mathbf{X}$, $\mu(P\mathbf{x}) = \mu(\mathbf{x})$.

Here, the vector \mathbf{x} in the conditions (ii) and (iii) does not contain any blank. Such an array is denoted by $BIA_\mu(t, s, d; v)$. Note that a $BIA_\mu(t, s, d; v)$ is not always of strength $t - 1$. By deleting i symbols of a $BA_\mu(t, s; v)$, a $BIA_\mu(t, s - i, d; v)$ is obtained for a suitable d . Gill [3] defined a balanced incomplete array and he gave a construction of such an array of strength 2 there.

Using the conditions $L(k, d)$ in Section 3, we characterize a balanced incomplete array.

Theorem 4.1. *A $BIA_\mu(t, s, d; v)$ is equivalent to a design (V, \mathcal{B}) with s subdesigns satisfying the following conditions:*

- (i) for any point $x \in V$, x is contained in exactly $b - d$ blocks of \mathcal{B} ,
- (ii) $L(t, d)$ are satisfied for every $d \in W_t^s$,

where $v = |V|$ and $b = |\mathcal{B}|$.

Proof. We define the similar correspondence to Theorem 3.2 between a incomplete array with s symbols and a design (V, \mathcal{B}) with s subdesigns replacing the symbol $s + 1$ of a balanced array in Theorem 3.2 to blank. Then, every row of the incomplete array has exactly d blanks if and only if any point $x \in V$ is contained in exactly $b - d$ blocks of \mathcal{B} . Futhermore, it is obvious that two conditions $C(t, d)$ and $L(t, d)$ are equivalent for every $d \in W_t^s$. Then $\mu(\mathbf{x}) = \eta_t(\mathbf{d})$ holds for a t -dimensional vector \mathbf{x} containing each symbol $i \in \mathcal{S}$ in d_i positions, where $\mathbf{d} = (d_1, d_2, \dots, d_s)$. \square

In the case $t = 2$, we have $W_2^s = \{\mathbf{e}_i + \mathbf{e}_j; i, j = 1, 2, \dots, s\}$. Hence Theorem 4.1 yields:

Corollary 4.2. *A $BIA_\mu(2, s, d; v)$ is equivalent to a design (V, \mathcal{B}) with s subdesigns satisfying the conditions:*

- (i) *for any point $x \in V$, x is contained in exactly $b - d$ blocks of \mathcal{B} ,*
- (ii) *each subdesign $(V, \mathcal{B}^{(i)})$ is a pairwise balanced design $S_\lambda(2, K; v)$ such that*

$$\lambda = \eta_2(2\mathbf{e}_i),$$
- (iii) *for any distinct points $x, y \in V$, the number of blocks $B = \{B^{(1)}; B^{(2)}; \dots; B^{(s)}\} \in \mathcal{B}$ such that $x \in B^{(i)}$ and $y \in B^{(j)}$ is $\eta_2(\mathbf{e}_i + \mathbf{e}_j)$,*

where $b = |\mathcal{B}|$ and $B^{(i)}$ denotes the i th subblock of $B \in \mathcal{B}$.

In order to construct a $BIA_\mu(2, s, d; v)$, we apply a pairwise balanced design instead of an (r, λ) -design to the constructions given in Theorems 3.3 and 3.4 and to satisfy the condition (i) of Corollary 4.2, we add some blocks containing only one point.

Theorem 4.3. *If there exists a $M_\eta(2, s; v)$ with the set of block sizes K and if there exist pairwise balanced designs $S_{\lambda'}(2, K'; v')$ for all v' in K , then there exists a $BIA_\mu(2, s, d; v)$ such that*

$$\mu(\{ij\}) = \lambda' \eta_2(\mathbf{e}_i + \mathbf{e}_j), \quad i, j = 1, 2, \dots, s,$$

for a suitable d .

Theorem 4.4 *If there exists a $M_{\eta}(2, s; v)$ with the set of block sizes K and if there exist pairwise balanced designs $S_1(2, K'; v')$ for some v' in K , then there exists a $BIA_{\mu}(2, s, d; v)$ such that*

$$\mu([ij]) = \eta_2(\mathbf{e}_i + \mathbf{e}_j), \quad i, j = 1, 2, \dots, s,$$

for a suitable d .

Note that a $BIA_{\mu}(2, s, d; v)$ constructed by Gill [3] is always a balanced array of strength 2 with $s + 1$ symbols.

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An Analogy Between Locally Hamming Graphs and Riemann Surfaces.

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December 16, 1989

Abstract

A locally hamming graph is an undirected simple graph locally isomorphic to a hamming scheme. It was known that every locally hamming graph has a covering which is a hamming scheme. This paper proves the universality of this covering, through which we obtain 1-to-1 correspondence of Galois type between conjugate classes of discrete subgroups of the automorphic group of the r -dimensional hamming scheme and isomorphic classes of r -dimensional locally hamming graphs, analogously to the case of a Riemann surface and its fundamental group.

1 Locally Hamming Graphs and Coverings

In this paper all graphs are undirected, simple, and connected. For a graph G , $V(G)$ denotes the vertex set and $E(G)$ denotes the edge set of G . For a finite set V , $\#(V)$ denotes the cardinality of V . By $H(r)$ we denote the r -dimensional hamming scheme for $r \geq 1$; that is, $H(r)$ is such a graph that its vertex set is the vector space \mathbb{F}_2^r and $u, v \in \mathbb{F}_2^r$ are adjacent if and only if the hamming distance $d(u, v) = 1$; i.e., $\#\{i \mid u_i \neq v_i\} = 1$ where $u = (u_1, \dots, u_r)$ and $v = (v_1, \dots, v_r)$.

Let us call $H(3)$ a *cubic*, and the graph shown in the next figure a *tulip* with petals p, q, r .



A connected graph G is said to be *locally hamming* if it satisfies:

1. G has no triangle.
2. For $u, v \in V(G)$ satisfying $d(u, v) = 2$, the span of u and v is a quadrilateral; in other words, there exist exactly two vertices adjacent to both u and v .
3. Let T be a subgraph of G isomorphic to a tulip with petals p, q, r . Then there exists a vertex $x \in V(G)$ adjacent to all p, q , and r . (The uniqueness of x follows from the condition 2.)

It can easily be proved that $H(r)$ with $r \geq 3$ is a locally hamming graph, and that a distance regular graph with parameters $a_1 = 0$, $a_2 = 0$, $c_2 = 2$, and $c_3 = 3$ (see [1] for the definition of a distance regular graph and its parameters) is also a locally hamming graph (see [5][7]). Recently, a comprehensive book on this material is also published[4].

A mapping $f : V(G) \rightarrow V(H)$ is a *morphism* $f : G \rightarrow H$ if $uv \in E(G)$ implies $f(u)f(v) \in E(H)$. A morphism $f : G \rightarrow H$ is a *locally injection* (or *locally isomorphism*) if for every $v \in V(G)$, $f|_{N(v)} : N(v) \rightarrow N(f(v))$ is injective (resp. bijective), where $N(v)$ denotes the set of vertices adjacent to v . A morphism f is a *covering* if it is locally isomorphic and surjective as a mapping $V(G) \rightarrow V(H)$.

Proposition 1. (Existence of a prolongation and its uniqueness.)

Let G be a locally hamming graph, and u be a vertex of G . Let $H(r)$ be the r -dimensional hamming scheme, and v be a vertex of $H(r)$. Given an injection $g : N(v) \rightarrow N(u)$, there exists a unique locally injection $f : H(r) \rightarrow G$ such that $f(v) = u$ and $f|_{N(v)} = g$.

This proposition and even its generalization were already proved in [3][5][7] somewhat implicitly as for uniqueness. So, this paper provides only a sketch of proof to lessen the readers' effort.

Sketch of proof. Put $H_i := \{x \in V(H(r)) \mid d(x, v) = i\}$ for $i = 0, 1, \dots, r$. We construct locally injections

$$f_i : \langle H_0 \cup H_1 \cup \dots \cup H_i \rangle \rightarrow G$$

with $f_i|_{N(v)} = g$ by induction on i . (Here $\langle \rangle$ denotes the induced subgraph.) $i = 1$. Clearly $f|_{H_1} = g$ gives the unique solution.

$i = 2$. Let $x \in H_2$. Then there exist $x_1, x_2 \in H_1$ such that x_1x, x_2x are edges. Condition 2 in the definition of locally hamming graphs assures that there exists a unique $y \in V(G)$ such that both $f_1(x_1)y$ and $f_1(x_2)y$ are edges in G . Clearly $f_2(x)$ must coincide with y , and this provides the unique f_2 . Locally injectiveness follows by a straightforward argument using Conditions 1 and 2.

$i \geq 3$. Let $x \in H_i$, then there exist exactly i vertices in H_{i-1} adjacent to x . We denote these vertices by x_1, \dots, x_i . There is an $x' \in H_{i-2}$ adjacent to both x_1 and x_2 . Then Condition 2 assures the existence of a unique $y \in V(G)$, $y \neq x'$, which is adjacent to both $f_{i-1}(x_1)$ and $f_{i-1}(x_2)$. The vertex $f_i(x)$ must coincide with y , and the uniqueness follows. For existence, well-definedness and locally injectiveness of f_i must be proved. For well-definedness, it is enough to prove that even if we choose x_2 and x_3 in place of x_1 and x_2 , we obtain the same y above. Clearly a tulip with petals x_1, x_2 , and x_3 is contained in $\langle H_{i-3} \cup H_{i-2} \cup H_{i-1} \rangle$. By locally injectiveness of f_{i-1} , this tulip is isomorphically mapped into G , and Condition 3 asserts that the existence of unique z which is adjacent to all of $f_{i-1}(x_1), f_{i-1}(x_2)$, and $f_{i-1}(x_3)$. It follows from Condition 2 that the above defined y coincides with this z independently of the choice of two of x_1, x_2 , and x_3 . Locally injectiveness straightforwardly follows from Condition 2. ■

Corollary 1. If G is a locally hamming graph, G is regular.

Proof. Let r be the maximum degree of G and take $u \in G$ with $\deg(u) = r$. Take the locally injection f of the proposition. The image of f in G is an r -regular subgraph of G , and the connectedness of G implies that the image equals G . ■

We call this r the *dimension* of G . It is clear that f in this proof is a covering. From now on, r is fixed; i.e., we consider only r -dimensional locally hamming graphs. Accordingly, locally injectiveness implies locally isomorphism and surjectiveness; that is, coveringness.

Corollary 2. (Universality of $f : H(r) \rightarrow G$)

Let $f : H(r) \rightarrow G$ be the above covering, H be an r -dimensional locally hamming graph, and $h : H \rightarrow G$ be a covering. Let w be a vertex of G , and take $v \in f^{-1}(w) \subset V(H(r))$ and $u \in h^{-1}(w) \subset V(H)$, arbitrarily. Then, there exists a unique covering $k : H(r) \rightarrow H$ such that $f = hk$ and $k(v) = u$.

Proof. Set $G' := H$ and $g' := (h|_{N(v)})^{-1} \circ (f|_{N(v)})$ and apply the proposition for $(G', u, H(r), v, g')$. Then we get $f' : H(r) \rightarrow G' = H$ such that $f'(v) = u$ and $f'|_{N(v)} = g'$. Let k be f' . Then, $(h \circ k)|_{N(v)} = h|_{N(u)} \circ k|_{N(v)} = h|_{N(u)} \circ g' = f|_{N(v)}$ implies $h \circ k = f$ by uniqueness. The uniqueness of such k follows from the fact that $k|_{N(v)}$ must be g' . ■

Corollary 3. (Uniqueness Theorem)

Let $g, h : H \rightarrow G$ be two coverings of locally hamming graphs. If $g|_{N(u) \cup \{u\}} = h|_{N(u) \cup \{u\}}$ for some $u \in H$, then $g = h$.

Proof. Take a covering $f : H(r) \rightarrow H$, and take a $v \in f^{-1}(u)$. Then $gf|_{N(v) \cup \{v\}} = hf|_{N(v) \cup \{v\}}$, and consequently, we have $gf = hf$ by Proposition 1, and since f is a surjection, we have $g = h$. ■

Let $h : H \rightarrow G$ be a covering. For $u \in G$, the set $h^{-1}(u)$ is called the *fiber* on u . The *Galois group* $\text{Gal}(H/G)$ of the covering $h : H \rightarrow G$ is defined to be $\{\gamma \in \text{Aut}(H) \mid h \circ \gamma = h\}$. It is obvious that if $v \in V(H)$ is a vertex in the fiber on u , then γv is again in the same fiber for any $\gamma \in \text{Gal}(H/G)$; i.e., $\text{Gal}(H/G)$ acts on every fiber. This action is easily proved to be faithful as follows. Suppose that $\gamma \in \text{Gal}(H/G)$ satisfies $\gamma(v) = v$ for some $v \in V(H)$. If $\gamma|_{N(v)} \neq \text{id}$, then there exists $s \in N(v)$ such that $\gamma(s) \neq s$, and since $h\gamma(s) = h(s)$, h is not locally injective. Thus, $\gamma|_{N(v)} = \text{id}$ holds, and $\gamma = \text{id}$ follows from Corollary 3.

A covering $h : H \rightarrow G$ is said to be *Galois* if $\text{Gal}(H/G)$ transitively acts on every fiber. This is an analogue of the Galois covering in the algebraic geometry[6].

Corollary 4. The universal covering $f : H(r) \rightarrow G$ is Galois.

Proof. This is a special case of Corollary 2 where $H = H(r)$. Note that a covering $h : H \rightarrow K$ is an isomorphism if $\#(V(H)) = \#(V(K))$.

2 Galois Groups

Let H be a locally hamming graph, and let Γ be a subgroup of $\text{Aut}(H)$. We define a graph H/Γ , which contains no multi-edges but may contain loops, and obtain a canonical projection $f : H \rightarrow H/\Gamma$. The vertex set of H/Γ is

$\{\Gamma u \mid u \in H\}$, where $\Gamma u = \{\gamma u \mid \gamma \in \Gamma\}$, and Γu is adjacent to Γv in H/Γ if and only if there exists a $\gamma \in \Gamma$ such that γu is adjacent to v in H . The canonical projection f is defined by $f : u \mapsto \Gamma u$. We define the *discreteness* d_Γ of Γ by

$$\min\{d(u, \gamma u) \mid \gamma \in \Gamma, \gamma \neq \text{id}, u \in V(H)\},$$

where d denotes the usual distance in the graph H . As usual, $d_{\{\text{id}\}}$ is defined to be ∞ .

Proposition 2. For a locally hamming graph H and a subgroup Γ of $\text{Aut}(H)$, H/Γ is a locally hamming graph if and only if $d_\Gamma \geq 5$. In this case, the canonical projection is a covering.

Proof. First, we prove the sufficiency. H/Γ contains a loop if and only if there exist two adjacent vertices $s, t \in \Gamma u$ for some $u \in V(H)$. This implies $\gamma s = t$ for some $\gamma \in \Gamma$; that is, $d_\Gamma \leq 1$. For the check of Conditions 1, 2, and 3, we use the next lemma.

Lemma 1. Let $\Gamma x_1, \Gamma x_2, \dots, \Gamma x_i$ be a walk in H/Γ ; i.e., Γx_j is adjacent to Γx_{j+1} for $j = 1, 2, \dots, i-1$. Then, there exist x'_2, \dots, x'_i such that $\Gamma x'_j = \Gamma x_j$ for $j = 2, \dots, i$ and that x_1, x'_2, \dots, x'_i is a walk in H . Moreover, if $\Gamma x_i = \Gamma x_1$, d_Γ , and $i \geq 5$, then $x'_i = x_1$.

Proof. Since Γx_1 is adjacent to Γx_2 , there exists an $x'_2 \in \Gamma x_2$ adjacent to x_1 . Thus, the existence of x'_j s is obvious. Suppose that $\Gamma x_i = \Gamma x_1$, that $i \leq 5$, and that $x'_i \neq x_1$. Then, since x_1, x'_2, \dots, x'_i is a path, $d(x_1, x'_i) \leq i-1 \leq 4$. Since x'_i is in $\Gamma x_i = \Gamma x_1$, $x'_i = \gamma x_1$ for some $\gamma \in \Gamma$, and consequently $d(x_1, \gamma x_1) \leq 4$, which is a contradiction. ■

Check of Condition 1. Suppose that H/Γ contains a triangle; in other words, that $i = 3$ holds in the latter half of Lemma 1. Then, Lemma 1 asserts that the triangle can be lifted up into H , a contradiction.

Check of Condition 2. Suppose that $\Gamma u, \Gamma w, \Gamma v$ is a path of length 2. Then, by Lemma 1, we may assume that uvw is also a path, by retaking w and v . Then, there exists an x such that $\langle u, w, v, x \rangle$ is a quadrilateral. It is easy to see that $\langle \Gamma u, \Gamma w, \Gamma v, \Gamma x \rangle$ is a quadrilateral. Suppose that all distinct $\Gamma u, \Gamma v, \Gamma w, \Gamma x, \Gamma y$ satisfy the condition that both $\langle \Gamma u, \Gamma w, \Gamma v, \Gamma x \rangle$

and $\langle \Gamma u, \Gamma w, \Gamma v, \Gamma y \rangle$ are quadrilaterals in H/Γ . Apply Lemma 1 on the $\langle \Gamma u, \Gamma w, \Gamma v, \Gamma x \rangle$, we get w', v' , and x' so that $\langle u, w', v', x' \rangle$ is a quadrilateral in H . Similarly we have y' so that $\langle u, w', v', y' \rangle$ is a quadrilateral in H , and since H is a locally hamming graph, $x' = y'$ holds; that is, $\Gamma x = \Gamma x' = \Gamma y' = \Gamma y$. This completes the check of Condition 2.

Check of Condition 3. Let T be a tulip in H/Γ . Using Lemma 1 on the three quadrilaterals in T , we can lift T up into H . Then, the existence of a vertex adjacent to all petals of T in H assures the existence of the one in H/Γ . This completes the proof of sufficiency.

For the necessity, it is enough to check the following easy statements. If $d_\Gamma = 0$, then there exists a $\gamma \neq \text{id}$ such that $\gamma(u) = u$. It follows from the r -regularity of H/Γ that $\gamma|_{N(u)}$ must be the identity function, and $\gamma = \text{id}$ follows from the uniqueness theorem. If $d_\Gamma = 1$, then H/Γ contains a loop. If $d_\Gamma = 2$, then H/Γ is not r -regular. If $d_\Gamma = 3$, then H/Γ contains a triangle. If $d_\Gamma = 4$, then H/Γ does not satisfy the uniqueness in Condition 2. ■

If $d_\Gamma \geq 5$, we call Γ a *discrete subgroup* of $\text{Aut}(H)$.

Lemma 2. Let Γ be a discrete subgroup of $\text{Aut}(H)$. Then the canonical projection $f : H \rightarrow H/\Gamma$ is Galois with $\text{Gal}(H/(H/\Gamma)) = \Gamma$.

Proof. Straightforward.

Lemma 3. Let G and H be locally hamming graphs and let $f : H \rightarrow G$ be a covering. Then there exists a unique covering $h : H/\text{Gal}(H/G) \rightarrow G$ such that $H \rightarrow H/\text{Gal}(H/G) \rightarrow G$ coincides with f . This covering h is an isomorphism if and only if $f : H \rightarrow G$ is Galois.

Proof. The covering h is defined by $h : \text{Gal}(H/G)u \mapsto fu$. The uniqueness follows from the surjectiveness of $H \rightarrow H/\text{Gal}(H/G)$. It is an isomorphism if and only if it is injective on the vertex sets; i.e., $f(u) = f(v)$ implies $\text{Gal}(H/G)u = \text{Gal}(H/G)v$. This is equivalent to saying that $\text{Gal}(H/G)$ transitively acts on every fibers; i.e., that $f : H \rightarrow G$ is Galois. ■

The Galois correspondence between coverings and discrete subgroups is as usual best described in terms of categorical framework. Let $f : H \rightarrow G$ be a Galois covering between two locally hamming graphs. We define a category $Sub(H/G)$ as follows. Its object set is the set of intermediate covering between H and G ; that is, $\{ \langle k, K, g \rangle \mid k : H \rightarrow K, g : K \rightarrow G : \text{coverings such that } f = gk \}$. Its arrow $l : \langle k, K, g \rangle \rightarrow \langle k', K', g' \rangle$ is a covering $l : K \rightarrow K'$ satisfying $k' = lk$ (and consequently $g = g'l$). Another category $Gal(H/G)$ is defined as follows. Its object set is the set of subgroups of $Gal(H/G)$. For $\Gamma, \Gamma' \in Gal(H/G)$, there exists at most one arrow $\Gamma \rightarrow \Gamma'$, and it exists if and only if $\Gamma \subset \sigma\Gamma'\sigma^{-1}$ for some $\sigma \in Gal(H/G)$. Obviously Γ and Γ' are isomorphic in $Gal(H/G)$ if and only if they are conjugate.

Theorem 1. For a Galois covering $f : H \rightarrow G$, $Sub(H/G)$ is categorically equivalent to $Gal(H/G)$ by

$$\begin{array}{ccc} Sub(H/G) & \cong & Gal(H/G) \\ \langle h, K, k \rangle & \mapsto & Gal(H/K) \\ H/\Gamma & \leftrightarrow & \Gamma. \end{array}$$

Moreover, the covering $k : H \rightarrow G$ is Galois if and only if $Gal(H/K)$ is a normal subgroup of $Gal(H/G)$, and in this case we have a canonical isomorphism $Gal(K/G) \cong Gal(H/G)/Gal(H/K)$.

We have another similar correspondence. For a locally hamming graph H , the category $Sub(H)$ is defined by: (i) the object set is $\{ \langle g, G \rangle \mid g : H \rightarrow G : \text{a Galois covering} \}$, (ii) an arrow $h : \langle g, G \rangle \rightarrow \langle g', G' \rangle$ is a covering $h : G \rightarrow G'$ with $g' = gh$. Also, the category $Daut(H)$ is defined by (i) the objects are the discrete subgroups of $Aut(H)$, (ii) there exists exactly one arrow $\Gamma \rightarrow \Gamma'$ if Γ is contained in some conjugate of Γ' in $Aut(H)$, and none exists otherwise.

Theorem 2. Let H be a locally hamming graph. Then, the categories $Sub(H)$ and $Daut(H)$ are categorically equivalent by

$$\begin{array}{ccc} Sub(H) & \cong & Daut(H) \\ \langle g, G \rangle & \mapsto & Gal(H/G) \\ H/\Gamma & \leftrightarrow & \Gamma. \end{array}$$

We can prove the above two theorems in exactly the same way with the case of Galois coverings in algebraic geometry (see [6]). The proofs are straightforward in the presence of Corollaries 1–4 of Proposition 1, Proposition 2, and Lemmas 2 and 3, and left to the readers as easy exercises. One may need an easy fact that if gh is a Galois covering then h is also. By setting $H := H(r)$ in Theorem 2, we have the next corollary.

Corollary 1. The isomorphic classes of r -dimensional hamming graphs are in 1-to-1 correspondence with the conjugate classes of discrete subgroups of $\text{Aut}(H(r))$.

Thus, the classification problem of locally hamming graphs is reduced to the one of discrete subgroups of $\text{Aut}(H(r))$. For an r -dimensional locally hamming graph G , the *fundamental group* $\pi_1(G)$ of G is defined to be $\text{Gal}(H(r)/G)$ by taking a covering $H(r) \rightarrow G$. In the next section, the structure of $\text{Aut}(H(r))$ is analyzed and examples of discrete subgroups are stated.

Remark. Any r -dimensional locally hamming graph G has a Galois covering $f : H(r) \rightarrow G$. For any $u \in V(G)$, $\pi_1(G) \cong \text{Gal}(H(r)/G)$ transitively and faithfully acts on every fiber $f^{-1}(u)$, and consequently $\#(f^{-1}(u)) = \#(\pi_1(G))$. Since $2^r = \#(V(H(r))) = \sum_{u \in V(G)} \#(f^{-1}(u)) = \#(V(G)) \times \#(\pi_1(G))$, we have that both $\#(V(G))$ and $\#(\pi_1(G))$ are power of 2. It follows that every group appeared so far is a 2-group.

3 Examples

Proposition 2. (The structure of $\text{Aut}(H(r))$)

The group $\text{Aut}(H(r))$ is isomorphic to the wreath product $\mathbf{F}_2 \text{wr} \mathcal{S}_r$.

Proof. The wreath product $\mathbf{F}_2 \text{wr} \mathcal{S}_r$ is by definition (i) as a set, it is $\mathcal{S}_r \times \mathbf{F}_2^r$, where \mathcal{S}_r is considered as the group of permutation matrices of size r over \mathbf{F}_2 , (ii) the operation \cdot is defined by $\langle \sigma, d \rangle \cdot \langle \sigma', d' \rangle = \langle \sigma\sigma', \sigma d' + d \rangle$ for $\sigma, \sigma' \in \mathcal{S}_r$ and $d, d' \in \mathbf{F}_2^r$.

We construct a homomorphism $\Phi : \mathbf{F}_2 \text{wr} \mathcal{S}_r \rightarrow \text{Aut}(H(r))$ by $\langle \sigma, d \rangle \mapsto \phi : \phi(z) = \sigma z + d$ for any $z \in \mathbf{F}_2^r$. It is easy to show that ϕ is in fact in $\text{Aut}(H(r))$. The kernel of Φ is easily proved to be $\{\text{id}\}$. To prove the surjectiveness, take an arbitrary $\gamma \in \text{Aut}(H(r))$. Put $d := \gamma(0)$ and define

$\sigma \in \mathcal{S}_r$ so that $\sigma(e_i) = \gamma(e_i) - \gamma(0)$ holds, where e_i denotes the i -th unit vector whose components are all zero except for the i -th. We can easily check that $\Phi(\langle \sigma, d \rangle) |_{N(0) \cup \{0\}} = \gamma |_{N(0) \cup \{0\}}$, and $\Phi(\langle \sigma, d \rangle) = \gamma$ follows from Uniqueness Theorem. ■

From now on, we identify $\text{Aut}(H(r))$ with $\mathbf{F}_2 \text{wr} \mathcal{S}_r$, and denote any element in $\text{Aut}(H(r))$ in the form of $\sigma z + d$. In case that $\sigma = \text{id}$, it is denoted by d . In the rest of this section, we enumerate some examples of discrete subgraph $\Gamma \subset \text{Aut}(H(r))$.

Example 1. Distance-regular case.

$\Gamma = \langle \mathbf{1} \rangle$, where $\mathbf{1}$ is the vector whose entries are all 1, and $\langle \rangle$ denotes the generated group. Γ is discrete if and only if $r \geq 5$. It is known that $H(r)/\Gamma$ is a distance-regular graph. The next conjecture is a paraphrase of a well-known conjecture.

Conjecture. If $H(r)/\Gamma$ is a locally hamming distance-regular graph, then $\Gamma = \langle \mathbf{1} \rangle$ and $r \geq 5$, or $\Gamma = \{\text{id}\}$.

Rifa[5] and Nomura[7] settled some special cases of this conjecture using code theory. For the relation between completely regular codes and distance-regular graphs, see [3].

Example 2. Linear codes. Suppose that $\Gamma = \langle d_1, d_2, \dots, d_i \rangle$. In this case, we can identify Γ with a subspace in \mathbf{F}_2^r . A vector space $V \subset \mathbf{F}_2^r$ is said to be k -error correcting linear code if $\min\{d(u, v) \mid u, v \in V, u \neq v\} = 2k + 1$. Thus, Γ is discrete if and only if Γ is a 2-error correcting linear code. (There are many kinds of such codes. See [2].)

Example 3. The case $\pi_1(G)$ is non-abelian. Suppose that $\Gamma \subset \text{Aut}(H(r))$, and let s be a positive integer. Then, it is clear that the s -ary cartesian products of Γ , $\Gamma \times \Gamma \times \dots \times \Gamma$, acts on $H(r \times s)$, for a positive integer s . Take the diagonal subgroup Δ of $\Gamma \times \dots \times \Gamma$. It is isomorphic to Γ , and clearly $d_\Delta = s \cdot d_\Gamma$ holds. Thus, if one has a Γ with $d_\Gamma > 0$; in other words, a $\Gamma \subset \text{Aut}(H(r))$ which has no fixed point, then there is a $\Delta \subset \text{Aut}(H(5r))$ such that $\Gamma \cong \Delta$ and that $H(5r)/\Delta$ is a locally hamming graph. Let D_8 be the group generated by a, b satisfying the equation $ba = a^3b$; i.e., the fourth dihedral group. This is the smallest non-abelian 2-group. It is clear

that D_8 is isomorphic to $\text{Aut}(H(2))$ and acts on $H(2)$ or \mathbb{F}_2^2 . Let D_8 act on $H(3)$ or $\mathbb{F}_2^3 = \mathbb{F}_2^2 \oplus \mathbb{F}_2$ by setting $\sigma \cdot \langle d, s \rangle := \langle \sigma \cdot d, s + \text{sign}(\sigma) \rangle$ for $\sigma \in D_8$, $d \in \mathbb{F}_2^2$, and $s \in \mathbb{F}_2$, where $\text{sign}(\sigma)$ is defined to be 0 if $\sigma = \text{id}$, a , a^2 , or a^3 , and to be 1 otherwise. With no difficulty we can check that D_8 is in fact a subgroup of $\text{Aut}(H(3))$ without fixed point. Thus, there is a $\Delta \subset \text{Aut}(H(15))$ isomorphic to D_8 such that $H(15)/\Delta$ is a locally hamming graph.

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Nearly triply regular symmetric designs

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1. $D = (P, B)$ を 2 - (ν, k, λ) 対称 \bar{t} - ガイン とする. ここで P, B はそれぞれ D の点およびブロックの集合を示す. 任意の相異なる $\alpha, \beta, \gamma \in B$ に対し $|\alpha \cap \beta \cap \gamma|$ がちょうど ν の値 μ, ν ($0 \leq \nu < \mu \leq \lambda$) を取るから, D は nearly triply regular (NTR) である. 次の事実はすぐに確かめられる.

(1) 上で $|\alpha \cap \beta \cap \gamma|$ が constant とすると, D は自明な \bar{t} - ガイン , 即ち $k = \nu - 1$ とする. 逆もよい. それで以下 D は自明でないとする.

(2) $\lambda = 1, 2$ ならば何時でも NTR である. それで以下 $\lambda \geq 3$ とする.

(3) D が NTR ならば D の complement も NTR である. 逆もよい. それで以下

$v \geq 2n + 1$ とする.

2. NTR とは 概念は 最初 Herzog-Reid [1] によつて, D の 特殊な アダマール デイザイン, 即ち アダマール トーナメント のとき, 有向グラフ の 言葉 を 使って 導入 された 様 に 思 はれる. 彼等 の 主要 な 結果 は $v=0$ から $\lambda=1, 2$ に なる と いう こ と で あつた. 然し NTR は 対称 デイザイン について の 概念 である. 著者 の 人は 最初に アダマール デイザイン の 時 を 考察 した. $\mu = \lambda$ の ときは $GF(2)$ 上 の 射影 幾何 型 の アダマール デイザイン として 特徴 付け られる. ~~相違~~ $v=0$ 以下 $\lambda > \mu$ と する.

(1) $\lambda \{ (v-2)(\mu+1) - (\lambda-1)(\lambda-2) \} = \mu v (v-2)$
という 容易 に 得 られる 関係 式 から 出発 する.

$v=0$ から

$$(2) \quad (v-2)(\mu-1) = (\lambda-1)(\lambda-2)$$

が 得 られる. アダマール デイザイン の とき $v=2\lambda+1$ である から, $v=0$ から $\lambda=1, 2$ は 自明 である. それ 以外 の Herzog-Reid の 主要 な 結果 の 簡単 な 証明 と なる. そこで $v > 0$ と する. (1) から, $v=4\lambda+3$ である から, $\mu v = a\lambda$, a は 正 整数, と 表わ せる.

$$\lambda > \mu > \nu > a \text{ のとき } 12a \leq 12\lambda - 33$$

こゝから (1) から

$$(3) \quad 12a + 3 = (2\lambda - 1)(4\mu + 4\nu + 1 - 8a - 2\lambda)$$

$$\text{そこで } 4\mu + 4\nu + 1 - 8a - 2\lambda = 1, 3 \text{ または } 5$$

が得られる。=5 のときは $6a = 5\lambda - 4$, したがって

$$6\mu\nu = (5\lambda - 4)\lambda, \quad 6\mu + 6\nu = 13\lambda - 2$$

と得るが, これから $\mu > \lambda$ が出てしまう, =3 の

$$\text{ときは, } 2a = \lambda - 1, \quad 2\mu\nu = (\lambda - 1)\lambda, \quad 2\mu + 2\nu$$

$$= 3\lambda - 1 \text{ と得るが, これから } \mu = \lambda \text{ が出てしまう}$$

$$\text{それで, } =1 \text{ とする. このとき } 6a = \lambda - 2, \quad 6\mu\nu =$$

$$(\lambda - 2)\lambda, \quad 6\mu + 6\nu = 5\lambda - 4 \text{ である. これから}$$

$$\mu = \frac{\lambda}{2}, \quad \nu = \frac{\lambda - 2}{3} \text{ が得られる. } \alpha, \beta \text{ を固定した}$$

とき, $|\alpha \cap \beta \cap \gamma| = \nu$ とする γ の個数を d とすると,

$$\text{このとき } d = 9 - \frac{36}{\lambda + 4} \text{ が得られる. ところで}$$

$$\lambda = 8, 14, 32, \text{ 然し } |\alpha \cap \beta \cap \gamma| = \nu \text{ とする } \gamma \text{ 才}$$

$$\{\alpha, \beta, \gamma\} \text{ 全部の個数 } t = \frac{\binom{\lambda+3}{2} d}{3} \text{ である.}$$

$$\lambda = 14, 32 \text{ のとき } t \text{ 整数とならないうので, } \lambda = 8$$

$$\text{を得る. この時は } d = 6, \quad t = 1190 \text{ となり, 数論}$$

的矛盾を得るのは困難である. 然し実際に

構成しようとしてみると, 存在する. これは比較的

簡単にはわかる ([5] 参照)

3. 一般の対称 Γ "サイン" について考察したとき、
その人の著者は、 $\mu = \lambda$ という假定の下では
NTR は Dembowski-Wagner の smooth という
概念の dual ということに気付いた。それで
 $\mu = \lambda$ のときは $GF(q)$ 上の射影幾何型の Γ "サイン" として特徴付けられた。

4. それでまた以下 $\lambda > \mu$ とする。著者は
NTR の全体像を未だ推測出来ていない。
例えば $\nu = 0$ なら (2) があるのであるが、
はっきりわかりにくい。それで アタマール Γ "サイン"
の次の Γ "サイン" 類として $v = 4(q - \lambda)$ (アタマール
は $v = 4(q - \lambda) - 1$ の時がある) のとき、
即ち正則 アタマール 行列型 (RH) のときを
考察してみた。このとき $v = 4m^2$, $q = 2m^2 - m$,
 $\lambda = m^2 - m$ とパラメタライズされる。もし NTR
があるとすると $\mu = \frac{m^2 - m}{2}$, $\nu = \frac{m^2 - 2m}{2}$ となる
ことだ。そこで アタマール Γ "サイン" の時に述べた
様子を仕方でわかる。とくに m は偶数となる。
それで無限列を構成することを考えた。もとより

普通のRH型デザインは $A_0 = \begin{pmatrix} 1 & -1 & -1 & -1 \\ -1 & 1 & -1 & -1 \\ -1 & -1 & 1 & -1 \\ -1 & -1 & -1 & 1 \end{pmatrix}$,

$$B_0 = -A_0, \quad A_{i+1} = \begin{pmatrix} A_i & B_i & B_i & B_i \\ B_i & A_i & B_i & B_i \\ B_i & B_i & A_i & B_i \\ B_i & B_i & B_i & A_i \end{pmatrix}, \quad B_{i+1} = -A_{i+1}$$

($i=0, 1, 2, \dots$) とするとき $C_{2l-1} = B_{2l-1} c^{-1} \rightarrow 0$

としたもの, $C_{2l} = A_{2l} c^{-1} \rightarrow 0$ として得られる

C_1, C_2, C_3, \dots と "5-リ-ク" c がある. このとき

$m = 2^n$ ($n=1, 2, 3, \dots$). 考えられるものは

NPR であることが induction で示される.

然し, RH型デザインで NPR であるものが

あることは, $m=4$ のときに (RH型デザインは

m をきめるとき 非常に沢山の "否" であることが知られて

いる) 示される.

さらに $v=4(n-1)+1$ という形のときは NPR であることが示される.

5. NPR の dual は NPR か? という様な問

にも答が得られていない. 著者達もはじめは

ばかりで, 文献なども充分に check されて

いない. この教示を待つ次第です.

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Products of Hadamard Matrices, Williamson Matrices and Other Orthogonal Matrices using M-Structures

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Abstract

The new concept of M-structures is used to unify and generalize a number of concepts in Hadamard matrices including Williamson matrices, Goethals-Seidel matrices, Wallis-Whiteman matrices and generalized quaternion matrices. The concept is used to find many new symmetric Williamson-type matrices, both in sets of four and eight, and many new Hadamard matrices. We give as corollaries “that the existence of Hadamard matrices of orders $4g$ and $4h$ implies the existence of an Hadamard matrix of order $8gh$ ” and “the existence of Williamson type matrices of orders u and v implies the existence of Williamson type matrices of order $2uv$ ”. This work generalizes and utilizes the work of Masahiko Miyamoto and Mieko Yamada. Lists of odd orders < 1000 for which Hadamard and Williamson type matrices are known are given.

1 Definitions and Introduction

For the definitions used in this paper, and for detailed proofs, we refer the reader to [47].

2 M-structures

An orthogonal matrix of order $4t$ can be divided into sixteen (16) $t \times t$ blocks M_{ij} . This partitioned matrix is said to be an M-structure. If the orthogonal matrix can be partitioned into sixty-four (64) $s \times s$ blocks M_{ij} it will be called a 64 block M-structure.

An Hadamard matrix made from (symmetric) Williamson matrices W_1, W_2, W_3, W_4 is an M-structure with

$$\begin{aligned} W_1 &= M_{11} = M_{22} = M_{33} = M_{44}, \\ W_2 &= M_{12} = -M_{21} = M_{34} = -M_{43}, \\ W_3 &= M_{13} = -M_{31} = -M_{24} = M_{42}, \text{ and} \\ W_4 &= M_{14} = -M_{41} = M_{23} = -M_{32}. \end{aligned}$$

An Hadamard matrix made from four (4) circulant (or type 1) matrices A_1, A_2, A_3, A_4 of order n , where R is the matrix which makes all the $A_i R$ back-circulant (or type 2), is an M-structure with

$$\begin{aligned} A_1 &= M_{11} = M_{22} = M_{33} = M_{44}, \\ A_2 &= M_{12}R = -M_{21}R = RM_{34}^T = -RM_{43}^T, \\ A_3 &= M_{13}R = -M_{31}R = -RM_{24}^T = RM_{42}^T, \text{ and} \\ A_4 &= M_{14}R = -M_{41}R = RM_{23}^T = -RM_{32}^T. \end{aligned}$$

The next theorem and corollary are easy to prove using M-structures.

Theorem 1 *Suppose there are T -matrices of order t . Further suppose there is an $OD(4s; u_1, \dots, u_n)$ constructed of sixteen circulant (or type 1) $s \times s$ blocks on the variables x_1, \dots, x_n . Then there is an $OD(4st; tu_1, \dots, tu_n)$. In particular if there is an $OD(4s; s, s, s, s)$ constructed of sixteen circulant (or type 1) $s \times s$ blocks then there is an $OD(4st; st, st, st, st)$.*

Corollary 2 *Suppose the T -matrices are of order t . Then there are orthogonal designs $OD(20t; 5t, 5t, 5t, 5t)$ and $OD(36t; 9t, 9t, 9t, 9t)$.*

Conjecture 3 *There exists an $OD(4t; t, t, t, t)$ for every positive integer t .*

We also conjecture

Conjecture 4 *There exists an M-structure $OD(4t; t, t, t, t)$ for every $t \equiv 1 \pmod{4}$ comprising sixteen circulant or type 1 blocks.*

3 Some properties of certain amicable orthogonal matrices

We use the following three lemmas proved in [47].

Lemma 5 Suppose there exist two amicable $(0, +1, -1)$ matrices U, V of order u satisfying $UU^T + VV^T = (2u - 1)I$. Then there exist matrices A, B, D of order u satisfying

$$\begin{aligned} AA^T + BB^T &= B^T B + D^T D = (2u - 1)I \\ A^T &= (-1)^{\frac{1}{2}(u-1)}A, D^T = (-1)^{\frac{1}{2}(u-1)}D, \end{aligned}$$

where A and D have zero diagonal.

Lemma 6 Let $q + 1$ be the order of a conference matrix. Then there exist four matrices C_1, C_2, C_3, C_4 , of order $\frac{1}{2}(q - 1)$ satisfying

$$\begin{aligned} C_1 C_1^T + C_2 C_2^T &= C_3 C_3^T + C_4 C_4^T = qI - 2J, \\ e C_1^T &= e C_4^T = e, \quad e C_2^T = e C_3^T = 0, \\ C_1 C_3^T - C_2 C_4^T &= 0, \quad C_1^T = C_1, \quad C_4^T = C_4, \quad C_3^T = C_2, \end{aligned}$$

where e is the $1 \times \frac{1}{2}(q - 1)$ matrix of ones, C_1 and C_4 have zero diagonal elements ± 1 , C_2 and C_3 have elements ± 1 .

Lemma 7 Suppose there exist two amicable $(0, +1, -1)$ matrices U, V of order u satisfying $UU^T + VV^T = (2u - 1)I$. Further suppose U has zero diagonal and U, V have other elements $+1$ or -1 . Then there exist matrices A, B of order $u - 1$ satisfying

$$\begin{aligned} AA^T + BB^T &= (2u - 1)I_{u-1} - 2J_{u-1}, \\ eA^T &= e, \quad eB^T = 0, \quad AB^T = BA^T, \end{aligned}$$

where A has one zero element per row and column and the other entries of A and B are ± 1 . Further if U and V are symmetric (or skew-type respectively) then A and B are symmetric (or skew-type respectively).

Furthermore if U and V satisfy $UU^T + VV^T = 2uI$ (U, V are $(1, -1)$ matrices), u even, then there exist matrices A, B of order $u - 1$, with entries ± 1 , satisfying

$$\begin{aligned} AA^T + BB^T &= 2uI_{u-1} - 2J_{u-1}, \\ eA^T &= e, \quad eB^T = e, \quad AB^T = BA^T, \end{aligned}$$

and if U and V are symmetric (or skew-type respectively) then A and B are symmetric (or skew-type respectively).

4 A multiplication Theorem using M-structures

Theorem 8 Let $N = (N_{ij})$, $i, j = 1, 2, 3, 4$ be an Hadamard matrix of order $4n$ of M-structure. Further let T_{ij} , $i, j = 1, 2, 3, 4$ be 16 $(0, +1, -1)$ type 1 or circulant matrices of order t which satisfy

- (i) $T_{ij} * T_{ik} = 0, T_{ji} * T_{ki} = 0, j \neq k, (* \text{ the Hadamard product})$
(ii) $\sum_{k=1}^4 T_{ik}$ is a $(1, -1)$ matrix,
(iii) $\sum_{k=1}^4 T_{ik} T_{ik}^T = tI_t = \sum_{k=1}^4 T_{ki} T_{ki}^T,$
(iv) $\sum_{k=1}^4 T_{ik} T_{jk}^T = 0 = \sum_{k=1}^4 T_{ki} T_{kj}^T, i \neq j.$

Then there is an M -structure Hadamard matrix of order $4nt$.

Corollary 9 *If there exists an Hadamard matrix of order $4h$ and an orthogonal design $OD(4u; u_1, u_2, u_3, u_4)$, then an $OD(8hu; 2hu_1, 2hu_2, 2hu_3, 2hu_4)$ exists.*

Corollary 10 *If there exists an Hadamard matrix of order $4h$ and an orthogonal design $OD(4u; u, u, u, u)$, then there exists an $OD(8hu; 2hu, 2hu, 2hu, 2hu)$.*

This gives the theorem of Agayan and Sarukhanyan [2] as a corollary by setting all variables equal to one:

Corollary 11 *If there exists Hadamard matrices of orders $4h$ and $4u$ then there exists an Hadamard matrix of order $8hu$.*

We now give as a corollary a result, motivated by, and a little stronger than that of Agayan and Sarukhanyan [2]:

Corollary 12 *Suppose there are Williamson or Williamson type matrices of orders u and v . Then there are Williamson type matrices of order $2uv$.*

If the matrices of orders u and v are symmetric the matrices of order $2uv$ are also symmetric.

If the matrices of orders u and v are circulant and/or type 1 the matrices of order $2uv$ are type 1.

5 Miyamoto's Theorem and Corollaries via M -structures

We reformulate Miyamoto's results so that symmetric Williamson-type matrices can be obtained.

Lemma 13 (Miyamoto's Lemma Reformulated) *Let $U_i, V_j, i, j = 1, 2, 3, 4$ be $(0, +1, -1)$ matrices of order n which satisfy*

- (i) $U_i, U_j, i \neq j$ are pairwise amicable,
(ii) $V_i, V_j, i \neq j$ are pairwise amicable,

- (iii) $U_i \pm V_i$, $(+1, -1)$ matrices, $i = 1, 2, 3, 4$,
- (iv) the row sum of U_1 is 1, and the row sum of U_j , $i = 2, 3, 4$ is zero,
- (v) $\sum_{i=1}^4 U_i U_i^T = (2n + 1)I - 2J$, $\sum_{i=1}^4 V_i V_i^T = (2n + 1)I$.

Then there are 4 Williamson type matrices of order $2n + 1$. If U_i and V_i are symmetric, $i = 1, 2, 3, 4$ then the Williamson-type matrices are symmetric. Hence there is a Williamson type Hadamard matrix of order $4(2n + 1)$.

Corollary 14 Let $q \equiv 1 \pmod{4}$ be a prime power then there are symmetric Williamson type matrices of order $q + 2$ whenever $\frac{1}{2}(q + 1)$ is a prime power or $\frac{1}{2}(q + 3)$ is the order of a symmetric conference matrix. Also there exists an Hadamard matrix of Williamson type of order $4(q + 2)$.

Remark 15 Some of the results in Corollary 14 are also due to A.L. Whiteman [35]. This gives symmetric Williamson-type matrices of orders

| | | | | | | | | | |
|------|------|------|------|------|------|------|------|------|------|
| 7 | 11 | 15 | 19 | 27 | 39 | 51 | 55 | 63 | 75 |
| 83 | 91 | 99 | 123 | 159 | 195 | 243 | 279 | 315 | 339 |
| 363 | 399 | 423 | 451 | 459 | 543 | 579 | 615 | 627 | 663 |
| 675 | 735 | 759 | 843 | 879 | 883 | 999 | 1095 | 1155 | 1203 |
| 1215 | 1239 | 1251 | 1323 | 1383 | 1455 | 1623 | 1659 | 1683 | 1755 |
| 1875 | 1935 | 1995 | | | | | | | |

(since Mathon found conference matrices of orders 46 and 442). Almost all these, with symmetry, are new though Miyamoto [12] has found Williamson-type matrices for these orders and hence Hadamard matrices for four times these orders.

Koukouvinos and Kounias [10] have shown there are no circulant symmetric Williamson matrices of order 39 but here a symmetric but not circulant Williamson matrix of order 39 is given.

Corollary 16 Let $q \equiv 1 \pmod{4}$ be a prime power. Then

- (i) if there are Williamson type matrices of order $(q - 1)/4$ or an Hadamard matrix of order $\frac{1}{2}(q - 1)$ there exist Williamson type matrices of order q ;
- (ii) if there exist symmetric conference matrices of order $\frac{1}{2}(q - 1)$ or a symmetric Hadamard matrix of order $\frac{1}{2}(q - 1)$ then there exist symmetric Williamson type matrices of order q .

Hence there exists an Hadamard matrix of Williamson type of order $4q$.

Remark 17 Part (i) of Corollary 16 for Williamson matrices of order $(q-1)/4$ was found by Miyamoto [12]. Part (i) with Hadamard matrices of order $\frac{1}{2}(q-1)$ is new. Part (ii) with symmetry is new.

Corollary 16 (ii) gives symmetric Williamson-type matrices of order q when $q \equiv 1 \pmod{4}$ is a prime power and $\frac{1}{2}(q-1)$ is the order of a symmetric conference matrix. This gives symmetric Williamson-type matrices for the following orders:

| | | | | | | | | | |
|------|------|------|------|------|------|------|------|------|------|
| 13 | 29 | 37 | 53 | 61 | 101 | 109 | 125 | 149 | 181 |
| 197 | 229 | 277 | 317 | 349 | 389 | 397 | 461 | 541 | 557 |
| 677 | 701 | 709 | 797 | 821 | 1021 | 1061 | 1117 | 1229 | 1237 |
| 1549 | 1597 | 1621 | 1709 | 1861 | 1877 | 1997 | | | |

Corollary 16 part (ii) gives symmetric Williamson-type matrices of order q when $q \equiv 1 \pmod{4}$ is a prime power and $\frac{1}{2}(q-1)$ is the order of a symmetric Hadamard matrix. Remembering that symmetric Hadamard matrices exist for orders $p+1$ when $p \equiv 3 \pmod{4}$ is a prime power we have symmetric Williamson-type matrices for the following orders:

| | | | | | | | | | |
|------|------|------|------|------|------|------|------|------|------|
| 5 | 9 | 17 | 25 | 41 | 49 | 73 | 81 | 89 | 97 |
| 113 | 121 | 169 | 193 | 241 | 257 | 281 | 289 | 337 | 353 |
| 361 | 401 | 409 | 433 | 449 | 457 | 529 | 569 | 577 | 593 |
| 601 | 617 | 625 | 641 | 673 | 729 | 761 | 769 | 841 | 881 |
| 929 | 937 | 961 | 977 | 1009 | 1033 | 1049 | 1097 | 1129 | 1153 |
| 1201 | 1217 | 1249 | 1289 | 1297 | 1321 | 1361 | 1369 | 1409 | 1481 |
| 1489 | 1553 | 1601 | 1609 | 1657 | 1681 | 1697 | 1721 | 1777 | 1801 |
| 1849 | 1873 | | | | | | | | |

Corollary 16 part (i) gives Williamson-type matrices of order q when $q \equiv 1 \pmod{4}$ is a prime power and $\frac{1}{2}(q-1)$ is the order of an Hadamard matrix. This gives Williamson-type matrices for the following orders not given above:

| | | | | | | | | | |
|-----|-----|-----|-----|-----|-----|------|------|------|------|
| 137 | 233 | 313 | 521 | 809 | 953 | 1193 | 1753 | 1889 | 1993 |
|-----|-----|-----|-----|-----|-----|------|------|------|------|

Corollary 16 part (i) gives Williamson-type matrices of order q when $q \equiv 1 \pmod{4}$ is a prime power and $(q-1)/4$ is the order of Williamson-type matrices. This result is also due to Miyamoto [12]. This gives Williamson-type matrices for the following orders:

| | | | | | | | | | |
|------|------|------|------|------|------|-----|-----|------|------|
| 157 | 173 | 293 | 373 | 613 | 757 | 757 | 773 | 1109 | 1301 |
| 1453 | 1493 | 1637 | 1693 | 1733 | 1741 | | | | |

Corollary 18 Let $q \equiv 1 \pmod{4}$ be a prime power or $q+1$ the order of a symmetric conference matrix. Let $2q-1$ be a prime power. Then there exist

symmetric Williamson type matrices of order $2q + 1$ and an Hadamard matrix of Williamson type of order $4(2q + 1)$.

Remark 19 Corollary 18 is satisfied for the appropriate primes or conference matrix orders to give symmetric Williamson-type matrices for the following orders:

| | | | | | | | | | |
|------|------|------|------|------|------|------|------|------|------|
| 11 | 19 | 27 | 51 | 75 | 83 | 91 | 99 | 123 | 195 |
| 243 | 315 | 339 | 363 | 451 | 459 | 579 | 627 | 675 | 843 |
| 883 | 1155 | 1203 | 1251 | 1323 | 1659 | 1683 | 1755 | 1875 | 1995 |
| 2019 | 2139 | 2403 | 2475 | 2595 | 2859 | 3043 | 3219 | 3315 | 3363 |
| 3483 | 3699 | 3723 | | | | | | | |

Note this last corollary is a modified version of Miyamoto's Corollary 5 (original manuscript). A new proof of Miyamoto's result, preserving symmetry, gives:

Corollary 20 Let $q \equiv 5 \pmod{8}$ be a prime power. Further let $\frac{1}{2}(q - 3)$ be a prime power or $\frac{1}{2}(q - 1)$ be the order of a symmetric conference matrix then there exist symmetric Williamson type matrices of order q and an Hadamard matrix of Williamson type of order $4q$.

Theorem 21 (Miyamoto's Theorem Reformulated) Let $U_{ij}, V_{ij}, i, j = 1, 2, 3, 4$ be $(0, +1, -1)$ matrices of order n which satisfy

- (i) $U_{ki}, U_{kj}, i \neq j$ are pairwise amicable, $k = 1, 2, 3, 4$,
- (ii) $V_{ki}, V_{kj}, i \neq j$ are pairwise amicable, $k = 1, 2, 3, 4$,
- (iii) $U_{ki} \pm V_{ki}, (+1, -1)$ matrices, $i, k = 1, 2, 3, 4$,
- (iv) the row sum of U_{ii} is 1, and the row sum of U_{ij} is zero, $i \neq j, i, j = 1, 2, 3, 4$,
- (v) $\sum_{i=1}^4 U_{ji}U_{ji}^T = (2n + 1)I - 2J, \sum_{i=1}^4 V_{ji}V_{ji}^T = (2n + 1)I, j = 1, 2, 3, 4$,
- (vi) $\sum_{i=1}^4 U_{ji}U_{ki}^T = 0, \sum_{i=1}^4 V_{ji}V_{ki}^T = 0, j \neq k, j, k = 1, 2, 3, 4$.

If conditions (i) to (v) hold, there are four Williamson matrices type of order $2n + 1$ and thus a Williamson type Hadamard matrix of order $4(2n + 1)$. Furthermore if the matrices U_{ki} and V_{ki} are symmetric for all $i, j = 1, 2, 3, 4$ the Williamson matrices obtained of order $2n + 1$ are also symmetric.

If conditons (iii) to (vi) hold, there is an M -structure Hadamard matrix of order $4(2n + 1)$.

Proof: Use

$$\begin{aligned}
 X_{11} &= \begin{bmatrix} -1 & -e \\ -e^T & S_{11} \end{bmatrix} & X_{12} &= \begin{bmatrix} 1 & e \\ e^T & S_{12} \end{bmatrix} & X_{13} &= \begin{bmatrix} 1 & e \\ e^T & S_{13} \end{bmatrix} & X_{14} &= \begin{bmatrix} -1 & e \\ e^T & S_{14} \end{bmatrix} \\
 X_{21} &= \begin{bmatrix} 1 & e \\ e^T & S_{21} \end{bmatrix} & X_{22} &= \begin{bmatrix} -1 & -e \\ -e^T & S_{22} \end{bmatrix} & X_{23} &= \begin{bmatrix} 1 & e \\ e^T & S_{23} \end{bmatrix} & X_{24} &= \begin{bmatrix} -1 & e \\ e^T & S_{24} \end{bmatrix} \\
 X_{31} &= \begin{bmatrix} 1 & e \\ e^T & S_{31} \end{bmatrix} & X_{32} &= \begin{bmatrix} 1 & e \\ e^T & S_{32} \end{bmatrix} & X_{33} &= \begin{bmatrix} -1 & -e \\ -e^T & S_{33} \end{bmatrix} & X_{34} &= \begin{bmatrix} -1 & e \\ e^T & S_{34} \end{bmatrix} \\
 X_{41} &= \begin{bmatrix} -1 & -e \\ e^T & -S_{41} \end{bmatrix} & X_{42} &= \begin{bmatrix} 1 & e \\ e^T & -S_{42} \end{bmatrix} & X_{43} &= \begin{bmatrix} -1 & -e \\ e^T & -S_{43} \end{bmatrix} & X_{44} &= \begin{bmatrix} -1 & -e \\ -e^T & -S_{44} \end{bmatrix}
 \end{aligned}$$

We note that the following always holds as it is just a case of Miyamoto's Lemma Reformulated:

$$\sum_{i=1}^4 S_{ji} S_{ji}^T = 4(2n+1)I_{2n} - 4J_{2n}. \quad (2)$$

In all cases though assumption (vi) assures us that

$$\sum_{i=1}^4 S_{ki} S_{ji}^T = 0, \quad j \neq k. \quad (3)$$

Note that if we write our M-structure from the theorem as

$$\begin{array}{cccccccc}
 -1 & 1 & 1 & -1 & -e & e & e & e \\
 1 & -1 & 1 & -1 & e & -e & e & e \\
 1 & 1 & -1 & -1 & e & e & -e & e \\
 1 & 1 & 1 & 1 & -e & -e & -e & e \\
 -e^T & e^T & e^T & e^T & S_{11} & S_{12} & S_{13} & S_{14} \\
 e^T & -e^T & e^T & e^T & S_{21} & S_{22} & S_{23} & S_{24} \\
 e^T & e^T & -e^T & e^T & S_{31} & S_{32} & S_{33} & S_{34} \\
 -e^T & -e^T & -e^T & e^T & S_{41} & S_{42} & S_{43} & S_{44}
 \end{array}$$

and we can see Yamada's matrix with trimming [46] or the J. Wallis-Whiteman [30] matrix with a border embodied in the construction.

Corollary 22 *Suppose there exists a symmetric conference matrix of order $q+1 = 4t+2$ and an Hadamard matrix of order $4t = q-1$. Then there is an Hadamard matrix with M-structure of order $4(4t+1) = 4q$. Further if the Hadamard matrix is symmetric the Hadamard matrix of order $4q$ is of the form*

$$\begin{bmatrix} X & Y \\ -Y & X \end{bmatrix},$$

where X, Y are amicable and symmetric.

We note that complex Hadamard matrices of order $n \equiv 2 \pmod{4}$ do exist when symmetric conference matrices cannot exist (see [22, Chapter VI]). These complex Hadamard matrices may be written as $K = X + iY$ where $KK^* = kI_k$ (* the Hermitian conjugate).

Hence we have

Corollary 23 *Let $q \equiv 4f + 1$ be a prime power. Suppose there is a complex Hadamard matrix of order $2f$. Then there is an Hadamard matrix of order $4(4f + 1)$.*

Note complex Hadamard matrices exist for orders 22, 34, 58, 86, 306, 650, 870, 1046, 2450, 3782, ..., for which either a symmetric conference matrix cannot exist or is not known. None of these orders give new Hadamard matrices.

6 Using 64 Block M-structures

In a similar fashion, we consider the following lemma so symmetric 8-Williamson-type matrices can be obtained.

Lemma 24 *Let $U_i, V_j, i, j = 1, \dots, 8$ be $(0, +1, -1)$ matrices of order n which satisfy*

- (i) $U_i, U_j, i \neq j$ are pairwise amicable,
- (ii) $V_i, V_j, i \neq j$ are pairwise amicable,
- (iii) $U_i \pm V_i, (+1, -1)$ matrices, $i = 1, \dots, 8$,
- (iv) the row(column) sums of U_1 and U_2 are both 1, and the row sum of $U_j, i = 3, \dots, 8$ is zero,
- (v) $\sum_{i=1}^8 U_i U_i^T = 2(2n + 1)I - 4J, \sum_{i=1}^8 V_i V_i^T = 2(2n + 1)I$.

Then there are 8-Williamson type matrices of order $2n + 1$. Furthermore, if the U_i and V_i are symmetric, $i = 1, \dots, 8$, then the 8-Williamson type matrices are symmetric. Hence there is a block type Hadamard matrix of order $8(2n + 1)$.

Corollary 25 *Let $q + 1$ be the order of amicable Hadamard matrices $I + S$ and P . Suppose there exist 4 Williamson type matrices of order q . Then there exist Williamson type matrices of order $2q + 1$. Furthermore there exists an Hadamard matrix of block type of order $8(2q + 1)$.*

Using the amicable Hadamard matrices given in [22] and [16, Table 1] we get 8 Williamson type matrices for the following orders for which 4 Williamson matrices are not known:

47, 111, 127, 167, 319, 487, 655, 831, ...

This gives new constructions for Hadamard matrices of orders 8.167 and 8.487.

Corollary 26 *Let q be a prime power and $(q-1)/2$ be the order of four (symmetric) Williamson type matrices. Then there exist (symmetric) 8-Williamson type matrices of order q and an Hadamard matrix of block structure of order $8q$.*

In particular we have 8-Williamson matrices for the following orders for which no Williamson type matrices are known:

59, 67, 103, 107, 151, 163, 179, 227, 251, 283, 347, 463, 467, 523, 563, 571, 587, 631, 643, 823, 859, 919, 947, ...

This gives new Hadamard matrices or new constructions for Hadamard matrices of orders 8.107, 8.163, 8.179, 8.251, 8.283, 8.347, 8.463, 8.523, 8.571, 8.631, 8.643, 8.823, 8.859, 8.919, 8.947, ...

Corollary 27 *Let $q \equiv 1 \pmod{4}$ be a prime power or $q+1$ the order of a symmetric conference matrix. Suppose there exist four (symmetric) Williamson type matrices of order q . Then there exist (symmetric) 8-Williamson type matrices of order $2q+1$ and an Hadamard matrix of block structure of order $8(2q+1)$.*

This corollary gives 8 Williamson type matrices for the following new orders: 219, 275, 299, 395, 483, 515, 579, 635, 699, 707, 723, 779, 795, 803, 899, 915, 923, ...

It does not give new Hadamard matrices for these orders.

Corollary 28 *Let $q = 9^t$, $t > 0$. Now there exist four (symmetric) Williamson type matrices of order 9^t , $t > 0$. Hence there exist (symmetric) 8-Williamson type matrices of order $2 \cdot 9^t + 1$, $t > 0$, and an Hadamard matrix of block structure of order $8(2 \cdot 9^t + 1)$.*

This gives symmetric 8-Williamson type matrices for the new order 163, 13123, ...

Also we have the following theorem:

Theorem 29 *Let U_{ij} , V_{ij} , $i, j = 1, \dots, 8$ be $(0, +1, -1)$ matrices of order n which satisfy*

- (i) U_{ki} , U_{kj} , $i \neq j$ are pairwise amicable, $k = 1, \dots, 8$,
- (ii) V_{ki} , V_{kj} , $i \neq j$ are pairwise amicable, $k = 1, \dots, 8$,

- (iii) $U_{ki} \pm V_{ki}$, $(+1, -1)$ matrices, $i, k = 1, \dots, 8$,
(iv) the row(column) sum of U_{ab} is 1 for $(a, b) \in \{(i, i), (i, i+1), (i+1, i)\}$,
 $i = 1, 3, 5, 7$, the row(column) sum of U_{aa} is -1 for $(a, a) = 2, 4, 6, 8$ and
otherwise, and the row(column) sum of U_{ij} , $i \neq j$ is zero,
(v) $\sum_{i=1}^8 U_{ji} U_{ji}^T = 2(2n+1)I - 4J$, $\sum_{i=1}^8 V_{ji} V_{ji}^T = 2(2n+1)I$, $j = 1, \dots, 8$,
(vi) $\sum_{i=1}^8 U_{ji} U_{ki}^T = 0$, $\sum_{i=1}^8 V_{ji} V_{ki}^T = 0$, $j \neq k$, $j, k = 1, \dots, 8$.

If (iii) to (vi) hold, there is a 64 block M -structure Hadamard matrix of order $8(2n+1)$.

Proof: Use

$$\begin{aligned}
X_{11} &= \begin{bmatrix} -1 & -e \\ -e^T & S_{11} \end{bmatrix}, & X_{12} &= \begin{bmatrix} -1 & -e \\ -e^T & S_{12} \end{bmatrix}, & X_{13} &= \begin{bmatrix} 1 & e \\ e^T & S_{13} \end{bmatrix}, & X_{14} &= \begin{bmatrix} 1 & e \\ e^T & S_{14} \end{bmatrix}, \\
X_{15} &= \begin{bmatrix} 1 & e \\ e^T & S_{15} \end{bmatrix}, & X_{16} &= \begin{bmatrix} 1 & e \\ e^T & S_{16} \end{bmatrix}, & X_{17} &= \begin{bmatrix} -1 & e \\ e^T & S_{17} \end{bmatrix}, & X_{18} &= \begin{bmatrix} -1 & e \\ e^T & S_{18} \end{bmatrix}, \\
X_{21} &= \begin{bmatrix} -1 & -e \\ -e^T & S_{21} \end{bmatrix}, & X_{22} &= \begin{bmatrix} 1 & e \\ e^T & S_{22} \end{bmatrix}, & X_{23} &= \begin{bmatrix} 1 & e \\ e^T & S_{23} \end{bmatrix}, & X_{24} &= \begin{bmatrix} -1 & -e \\ -e^T & S_{24} \end{bmatrix}, \\
X_{25} &= \begin{bmatrix} 1 & e \\ e^T & S_{25} \end{bmatrix}, & X_{26} &= \begin{bmatrix} -1 & -e \\ -e^T & S_{26} \end{bmatrix}, & X_{27} &= \begin{bmatrix} -1 & e \\ e^T & S_{27} \end{bmatrix}, & X_{28} &= \begin{bmatrix} 1 & -e \\ -e^T & S_{28} \end{bmatrix}, \\
X_{31} &= \begin{bmatrix} 1 & e \\ e^T & S_{31} \end{bmatrix}, & X_{32} &= \begin{bmatrix} 1 & e \\ e^T & S_{32} \end{bmatrix}, & X_{33} &= \begin{bmatrix} -1 & -e \\ -e^T & S_{33} \end{bmatrix}, & X_{34} &= \begin{bmatrix} -1 & -e \\ -e^T & S_{34} \end{bmatrix}, \\
X_{35} &= \begin{bmatrix} 1 & e \\ e^T & S_{35} \end{bmatrix}, & X_{36} &= \begin{bmatrix} 1 & e \\ e^T & S_{36} \end{bmatrix}, & X_{37} &= \begin{bmatrix} -1 & e \\ e^T & S_{37} \end{bmatrix}, & X_{38} &= \begin{bmatrix} -1 & e \\ e^T & S_{38} \end{bmatrix}, \\
X_{41} &= \begin{bmatrix} 1 & e \\ e^T & S_{41} \end{bmatrix}, & X_{42} &= \begin{bmatrix} -1 & -e \\ -e^T & S_{42} \end{bmatrix}, & X_{43} &= \begin{bmatrix} -1 & -e \\ -e^T & S_{43} \end{bmatrix}, & X_{44} &= \begin{bmatrix} 1 & e \\ e^T & S_{44} \end{bmatrix}, \\
X_{45} &= \begin{bmatrix} 1 & e \\ e^T & S_{45} \end{bmatrix}, & X_{46} &= \begin{bmatrix} -1 & -e \\ -e^T & S_{46} \end{bmatrix}, & X_{47} &= \begin{bmatrix} -1 & e \\ e^T & S_{47} \end{bmatrix}, & X_{48} &= \begin{bmatrix} 1 & -e \\ -e^T & S_{48} \end{bmatrix}, \\
X_{51} &= \begin{bmatrix} 1 & e \\ e^T & S_{51} \end{bmatrix}, & X_{52} &= \begin{bmatrix} 1 & e \\ e^T & S_{52} \end{bmatrix}, & X_{53} &= \begin{bmatrix} 1 & e \\ e^T & S_{53} \end{bmatrix}, & X_{54} &= \begin{bmatrix} 1 & e \\ e^T & S_{54} \end{bmatrix}, \\
X_{55} &= \begin{bmatrix} -1 & -e \\ -e^T & S_{55} \end{bmatrix}, & X_{56} &= \begin{bmatrix} -1 & -e \\ -e^T & S_{56} \end{bmatrix}, & X_{57} &= \begin{bmatrix} -1 & e \\ e^T & S_{57} \end{bmatrix}, & X_{58} &= \begin{bmatrix} -1 & e \\ e^T & S_{58} \end{bmatrix}, \\
X_{61} &= \begin{bmatrix} 1 & e \\ e^T & S_{61} \end{bmatrix}, & X_{62} &= \begin{bmatrix} -1 & -e \\ -e^T & S_{62} \end{bmatrix}, & X_{63} &= \begin{bmatrix} 1 & e \\ e^T & S_{63} \end{bmatrix}, & X_{64} &= \begin{bmatrix} -1 & -e \\ -e^T & S_{64} \end{bmatrix}, \\
X_{65} &= \begin{bmatrix} -1 & -e \\ -e^T & S_{65} \end{bmatrix}, & X_{66} &= \begin{bmatrix} 1 & e \\ e^T & S_{66} \end{bmatrix}, & X_{67} &= \begin{bmatrix} -1 & e \\ e^T & S_{67} \end{bmatrix}, & X_{68} &= \begin{bmatrix} 1 & -e \\ -e^T & S_{68} \end{bmatrix},
\end{aligned}$$

$$\begin{aligned}
X_{71} &= \begin{bmatrix} 1 & -e \\ -e^T & S_{71} \end{bmatrix}, & X_{72} &= \begin{bmatrix} 1 & -e \\ -e^T & S_{72} \end{bmatrix}, & X_{73} &= \begin{bmatrix} 1 & -e \\ -e^T & S_{73} \end{bmatrix}, & X_{74} &= \begin{bmatrix} -1 & -e \\ -e^T & S_{74} \end{bmatrix}, \\
X_{75} &= \begin{bmatrix} 1 & -e \\ -e^T & S_{75} \end{bmatrix}, & X_{76} &= \begin{bmatrix} 1 & -e \\ -e^T & S_{76} \end{bmatrix}, & X_{77} &= \begin{bmatrix} 1 & e \\ e^T & S_{77} \end{bmatrix}, & X_{78} &= \begin{bmatrix} 1 & e \\ e^T & S_{78} \end{bmatrix}, \\
X_{81} &= \begin{bmatrix} 1 & -e \\ -e^T & S_{81} \end{bmatrix}, & X_{82} &= \begin{bmatrix} -1 & e \\ e^T & S_{82} \end{bmatrix}, & X_{83} &= \begin{bmatrix} 1 & -e \\ -e^T & S_{83} \end{bmatrix}, & X_{84} &= \begin{bmatrix} -1 & e \\ e^T & S_{84} \end{bmatrix}, \\
X_{85} &= \begin{bmatrix} 1 & -e \\ -e^T & S_{85} \end{bmatrix}, & X_{86} &= \begin{bmatrix} -1 & e \\ e^T & S_{86} \end{bmatrix}, & X_{87} &= \begin{bmatrix} 1 & e \\ e^T & S_{87} \end{bmatrix}, & X_{88} &= \begin{bmatrix} -1 & -e \\ -e^T & S_{88} \end{bmatrix},
\end{aligned}$$

Then provided conditions (i) to (v) hold and $S_{7i}^T = S_{7i}$, $i = 1, \dots, 8$ are symmetric, X_{7i} , $i = 1, \dots, 8$ are symmetric 8-Williamson type matrices. Otherwise X_{7i} , $i = 1, \dots, 8$ are 8-Williamson type matrices. This can be verified by straightforward checking. Hence there is an Hadamard matrix of block structure of order $8(2n + 1)$.

If conditions (iii) to (vi) hold then straightforward verification shows the 64 block M-structure X_{ij} is an Hadamard matrix of order $8(2n + 1)$. \square

Corollary 30 *Let q be an odd prime power and suppose there exist Williamson-type matrices of order $\frac{1}{2}(q - 1)$. Then there exists an M-structure Hadamard matrix of order $8q$.*

Remark 31 This corollary gives new Hadamard matrices of order $8q$ for $q = 179, 1087, 1283, 1327, 1619, 1907, 2099, 2459, 2579, 2647, \dots$

Corollary 32 *Let $q = 2m + 1 \equiv 9 \pmod{16}$ be a prime power. Suppose there are Williamson-type matrices of order q . Then there is a M-structure Hadamard matrix of order $8(2q + 1)$.*

The analogous Yamada-J. Wallis-Whiteman structure to Theorem 29 is:

| | | | | | | | | | | | | | | | | | |
|-----------------|-----------------|-----------------|-----------------|-----------------|-----------------|----------------|-----------------|-----------------|-----------------|-----------------|-----------------|-----------------|-----------------|-----------------|-----------------|-----------------|-----------------|
| -1 | -1 | 1 | 1 | 1 | 1 | -1 | -1 | -e | -e | e | e | e | e | e | e | e | e |
| -1 | 1 | 1 | 1 | -1 | 1 | -1 | 1 | -e | e | e | -e | e | e | -e | e | e | -e |
| 1 | 1 | -1 | -1 | 1 | 1 | -1 | -1 | e | e | -e | -e | e | e | -e | e | e | e |
| 1 | -1 | -1 | 1 | 1 | -1 | -1 | 1 | e | -e | -e | e | e | e | -e | -e | e | -e |
| 1 | 1 | 1 | 1 | 1 | -1 | -1 | -1 | e | e | e | e | e | e | -e | -e | e | e |
| 1 | -1 | 1 | -1 | -1 | 1 | -1 | 1 | e | -e | e | -e | -e | -e | e | e | e | -e |
| 1 | 1 | 1 | 1 | 1 | 1 | 1 | 1 | -e | -e | -e | -e | -e | -e | -e | e | e | e |
| 1 | -1 | 1 | -1 | 1 | -1 | 1 | -1 | -e | e | -e | e | -e | e | -e | e | -e | e |
| -e ^T | -e ^T | e ^T | e ^T | e ^T | e ^T | e ^T | e ^T | S ₁₁ | S ₁₂ | S ₁₃ | S ₁₄ | S ₁₅ | S ₁₆ | S ₁₇ | S ₁₈ | S ₁₉ | S ₂₀ |
| -e ^T | e ^T | e ^T | -e ^T | e ^T | -e ^T | e ^T | -e ^T | S ₂₁ | S ₂₂ | S ₂₃ | S ₂₄ | S ₂₅ | S ₂₆ | S ₂₇ | S ₂₈ | S ₂₉ | S ₃₀ |
| e ^T | e ^T | -e ^T | -e ^T | e ^T | e ^T | e ^T | e ^T | S ₃₁ | S ₃₂ | S ₃₃ | S ₃₄ | S ₃₅ | S ₃₆ | S ₃₇ | S ₃₈ | S ₃₉ | S ₄₀ |
| e ^T | -e ^T | -e ^T | e ^T | e ^T | -e ^T | e ^T | -e ^T | S ₄₁ | S ₄₂ | S ₄₃ | S ₄₄ | S ₄₅ | S ₄₆ | S ₄₇ | S ₄₈ | S ₄₉ | S ₅₀ |
| e ^T | e ^T | e ^T | e ^T | -e ^T | -e ^T | e ^T | e ^T | S ₅₁ | S ₅₂ | S ₅₃ | S ₅₄ | S ₅₅ | S ₅₆ | S ₅₇ | S ₅₈ | S ₅₉ | S ₆₀ |
| e ^T | -e ^T | e ^T | -e ^T | -e ^T | e ^T | e ^T | -e ^T | S ₆₁ | S ₆₂ | S ₆₃ | S ₆₄ | S ₆₅ | S ₆₆ | S ₆₇ | S ₆₈ | S ₆₉ | S ₇₀ |
| -e ^T | -e ^T | e ^T | -e ^T | -e ^T | e ^T | e ^T | -e ^T | S ₇₁ | S ₇₂ | S ₇₃ | S ₇₄ | S ₇₅ | S ₇₆ | S ₇₇ | S ₇₈ | S ₇₉ | S ₈₀ |
| -e ^T | e ^T | -e ^T | e ^T | -e ^T | e ^T | e ^T | -e ^T | S ₈₁ | S ₈₂ | S ₈₃ | S ₈₄ | S ₈₅ | S ₈₆ | S ₈₇ | S ₈₈ | S ₈₉ | S ₉₀ |

We can see Yamada's matrix with trimming [46] or the J. Wallis-Whiteman [30] matrix with a border embodied in the construction.

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Index of Williamson Matrices

This table contains odd integers $q < 40000$ for which Williamson matrices exist. The following legend gives the method of construction used

| Key | Method | Explanation |
|-----|--------------------------------|--|
| w1 | $\{1, \dots, 33, 37, 41, 43\}$ | |
| w2 | $\frac{p+1}{2}$ | $p \equiv 1 \pmod{4}$ a prime power |
| w3 | 9^d | d a natural number |
| w4 | $\frac{p(p+1)}{2}$ | $p \equiv 1 \pmod{4}$ a prime power |
| w5 | $s(4s+3), s(4s-1)$ | $s \in \{1, 3, 5, \dots, 31\}$ |
| w6 | 93 | |
| w7 | $\frac{(f-1)(4f+1)}{4}$ | $p = 4f + 1, f$ odd, is a prime power of the form $1 + 4t^2$, $\frac{f-1}{8}$ is the order of a good matrix |
| w8 | $\frac{(f+1)(4f+1)}{4}$ | $p = 4f + 1, f$ odd, is a prime power of the form $25 + 4t^2$, $\frac{f+1}{8}$ is the order of a good matrix |
| w9 | $\frac{p(p-1)}{2}$ | $p = 4f + 1$ is a prime power and $\frac{p-1}{4}$ is the order of a good matrix |
| w0 | $(p+2)(p+1)$ | $p \equiv 1 \pmod{4}$ a prime power, $p+3$ is the order of a symmetric Hadamard matrix |
| wa | $\frac{(f+1)(4f+1)}{2}$ | $p = 4f + 1, f$ odd, is a prime power of the form $9 + 4t^2$, $\frac{f-1}{2} \equiv 1 \pmod{4}$ a prime power |
| wb | $\frac{(f-1)(4f+1)}{2}$ | $p = 4f + 1, f$ odd, is a prime power of the form $49 + 4t^2$, $\frac{f-3}{2} \equiv 1 \pmod{4}$ a prime power |
| wc | $2p + 1$ | $q = 2p - 1$ is a prime power p is a prime |
| wd | $7 \cdot 3^i$ | $i \geq 0$ |
| w#e | $7^{i+1}, 11 \cdot 7^i$ | $i \geq 0$ (Gives 8-Williamson matrices) |
| wf | $\frac{q^d(q+1)}{2}$ | $q \equiv 1 \pmod{4}$ is a prime $d \geq 2$ |
| wg | $\frac{p^2(p+1)}{2}$ | $p \equiv 1 \pmod{4}$ is a prime power |
| wh | $\frac{p^2(p+1)}{4}$ | $p \equiv 3 \pmod{4}$ is a prime power and $\frac{p+1}{4}$ is the order of a Williamson type matrix |
| wi | $q + 2$ | $q \equiv 1 \pmod{4}$, is a prime power and $\frac{q+1}{2}$ is a prime power |
| wj | $q + 2$ | $q \equiv 1 \pmod{4}$, is a prime power and $\frac{q+3}{2}$ is the order of a symmetric conference matrix |

| Key | Method | Explanation |
|-----|-------------------|---|
| wk | q | $q \equiv 1(mod4)$ is a prime power and $\frac{q-1}{2}$ is the order of a symmetric conference matrix or the order of a symmetric hadamard matrix |
| wl | q | $q \equiv 1(mod4)$, is a prime power and $\frac{q-1}{4}$ is the order of a williamson type matrix |
| wm | q | $q \equiv 1(mod4)$, is a prime power and $\frac{q-1}{2}$ is the order of a hadamard matrix |
| wn | wn | w is the order of a williamson type matrix n is the order of a symmetric conference matrix |
| wo | $2wu$ | w and u are the orders of williamson type matrices |
| w#p | $2q + 1$ | $q + 1$ is the order of an amicable hadamard matrix and q is the order of a williamson type matrix |
| w#q | q | q is a prime power and $\frac{q-1}{2}$ is the order of a williamson type matrix |
| w#r | $2q + 1$ | $q \equiv 1(mod4)$ is a prime power or $q + 1$ is the order of a symmetric conference matrix and q is the order of a williamson type matrix |
| w#s | $2 \cdot 9^t + 1$ | $t > 0$ |

$S = \{1, \dots, 31\}$ is the set of good matrices.

$q - 1$ is a Hadamard matrix of SC-form if one of the following is true

- (i) $\frac{q-1}{4}$ is a Williamson matrix.
- (ii) $\frac{q-1}{2}$ is a Conference matrix.
- (iii) $\frac{q-1}{4}$ is a Hadamard matrix. Note: The fact that if there is a

Williamson matrix of order n then there is a Williamson matrix of order $2n$, is used in the calculation of wg.

[The references for these papers are w1 [48], [4], [49], [40], [22], w2 [51], [33], w3 [52], w4 [19], w5, w6 [21], w7, w8, w9, w0, we, wl [20], wf, wg [17], wh, wi, wj, wk [12]]

| q | t | Method | q | t | Method | q | t | Method | q | t | Method | q | t | Method |
|----|---|--------|-----|---|--------|-----|---|--------|-----|---|--------|-----|---|--------|
| 1 | 0 | w1 | 101 | 0 | wk | 201 | 0 | w2 | 301 | 0 | w2 | 401 | 0 | wk |
| 3 | 0 | w1 | 103 | 0 | w#q | 203 | 1 | w9 | 303 | 1 | w7 | 403 | 1 | wn |
| 5 | 0 | w1 | 105 | 1 | wn | 205 | 0 | w2 | 305 | 1 | wn | 405 | 0 | w2 |
| 7 | 0 | w1 | 107 | 0 | w#q | 207 | 1 | wn | 307 | 0 | w2 | 407 | 1 | wn |
| 9 | 0 | w1 | 109 | 0 | wk | 209 | 1 | wn | 309 | 0 | w2 | 409 | 0 | wk |
| 11 | 0 | w1 | 111 | 1 | wn | 211 | 0 | w2 | 311 | | | 411 | 0 | w2 |
| 13 | 0 | w1 | 113 | 0 | wk | 213 | | | 313 | 0 | w2 | 413 | | |
| 15 | 0 | w1 | 115 | 0 | w2 | 215 | 1 | wn | 315 | 0 | w5 | 415 | 0 | w2 |
| 17 | 0 | w1 | 117 | 0 | w2 | 217 | 0 | w2 | 317 | 0 | wk | 417 | 1 | wn |
| 19 | 0 | w1 | 119 | 1 | wn | 219 | 1 | wn | 319 | 1 | wo | 419 | | |
| 21 | 0 | w1 | 121 | 0 | w2 | 221 | 1 | wn | 321 | 0 | w2 | 421 | 0 | w2 |
| 23 | 0 | w1 | 123 | 0 | wi | 223 | | | 323 | 1 | wn | 423 | 0 | wi |
| 25 | 0 | w1 | 125 | 0 | wk | 225 | 0 | w2 | 325 | 0 | w4 | 425 | 1 | wn |
| 27 | 0 | w1 | 127 | 0 | w#p | 227 | 0 | w#q | 327 | 0 | w2 | 427 | 0 | w2 |
| 29 | 0 | w1 | 129 | 0 | w2 | 229 | 0 | w2 | 329 | | | 429 | 0 | w2 |
| 31 | 0 | w1 | 131 | | | 231 | 0 | w2 | 331 | 0 | w2 | 431 | | |
| 33 | 0 | w1 | 133 | 1 | wn | 233 | 0 | w1 | 333 | 1 | w9 | 433 | 0 | wk |
| 35 | 1 | wn | 135 | 0 | w2 | 235 | | | 335 | | | 435 | 0 | w4 |
| 37 | 0 | w1 | 137 | 0 | w1 | 237 | 1 | wn | 337 | 0 | w2 | 437 | 1 | wn |
| 39 | 0 | wi | 139 | 0 | w2 | 239 | | | 339 | 0 | w2 | 439 | 0 | w2 |
| 41 | 0 | w1 | 141 | 0 | w2 | 241 | 0 | wk | 341 | 1 | wn | 441 | 0 | w2 |
| 43 | 0 | w1 | 143 | 1 | wn | 243 | 0 | wj | 343 | 1 | wn | 443 | | |
| 45 | 0 | w2 | 145 | 0 | w2 | 245 | 1 | wn | 345 | 1 | wn | 445 | 1 | wn |
| 47 | 0 | w#p | 147 | 0 | w2 | 247 | 1 | wn | 347 | 0 | w#q | 447 | 1 | wn |
| 49 | 0 | w2 | 149 | 0 | wk | 249 | 1 | wn | 349 | 0 | wk | 449 | 0 | wk |
| 51 | 0 | w2 | 151 | 0 | w#q | 251 | 0 | w#q | 351 | 0 | w2 | 451 | 0 | wj |
| 53 | 0 | wk | 153 | 0 | w4 | 253 | 1 | wn | 353 | 0 | w1 | 453 | | |
| 55 | 0 | w2 | 155 | 1 | wn | 255 | 0 | w2 | 355 | 0 | w2 | 455 | 1 | wn |
| 57 | 0 | w2 | 157 | 0 | w2 | 257 | 0 | w1 | 357 | 1 | wn | 457 | 0 | wk |
| 59 | 0 | w#q | 159 | 0 | w2 | 259 | 1 | wn | 359 | | | 459 | 0 | wi |
| 61 | 0 | w1 | 161 | 1 | wn | 261 | 0 | w2 | 361 | 0 | wk | 461 | 0 | wk |
| 63 | 0 | w2 | 163 | 0 | w#q | 263 | | | 363 | 0 | wi | 463 | 0 | w#q |
| 65 | 1 | wn | 165 | 1 | wn | 265 | 0 | w2 | 365 | 0 | w2 | 465 | 0 | w2 |
| 67 | 0 | w#q | 167 | 0 | w#p | 267 | 1 | wn | 367 | 0 | w2 | 467 | 0 | w#q |
| 69 | 0 | w2 | 169 | 0 | w2 | 269 | | | 369 | 1 | wn | 469 | 0 | w2 |
| 71 | | | 171 | 1 | wn | 271 | 0 | w2 | 371 | 1 | wn | 471 | 0 | w2 |
| 73 | 0 | wk | 173 | 0 | w1 | 273 | 1 | wn | 373 | 0 | w1 | 473 | 0 | w5 |
| 75 | 0 | w2 | 175 | 0 | w2 | 275 | 1 | wn | 375 | 0 | wf | 475 | 1 | wn |
| 77 | 1 | wn | 177 | 0 | w2 | 277 | 0 | wk | 377 | 1 | wn | 477 | 0 | w2 |
| 79 | 0 | w2 | 179 | 0 | w#q | 279 | 0 | w2 | 379 | 0 | w2 | 479 | | |
| 81 | 0 | w3 | 181 | 0 | w2 | 281 | 0 | w1 | 381 | 0 | w2 | 481 | 0 | w2 |
| 83 | 0 | wi | 183 | 1 | wn | 283 | 0 | w#q | 383 | | | 483 | 1 | wn |
| 85 | 0 | w2 | 185 | 1 | wn | 285 | 0 | w2 | 385 | 0 | w2 | 485 | 1 | wn |
| 87 | 0 | w2 | 187 | 0 | w2 | 287 | 1 | wn | 387 | 0 | w2 | 487 | 0 | w#p |
| 89 | 0 | w1 | 189 | 0 | w5 | 289 | 0 | w2 | 389 | 0 | wk | 489 | 0 | w2 |
| 91 | 0 | w2 | 191 | 0 | w#p | 291 | 1 | wn | 391 | 1 | wn | 491 | | |
| 93 | 0 | w6 | 193 | 0 | wk | 293 | 0 | w1 | 393 | | | 493 | 1 | wo |
| 95 | 0 | w5 | 195 | 0 | w2 | 295 | | | 395 | 1 | wn | 495 | 1 | wn |
| 97 | 0 | w2 | 197 | 0 | wk | 297 | 0 | w2 | 397 | 0 | wk | 497 | | |
| 99 | 0 | w2 | 199 | 0 | w2 | 299 | 1 | wn | 399 | 0 | w2 | 499 | 0 | w2 |

Existence of Williamson Matrices

| q | t | Method | q | t | Method | q | t | Method | q | t | Method | q | t | Method |
|-----|---|--------|-----|---|--------|-----|---|--------|-----|---|--------|-----|---|--------|
| 501 | | | 601 | 0 | w2 | 701 | 0 | wk | 801 | 0 | w2 | 901 | 0 | w2 |
| 503 | | | 603 | 1 | wn | 703 | 0 | w4 | 803 | 1 | wo | 903 | 1 | wn |
| 505 | 0 | w2 | 605 | 1 | wn | 705 | 0 | w2 | 805 | 0 | w2 | 905 | 1 | wn |
| 507 | 0 | w2 | 607 | 0 | w2 | 707 | 1 | wn | 807 | 0 | w2 | 907 | | |
| 509 | | | 609 | 0 | w2 | 709 | 0 | wk | 809 | 0 | wl | 909 | 1 | wn |
| 511 | 0 | w2 | 611 | | | 711 | 1 | wn | 811 | 0 | w2 | 911 | | |
| 513 | 1 | wn | 613 | 0 | wl | 713 | 1 | wn | 813 | 1 | wn | 913 | 1 | wo |
| 515 | 0 | w#r | 615 | 0 | w2 | 715 | 0 | w2 | 815 | | | 915 | 1 | w9 |
| 517 | 0 | w2 | 617 | 0 | wl | 717 | 0 | w2 | 817 | 1 | wn | 917 | | |
| 519 | 1 | wn | 619 | 0 | w2 | 719 | | | 819 | 0 | w2 | 919 | 0 | w#q |
| 521 | 0 | wl | 621 | 1 | wn | 721 | | | 821 | 0 | wk | 921 | 1 | wn |
| 523 | 0 | w#q | 623 | 1 | wn | 723 | 1 | wn | 823 | 0 | w#q | 923 | 0 | w#r |
| 525 | 0 | w2 | 625 | 0 | w2 | 725 | 1 | wn | 825 | 1 | wn | 925 | 0 | w2 |
| 527 | 1 | wn | 627 | 0 | wi | 727 | 0 | w2 | 827 | | | 927 | 1 | wn |
| 529 | 0 | wl | 629 | 1 | wn | 729 | 0 | w3 | 829 | 0 | w2 | 929 | 0 | wl |
| 531 | 0 | w2 | 631 | 0 | w#q | 731 | 1 | wo | 831 | 1 | wn | 931 | 0 | w2 |
| 533 | 1 | wn | 633 | 1 | wn | 733 | 0 | w#q | 833 | 1 | wn | 933 | | |
| 535 | 0 | w2 | 635 | 0 | w#r | 735 | 0 | wi | 835 | 0 | w2 | 935 | 1 | wn |
| 537 | | | 637 | 1 | wn | 737 | | | 837 | 1 | wn | 937 | 0 | w2 |
| 539 | 1 | wn | 639 | 0 | w2 | 739 | | | 839 | | | 939 | 0 | w2 |
| 541 | 0 | wk | 641 | 0 | wk | 741 | 0 | w2 | 841 | 0 | w2 | 941 | | |
| 543 | 0 | wi | 643 | 0 | w#q | 743 | | | 843 | 0 | wi | 943 | 1 | wn |
| 545 | 1 | wn | 645 | 0 | w2 | 745 | 0 | w2 | 845 | 1 | wn | 945 | 0 | w2 |
| 547 | 0 | w2 | 647 | | | 747 | 0 | w2 | 847 | 0 | w2 | 947 | 0 | w#q |
| 549 | 0 | w2 | 649 | 0 | w2 | 749 | | | 849 | 0 | w2 | 949 | 1 | wn |
| 551 | 1 | wn | 651 | 0 | w2 | 751 | 0 | w#q | 851 | 1 | wn | 951 | 0 | w2 |
| 553 | 1 | wn | 653 | | | 753 | | | 853 | | | 953 | 0 | wl |
| 555 | 0 | w2 | 655 | 0 | w#p | 755 | | | 855 | 0 | w2 | 955 | | |
| 557 | 0 | wk | 657 | 1 | wn | 757 | 0 | wl | 857 | | | 957 | 0 | w2 |
| 559 | 0 | w2 | 659 | | | 759 | 0 | wi | 859 | 0 | w#q | 959 | 1 | wn |
| 561 | 1 | wn | 661 | 0 | w2 | 761 | 0 | wl | 861 | 0 | w2 | 961 | 0 | wk |
| 563 | 0 | w#q | 663 | 0 | w5 | 763 | 1 | wa | 863 | | | 963 | 1 | wn |
| 565 | 0 | w2 | 665 | 1 | wn | 765 | 1 | wn | 865 | 1 | wn | 965 | 1 | wn |
| 567 | 1 | wn | 667 | 1 | wn | 767 | | | 867 | 0 | w2 | 967 | 0 | w2 |
| 569 | 0 | wm | 669 | | | 769 | 0 | wk | 869 | 1 | wn | 969 | 1 | wn |
| 571 | 0 | w#q | 671 | 1 | wn | 771 | 1 | wn | 871 | 0 | w2 | 971 | | |
| 573 | | | 673 | 0 | wk | 773 | 0 | wl | 873 | 1 | wn | 973 | 1 | wn |
| 575 | 1 | wn | 675 | 0 | wi | 775 | 0 | w2 | 875 | 1 | wn | 975 | 0 | w2 |
| 577 | 0 | w2 | 677 | 0 | wk | 777 | 0 | w2 | 877 | 0 | w2 | 977 | 0 | wk |
| 579 | 0 | wj | 679 | 1 | wn | 779 | 1 | wn | 879 | 0 | wi | 979 | 1 | wo |
| 581 | 1 | wn | 681 | 0 | w2 | 781 | | | 881 | 0 | wk | 981 | 1 | wn |
| 583 | 1 | wo | 683 | | | 783 | 1 | wn | 883 | 0 | w#q | 983 | | |
| 585 | 1 | wn | 685 | 0 | w2 | 785 | 1 | wn | 885 | 0 | w5 | 985 | 1 | wn |
| 587 | 0 | w#q | 687 | 0 | w2 | 787 | | | 887 | | | 987 | 0 | w2 |
| 589 | 1 | wn | 689 | 1 | w9 | 789 | | | 889 | 0 | w2 | 989 | 1 | wn |
| 591 | 0 | w2 | 691 | 0 | w2 | 791 | 1 | wn | 891 | 1 | wn | 991 | | |
| 593 | 0 | wk | 693 | 1 | wn | 793 | 1 | wn | 893 | | | 993 | 1 | wn |
| 595 | 1 | wn | 695 | 1 | wn | 795 | 1 | wn | 895 | 0 | w2 | 995 | 1 | wn |
| 597 | 0 | w2 | 697 | 1 | wn | 797 | 0 | wk | 897 | 1 | wn | 997 | 0 | w2 |
| 599 | | | 699 | 1 | wn | 799 | 0 | w2 | 899 | 1 | wn | 999 | 0 | w2 |

Existence of Williamson Matrices

| q | t | Method | q | t | Method | q | t | Method | q | t | Method | q | t | Method |
|------|---|--------|------|---|--------|------|---|--------|------|---|--------|------|---|--------|
| 1001 | 1 | wn | 1101 | 1 | wn | 1201 | 0 | w2 | 1301 | 0 | wl | 1401 | 0 | w2 |
| 1003 | | | 1103 | | | 1203 | 0 | wi | 1303 | 0 | w#q | 1403 | 1 | wn |
| 1005 | 1 | wn | 1105 | 0 | w2 | 1205 | 1 | wn | 1305 | 0 | w2 | 1405 | 0 | w2 |
| 1007 | 1 | wn | 1107 | 0 | w2 | 1207 | 0 | w5 | 1307 | 0 | w#r | 1407 | 1 | wn |
| 1009 | 0 | w2 | 1109 | 0 | wl | 1209 | 0 | w2 | 1309 | 0 | w2 | 1409 | 0 | wl |
| 1011 | 1 | wn | 1111 | 0 | w2 | 1211 | 1 | wn | 1311 | 0 | w2 | 1411 | 1 | wo |
| 1013 | 1 | wn | 1113 | 1 | wn | 1213 | 0 | w#q | 1313 | 1 | wn | 1413 | 1 | wn |
| 1015 | 0 | w2 | 1115 | 0 | w#r | 1215 | 0 | wi | 1315 | | | 1415 | | |
| 1017 | 1 | wn | 1117 | 0 | wk | 1217 | 0 | wk | 1317 | 0 | w2 | 1417 | 0 | w2 |
| 1019 | 0 | w#r | 1119 | 0 | w2 | 1219 | 0 | w2 | 1319 | | | 1419 | 0 | w2 |
| 1021 | 0 | wk | 1121 | | | 1221 | 0 | w2 | 1321 | 0 | wl | 1421 | 1 | wn |
| 1023 | 1 | wn | 1123 | | | 1223 | | | 1323 | 0 | wi | 1423 | | |
| 1025 | 1 | wn | 1125 | 1 | wn | 1225 | 0 | w4 | 1325 | 1 | wn | 1425 | 0 | w5 |
| 1027 | 0 | w2 | 1127 | 1 | wn | 1227 | 1 | wn | 1327 | 0 | w#p | 1427 | | |
| 1029 | 1 | wn | 1129 | 0 | wk | 1229 | 0 | wk | 1329 | 0 | w2 | 1429 | 0 | w2 |
| 1031 | | | 1131 | 1 | wn | 1231 | 0 | w#p | 1331 | 1 | wn | 1431 | 0 | w2 |
| 1033 | 0 | wk | 1133 | | | 1233 | 1 | wn | 1333 | 1 | wn | 1433 | | |
| 1035 | 0 | w2 | 1135 | 0 | w2 | 1235 | 1 | wn | 1335 | 1 | wn | 1435 | 1 | wn |
| 1037 | 1 | wn | 1137 | 0 | w2 | 1237 | 0 | w2 | 1337 | | | 1437 | | |
| 1039 | | | 1139 | 0 | w5 | 1239 | 0 | w2 | 1339 | 0 | w2 | 1439 | | |
| 1041 | 0 | w2 | 1141 | 0 | w2 | 1241 | 1 | wo | 1341 | 1 | wb | 1441 | | |
| 1043 | 1 | wn | 1143 | 1 | wn | 1243 | 1 | wn | 1343 | 1 | wn | 1443 | 1 | wn |
| 1045 | 0 | w2 | 1145 | 1 | wn | 1245 | 1 | wn | 1345 | 0 | w2 | 1445 | 1 | wn |
| 1047 | 1 | wn | 1147 | 0 | w2 | 1247 | 1 | wo | 1347 | 0 | w2 | 1447 | | |
| 1049 | 0 | wm | 1149 | 0 | w2 | 1249 | 0 | wl | 1349 | | | 1449 | 0 | w2 |
| 1051 | 0 | w#q | 1151 | | | 1251 | 0 | wi | 1351 | 1 | wn | 1451 | | |
| 1053 | 1 | wn | 1153 | 0 | wk | 1253 | | | 1353 | 1 | wn | 1453 | 0 | wl |
| 1055 | 1 | wn | 1155 | 0 | w2 | 1255 | | | 1355 | 1 | wn | 1455 | 0 | w2 |
| 1057 | 0 | w2 | 1157 | 1 | wn | 1257 | | | 1357 | 0 | w2 | 1457 | | |
| 1059 | 1 | wn | 1159 | 1 | wn | 1259 | | | 1359 | | | 1459 | 0 | w2 |
| 1061 | 0 | wk | 1161 | 1 | wn | 1261 | 0 | w2 | 1361 | 0 | wk | 1461 | | |
| 1063 | 0 | w#q | 1163 | | | 1263 | 1 | wn | 1363 | | | 1463 | 2 | wn |
| 1065 | 0 | w2 | 1165 | 1 | wn | 1265 | 1 | wn | 1365 | 0 | w2 | 1465 | 1 | wn |
| 1067 | 1 | wn | 1167 | 0 | w2 | 1267 | 1 | wn | 1367 | | | 1467 | 1 | wn |
| 1069 | 0 | w2 | 1169 | | | 1269 | 1 | wn | 1369 | 0 | wl | 1469 | 1 | wn |
| 1071 | 0 | w2 | 1171 | 0 | w2 | 1271 | 1 | wn | 1371 | 0 | w2 | 1471 | 0 | w#p |
| 1073 | 1 | wn | 1173 | 1 | wn | 1273 | | | 1373 | 0 | w#q | 1473 | | |
| 1075 | 1 | wn | 1175 | | | 1275 | 0 | w2 | 1375 | 0 | w2 | 1475 | | |
| 1077 | 0 | w2 | 1177 | | | 1277 | 0 | w#q | 1377 | 0 | w2 | 1477 | 0 | w2 |
| 1079 | 1 | wn | 1179 | 0 | w2 | 1279 | 0 | w2 | 1379 | 1 | wn | 1479 | 0 | w2 |
| 1081 | 0 | w2 | 1181 | | | 1281 | 1 | wn | 1381 | 0 | w#q | 1481 | 0 | wl |
| 1083 | 1 | wn | 1183 | 0 | wf | 1283 | 0 | w#q | 1383 | 0 | wi | 1483 | 0 | w#q |
| 1085 | 1 | wn | 1185 | 1 | wn | 1285 | 1 | wn | 1385 | 1 | wn | 1485 | 0 | w2 |
| 1087 | 0 | w#p | 1187 | 0 | w#q | 1287 | 1 | wn | 1387 | 1 | wn | 1487 | | |
| 1089 | 1 | wn | 1189 | 0 | w2 | 1289 | 0 | wl | 1389 | 0 | w2 | 1489 | 0 | wl |
| 1091 | | | 1191 | 0 | w2 | 1291 | 0 | w#q | 1391 | | | 1491 | | |
| 1093 | 0 | w#q | 1193 | 0 | wl | 1293 | | | 1393 | 1 | wn | 1493 | 0 | wl |
| 1095 | 0 | wi | 1195 | 0 | w2 | 1295 | 2 | wn | 1395 | 0 | w2 | 1495 | 1 | wn |
| 1097 | 0 | wl | 1197 | 0 | w2 | 1297 | 0 | w2 | 1397 | | | 1497 | 1 | wn |
| 1099 | 0 | w2 | 1199 | 1 | wo | 1299 | 1 | wn | 1399 | 0 | w2 | 1499 | | |

Existence of Williamson Matrices

| q | t | Method | q | t | Method | q | t | Method | q | t | Method | q | t | Method |
|------|---|--------|------|---|--------|------|---|--------|------|---|--------|------|---|--------|
| 1501 | 0 | w2 | 1601 | 0 | wk | 1701 | 1 | wn | 1801 | 0 | wk | 1901 | 0 | w#q |
| 1503 | | | 1603 | 1 | wn | 1703 | | | 1803 | 1 | wn | 1903 | 1 | wo |
| 1505 | 1 | wn | 1605 | 0 | w2 | 1705 | 1 | wn | 1805 | 0 | wh | 1905 | 1 | wn |
| 1507 | 1 | wo | 1607 | | | 1707 | 0 | w2 | 1807 | 0 | w2 | 1907 | 0 | w#q |
| 1509 | | | 1609 | 0 | w2 | 1709 | 0 | wk | 1809 | 0 | w2 | 1909 | 1 | wn |
| 1511 | | | 1611 | 0 | w2 | 1711 | | | 1811 | | | 1911 | 0 | w2 |
| 1513 | 1 | wn | 1613 | 0 | w#q | 1713 | | | 1813 | 1 | wn | 1913 | | |
| 1515 | 1 | wn | 1615 | 0 | w2 | 1715 | 0 | w#r | 1815 | 1 | wn | 1915 | | |
| 1517 | 1 | wn | 1617 | 1 | wn | 1717 | 0 | w2 | 1817 | 1 | wn | 1917 | 0 | w2 |
| 1519 | 0 | w2 | 1619 | 0 | w#q | 1719 | | | 1819 | 0 | w2 | 1919 | 1 | wn |
| 1521 | 0 | w2 | 1621 | 0 | wk | 1721 | 0 | wl | 1821 | 1 | wn | 1921 | 1 | wn |
| 1523 | 0 | w#q | 1623 | 0 | wi | 1723 | 0 | w#q | 1823 | 1 | wn | 1923 | 1 | wn |
| 1525 | 0 | w2 | 1625 | 1 | wn | 1725 | 0 | w2 | 1825 | 1 | wn | 1925 | 1 | wn |
| 1527 | | | 1627 | 0 | w2 | 1727 | 1 | wn | 1827 | 0 | w5 | 1927 | 0 | w2 |
| 1529 | 1 | wn | 1629 | 0 | w2 | 1729 | 0 | w2 | 1829 | | | 1929 | | |
| 1531 | 0 | w2 | 1631 | 1 | wn | 1731 | 0 | w2 | 1831 | | | 1931 | | |
| 1533 | 1 | wn | 1633 | | | 1733 | 0 | wl | 1833 | 1 | wn | 1933 | 0 | w#q |
| 1535 | 1 | wn | 1635 | 1 | wn | 1735 | 0 | w2 | 1835 | 1 | wn | 1935 | 0 | wi |
| 1537 | 1 | wo | 1637 | 0 | wl | 1737 | 1 | wn | 1837 | 0 | w2 | 1937 | 1 | wn |
| 1539 | 1 | wn | 1639 | 1 | wo | 1739 | | | 1839 | 0 | w2 | 1939 | 0 | w2 |
| 1541 | | | 1641 | 1 | wn | 1741 | 0 | w2 | 1841 | | | 1941 | 0 | w2 |
| 1543 | | | 1643 | 1 | wn | 1743 | 0 | w5 | 1843 | 1 | wn | 1943 | | |
| 1545 | 0 | w2 | 1645 | | | 1745 | 1 | wn | 1845 | 1 | wn | 1945 | 0 | w2 |
| 1547 | 1 | wn | 1647 | 1 | wn | 1747 | | | 1847 | | | 1947 | 1 | wn |
| 1549 | 0 | wk | 1649 | 1 | wn | 1749 | 1 | wn | 1849 | 0 | w2 | 1949 | | |
| 1551 | 1 | wn | 1651 | 0 | w2 | 1751 | | | 1851 | 0 | w2 | 1951 | 0 | w#p |
| 1553 | 0 | wk | 1653 | 1 | wn | 1753 | 0 | wl | 1853 | 1 | wo | 1953 | 1 | wn |
| 1555 | 0 | w2 | 1655 | 1 | wn | 1755 | 0 | wi | 1855 | 0 | w2 | 1955 | 1 | wn |
| 1557 | 1 | wn | 1657 | 0 | w2 | 1757 | | | 1857 | 1 | wn | 1957 | | |
| 1559 | | | 1659 | 0 | wi | 1759 | 0 | w2 | 1859 | 1 | wn | 1959 | 0 | w2 |
| 1561 | 0 | w2 | 1661 | | | 1761 | | | 1861 | 0 | w2 | 1961 | 1 | wn |
| 1563 | 0 | w2 | 1663 | | | 1763 | 1 | wn | 1863 | 1 | wn | 1963 | | |
| 1565 | 1 | wn | 1665 | 0 | w2 | 1765 | 0 | w2 | 1865 | 1 | wn | 1965 | 0 | w2 |
| 1567 | | | 1667 | | | 1767 | 0 | w2 | 1867 | 0 | w2 | 1967 | 1 | wn |
| 1569 | 0 | w2 | 1669 | 0 | w#q | 1769 | 1 | wn | 1869 | 1 | wn | 1969 | | |
| 1571 | | | 1671 | 1 | wn | 1771 | 0 | w2 | 1871 | | | 1971 | 1 | wn |
| 1573 | 1 | wn | 1673 | | | 1773 | 1 | wn | 1873 | 0 | wk | 1973 | 0 | w#q |
| 1575 | 1 | wn | 1675 | | | 1775 | 1 | wn | 1875 | 0 | wf | 1975 | 1 | wn |
| 1577 | 1 | wn | 1677 | 1 | wn | 1777 | 0 | wl | 1877 | 0 | wk | 1977 | | |
| 1579 | | | 1679 | 1 | wn | 1779 | 0 | w2 | 1879 | 0 | w#q | 1979 | | |
| 1581 | 1 | wn | 1681 | 0 | w2 | 1781 | 1 | wn | 1881 | 0 | w2 | 1981 | | |
| 1583 | | | 1683 | 0 | wj | 1783 | | | 1883 | 0 | w#r | 1983 | 1 | wn |
| 1585 | 0 | w2 | 1685 | 1 | wn | 1785 | 1 | wn | 1885 | 0 | w2 | 1985 | 1 | wn |
| 1587 | 1 | wh | 1687 | 0 | w2 | 1787 | | | 1887 | 1 | wn | 1987 | | |
| 1589 | | | 1689 | | | 1789 | 0 | w#q | 1889 | 0 | wm | 1989 | 1 | wn |
| 1591 | 0 | w2 | 1691 | 1 | wn | 1791 | 0 | w2 | 1891 | 0 | w4 | 1991 | 1 | wn |
| 1593 | 1 | wn | 1693 | 0 | wl | 1793 | | | 1893 | | | 1993 | 0 | wl |
| 1595 | 1 | wn | 1695 | 0 | w2 | 1795 | | | 1895 | 1 | wn | 1995 | 0 | w2 |
| 1597 | 0 | wk | 1697 | 0 | wl | 1797 | 0 | w2 | 1897 | 0 | w2 | 1997 | 0 | wk |
| 1599 | 1 | wn | 1699 | 0 | w#q | 1799 | 1 | wn | 1899 | 0 | w2 | 1999 | 0 | w#p |

Existence of Williamson Matrices

Index of Hadamard Matrices

This table contains odd integers $q < 40000$ for which Hadamard matrices of the form $2^t q$ exist. The key for the methods of construction follows.

Amicable Hadamard Matrices.

| Key | Method | Explanation |
|-----|------------|---|
| a1 | $p^r + 1$ | $p^r \equiv 3(\text{mod}4)$, is a prime power |
| a2 | $2(q + 1)$ | $2q + 1$ is a prime power, $q \equiv 1(\text{mod}4)$, is a prime |
| a5 | nh | n, h , are amicable hadamard matrices |

Skew Hadamard Matrices.

| Key | Method | Explanation |
|-----|--------------------|---|
| s1 | $2^t \prod k_i$ | t all positive integers, $k_i - 1 \equiv 3(\text{mod}4)$ a prime power |
| s2 | $(p - 1)^u + 1$ | p is a skew Hadamard matrix, $u > 0$ is an odd integer |
| s3 | $2(q + 1)$ | $q \equiv 5(\text{mod}8)$ is a prime power |
| s4 | $2(q + 1)$ | $q = p^t$ is a prime power where $p \equiv 5(\text{mod}8)$ and $t \equiv 2(\text{mod}4)$. |
| s5 | $4m$ | $3 \leq m \leq 25$ |
| s6 | $4(q + 1)$ | $q \equiv 9(\text{mod}16)$ is a prime power |
| s7 | $(t + 1)(q + 1)$ | $q = s^2 + 4t^2 \equiv 5(\text{mod}8)$ is a prime power and $ t + 1$ is a skew Hadamard matrix |
| s8 | $4(q^2 + q + 1)$ | q is a prime power, $q^2 + q + 1 \equiv 3, 5, 7(\text{mod}8)$ a prime or $2(q^2 + q + 1) + 1$ is a prime power |
| s0 | hm | h is a skew hadamard matrix and m is an amicable hadamard matrix |

Spence Hadamard Matrices.

| Key | Method | Explanation |
|-----|------------------|---|
| p1 | $4(q^2 + q + 1)$ | $q^2 + q + 1 \equiv 1(\text{mod}8)$ is a prime |
| p2 | $4n$ or $8n$ | $n, n - 2$ are prime powers, if $n \equiv 1(\text{mod}4)$ there exists a Hadamard matrix of order $4n$, if $n \equiv 3(\text{mod}4)$ there exists a Hadamard matrix of order $8n$ |
| p3 | $4m$ | m is an odd prime power for which an integer $s \geq 0$ such that $\frac{(m - (2^{s+1} + 1))}{2^{s+1}}$ is an odd prime power, exists |

Symmetric Hadamard Matrices.

If there exists a Conference matrix of order n then there is symmetric Hadamard matrix of order $2n$, for this reason symmetric hadamard matrices indexed according to the method used to derive the order of a conference matrix with the exception of c6 which produces a symmetric Hadamard matrix.

| Key | Method | Explanation |
|-----|------------------------|--|
| c1 | $p^r + 1$ | $p^r \equiv 1(\text{mod}4)$ is a prime power |
| c2 | $(h - 1)^2 + 1$ | h is a skew Hadamard matrix |
| c3 | $q^2(q - 2) + 1$ | $q \equiv 3(\text{mod}4)$ is a prime power $q - 2$ is a prime power |
| c4 | $5 \cdot 9^{2t+1} + 1$ | $t \geq 0$ |
| c5 | $(n - 1)^s + 1$ | n is a conference matrix $s \geq 2$ |
| c6 | nh | n is a conference matrix h is a Hadamard matrix |

Note: a conference matrix of order n exists only if $n - 1$ is the sum of two squares.

Hadamard Matrices Obtained From Williamson Matrices.

If a Williamson matrix of order $2^t q$ exists then there is a Hadamard matrix of order $2^{t+2} q$, the same key as in the Index of Williamson Matrices is used to index the Hadamard matrices produced from them.

OD Hadamard Matrices.

| Key | Method | Explanation |
|-----|--------|--|
| o1 | | If a T-matrix of order $2^t q$ exists then there is a hadamard matrix of order $2^{t+2} q$ |
| o2 | ow | o is an OD-hadamard matrix and w is a Williamson matrix |

Yamada Hadamard Matrices.

| Key | Method | Explanation |
|-----|------------|---|
| y1 | $4q$ | $q \equiv 1(\text{mod}8)$ is a prime power $\frac{q-1}{2}$ is a Hadamard matrix |
| y2 | $4(q + 2)$ | $q \equiv 5(\text{mod}8)$ is a prime power $\frac{q+3}{2}$ is a skew Hadamard matrix |
| y3 | $4(q + 2)$ | $q \equiv 1(\text{mod}8)$ is a prime power $\frac{q+3}{2}$ is a conference matrix |

Miyamoto Hadamard Matrices.

| Key | Method | Explanation |
|-----|--------|--|
| m1 | $4q$ | $q \equiv 1(\text{mod}4)$ is a prime power $q - 1$ is a Hadamard matrix |
| m2 | $8q$ | $q \equiv 3(\text{mod}4)$ is a prime power $2q - 3$ is a prime power |

Seberry.

| Key | Method | Explanation |
|-----|---------|---|
| se | $2^t q$ | where t is the smallest integer such that for given odd q , $a(q + 1) + b(q - 3) = 2^t$ has a solution for a, b non-negative integers |

| q | t | Method | q | t | Method | q | t | Method | q | t | Method | q | t | Method |
|----|---|--------|-----|---|--------|-----|---|--------|-----|---|--------|-----|----|--------|
| 1 | | a2 | 101 | | wk | 201 | | o2 | 301 | | o2 | 401 | | wk |
| 3 | | a2 | 103 | | y2 | 203 | | a2 | 303 | | o2 | 403 | | o2 |
| 5 | | a2 | 105 | | a2 | 205 | | o2 | 305 | | o2 | 405 | | a2 |
| 7 | | a2 | 107 | 3 | w#q | 207 | | a2 | 307 | | s3 | 407 | | a2 |
| 9 | | o2 | 109 | | wk | 209 | | o2 | 309 | | c1 | 409 | | wk |
| 11 | | a2 | 111 | | a2 | 211 | | s3 | 311 | 3 | m2 | 411 | | o2 |
| 13 | | s4 | 113 | | c2 | 213 | 3 | c6 | 313 | | c1 | 413 | | o2 |
| 15 | | a2 | 115 | | o2 | 215 | | a2 | 315 | | a2 | 415 | | o2 |
| 17 | | a2 | 117 | | a2 | 217 | | o2 | 317 | | wk | 417 | | a2 |
| 19 | | s3 | 119 | | o2 | 219 | | o2 | 319 | | o2 | 419 | 4 | a2 |
| 21 | | a2 | 121 | | o2 | 221 | | a2 | 321 | | a2 | 421 | | s4 |
| 23 | | s5 | 123 | | a2 | 223 | 3 | a2 | 323 | | a2 | 423 | | o2 |
| 25 | | o2 | 125 | | a2 | 225 | | o2 | 325 | | o2 | 425 | | a2 |
| 27 | | a2 | 127 | | y2 | 227 | | a2 | 327 | | a2 | 427 | | o2 |
| 29 | | w1 | 129 | | o2 | 229 | | c1 | 329 | | o2 | 429 | | o2 |
| 31 | | s3 | 131 | | a2 | 231 | | o2 | 331 | | s3 | 431 | | a2 |
| 33 | | a2 | 133 | | o2 | 233 | | w1 | 333 | | a2 | 433 | | wk |
| 35 | | a2 | 135 | | o2 | 235 | | o2 | 335 | | o2 | 435 | | o2 |
| 37 | | c1 | 137 | | a2 | 237 | | a2 | 337 | | c1 | 437 | | a2 |
| 39 | | o2 | 139 | | s3 | 239 | 4 | a2 | 339 | | o2 | 439 | | s3 |
| 41 | | a2 | 141 | | a2 | 241 | | wk | 341 | | o2 | 441 | | o2 |
| 43 | | w1 | 143 | | a2 | 243 | | a2 | 343 | | o2 | 443 | 3 | m2 |
| 45 | | a2 | 145 | | o2 | 245 | | o2 | 345 | | o2 | 445 | | o2 |
| 47 | | o1 | 147 | | a2 | 247 | | o2 | 347 | 3 | w#q | 447 | | a2 |
| 49 | | o2 | 149 | | wk | 249 | | o2 | 349 | | wk | 449 | | wk |
| 51 | | o2 | 151 | | y2 | 251 | 3 | w#q | 351 | | o2 | 451 | | o2 |
| 53 | | a2 | 153 | | o2 | 253 | | o2 | 353 | | w1 | 453 | | a2 |
| 55 | | o2 | 155 | | a2 | 255 | | a2 | 355 | | s3 | 455 | | o2 |
| 57 | | a2 | 157 | | c1 | 257 | | w1 | 357 | | a2 | 457 | | wk |
| 59 | | o1 | 159 | | o2 | 259 | | o2 | 359 | 4 | a2 | 459 | | o2 |
| 61 | | a2 | 161 | | a2 | 261 | | o2 | 361 | | o2 | 461 | | wk |
| 63 | | a2 | 163 | 3 | a2 | 263 | | a2 | 363 | | a2 | 463 | 3 | w#q |
| 65 | | o2 | 165 | | a2 | 265 | | o2 | 365 | | a2 | 465 | | o2 |
| 67 | | o1 | 167 | 3 | w#p | 267 | | o2 | 367 | | s3 | 467 | | a2 |
| 69 | | o2 | 169 | | o2 | 269 | | m1 | 369 | | o2 | 469 | | o2 |
| 71 | | a2 | 171 | | a2 | 271 | | s3 | 371 | | a2 | 471 | | o2 |
| 73 | | wk | 173 | | a2 | 273 | | a2 | 373 | | w1 | 473 | | o2 |
| 75 | | o2 | 175 | | o2 | 275 | | o2 | 375 | | a2 | 475 | | o2 |
| 77 | | a2 | 177 | | o2 | 277 | | wk | 377 | | o2 | 477 | | a2 |
| 79 | | s3 | 179 | 3 | w#q | 279 | | o2 | 379 | | s3 | 479 | 16 | se |
| 81 | | o2 | 181 | | c1 | 281 | | a2 | 381 | | a2 | 481 | | o2 |
| 83 | | a2 | 183 | | o2 | 283 | 3 | w#q | 383 | | a2 | 483 | | a2 |
| 85 | | o2 | 185 | | a2 | 285 | | o2 | 385 | | o2 | 485 | | o2 |
| 87 | | a2 | 187 | | o2 | 287 | | o2 | 387 | | o2 | 487 | 3 | w#p |
| 89 | | w1 | 189 | | o2 | 289 | | o2 | 389 | | wk | 489 | | c1 |
| 91 | | o2 | 191 | 3 | w#p | 291 | | a2 | 391 | | o2 | 491 | 15 | se |
| 93 | | o2 | 193 | | wk | 293 | | a2 | 393 | | a2 | 493 | | o2 |
| 95 | | a2 | 195 | | o2 | 295 | | o2 | 395 | | a2 | 495 | | a2 |
| 97 | | c1 | 197 | | a2 | 297 | | a2 | 397 | | wk | 497 | | a2 |
| 99 | | o2 | 199 | | s3 | 299 | | o2 | 399 | | o2 | 499 | | s3 |

Existence of Hadamard Matrices

| q | t | Method | q | t | Method | q | t | Method | q | t | Method | q | t | Method |
|-----|----|--------|-----|----|--------|-----|----|--------|-----|----|--------|-----|---|--------|
| 501 | | a2 | 601 | | c1 | 701 | | a2 | 801 | | a2 | 901 | | o2 |
| 503 | | a2 | 603 | | a2 | 703 | | o2 | 803 | | o2 | 903 | | o2 |
| 505 | | o2 | 605 | | o2 | 705 | | a2 | 805 | | o2 | 905 | | o2 |
| 507 | | a2 | 607 | | s3 | 707 | | o2 | 807 | | s3 | 907 | 3 | m2 |
| 509 | | m1 | 609 | | o2 | 709 | | wk | 809 | | wl | 909 | | o2 |
| 511 | | o2 | 611 | | o2 | 711 | | a2 | 811 | | s3 | 911 | | a2 |
| 513 | | o2 | 613 | | c2 | 713 | | a2 | 813 | | a2 | 913 | | o2 |
| 515 | 3 | w#r | 615 | | a2 | 715 | | o2 | 815 | | a2 | 915 | | a2 |
| 517 | | o2 | 617 | | a2 | 717 | | c1 | 817 | | o2 | 917 | 4 | a5 |
| 519 | | o2 | 619 | | s3 | 719 | 4 | a2 | 819 | | o2 | 919 | 3 | a2 |
| 521 | | a2 | 621 | | o2 | 721 | 4 | o2 | 821 | | wk | 921 | | o2 |
| 523 | 3 | w#q | 623 | | o2 | 723 | | o2 | 823 | 3 | s7 | 923 | | a2 |
| 525 | | a2 | 625 | | o2 | 725 | | o2 | 825 | | a2 | 925 | | o2 |
| 527 | | o2 | 627 | | o2 | 727 | | s3 | 827 | | a2 | 927 | | o2 |
| 529 | | o2 | 629 | | o2 | 729 | | o2 | 829 | | c1 | 929 | | wl |
| 531 | | o2 | 631 | 3 | w#q | 731 | | o2 | 831 | | a2 | 931 | | o2 |
| 533 | | a2 | 633 | | a2 | 733 | | m1 | 833 | | a2 | 933 | 4 | c6 |
| 535 | | s3 | 635 | | a2 | 735 | | a2 | 835 | | s3 | 935 | | a2 |
| 537 | 4 | c6 | 637 | | o2 | 737 | | o2 | 837 | | a2 | 937 | | c1 |
| 539 | | o2 | 639 | | s3 | 739 | 16 | se | 839 | 18 | se | 939 | | o2 |
| 541 | | wk | 641 | | wk | 741 | | a2 | 841 | | o2 | 941 | | m1 |
| 543 | | o2 | 643 | 3 | w#q | 743 | | a2 | 843 | | a2 | 943 | | o2 |
| 545 | | a2 | 645 | | a2 | 745 | | o2 | 845 | | o2 | 945 | | a2 |
| 547 | | a2 | 647 | 3 | m2 | 747 | | o2 | 847 | | o2 | 947 | 3 | w#q |
| 549 | | o2 | 649 | | o2 | 749 | 5 | o2 | 849 | | c1 | 949 | | o2 |
| 551 | | a2 | 651 | | o2 | 751 | 3 | a2 | 851 | | o2 | 951 | | a2 |
| 553 | | o2 | 653 | 4 | o1 | 753 | | a2 | 853 | 3 | a2 | 953 | | wl |
| 555 | | o2 | 655 | | y2 | 755 | | a2 | 855 | | o2 | 955 | 3 | a2 |
| 557 | | wk | 657 | | o2 | 757 | | s8 | 857 | | m1 | 957 | | o2 |
| 559 | | o2 | 659 | 17 | se | 759 | | o2 | 859 | 3 | a2 | 959 | | o2 |
| 561 | | a2 | 661 | | c1 | 761 | | c2 | 861 | | o2 | 961 | | o2 |
| 563 | | a2 | 663 | | o2 | 763 | | o2 | 863 | 3 | m2 | 963 | | a2 |
| 565 | | o2 | 665 | | a2 | 765 | | o2 | 865 | | o2 | 965 | | o2 |
| 567 | | a2 | 667 | | o2 | 767 | | a2 | 867 | | a2 | 967 | | s3 |
| 569 | | p3 | 669 | 3 | a2 | 769 | | wk | 869 | | o2 | 969 | | o2 |
| 571 | 3 | a2 | 671 | | a2 | 771 | | a2 | 871 | | o2 | 971 | 6 | a2 |
| 573 | 3 | a2 | 673 | | wk | 773 | | wl | 873 | | a2 | 973 | | o2 |
| 575 | | o2 | 675 | | a2 | 775 | | o2 | 875 | | a2 | 975 | | o2 |
| 577 | | c1 | 677 | | a2 | 777 | | o2 | 877 | | c1 | 977 | | a2 |
| 579 | | o2 | 679 | | o2 | 779 | | o2 | 879 | | o2 | 979 | | o2 |
| 581 | | o2 | 681 | | c1 | 781 | 3 | a2 | 881 | | wk | 981 | | a2 |
| 583 | | o2 | 683 | | a2 | 783 | | o2 | 883 | 3 | w#q | 983 | | a2 |
| 585 | | a2 | 685 | | o2 | 785 | | o2 | 885 | | a2 | 985 | | o2 |
| 587 | | a2 | 687 | | o2 | 787 | 3 | m2 | 887 | | a2 | 987 | | a2 |
| 589 | | o2 | 689 | | o2 | 789 | 3 | a2 | 889 | | c1 | 989 | | o2 |
| 591 | | o2 | 691 | | s3 | 791 | | a2 | 891 | | o2 | 991 | 3 | a2 |
| 593 | | a2 | 693 | | o2 | 793 | | o2 | 893 | | a2 | 993 | | o2 |
| 595 | | o2 | 695 | | o2 | 795 | | o2 | 895 | | s3 | 995 | | o2 |
| 597 | | o2 | 697 | | o2 | 797 | | a2 | 897 | | o2 | 997 | | c1 |
| 599 | 17 | se | 699 | | o2 | 799 | | o2 | 899 | | o2 | 999 | | o2 |

Existence of Hadamard Matrices

| q | t | Method | q | t | Method | q | t | Method | q | t | Method | q | t | Method |
|------|---|--------|------|---|--------|------|----|--------|------|----|--------|------|----|--------|
| 1001 | | a2 | 1101 | | o2 | 1201 | | c1 | 1301 | | c2 | 1401 | | c1 |
| 1003 | | o2 | 1103 | 3 | m2 | 1203 | | o2 | 1303 | 3 | w#q | 1403 | | o2 |
| 1005 | | a2 | 1105 | | o2 | 1205 | | o2 | 1305 | | o2 | 1405 | | o2 |
| 1007 | | a2 | 1107 | | o2 | 1207 | | w5 | 1307 | | a2 | 1407 | | o2 |
| 1009 | | c1 | 1109 | | wl | 1209 | | o2 | 1309 | | o2 | 1409 | | wl |
| 1011 | | o2 | 1111 | | o2 | 1211 | | o2 | 1311 | | o2 | 1411 | | o2 |
| 1013 | | a2 | 1113 | | a2 | 1213 | | m1 | 1313 | | o2 | 1413 | | a2 |
| 1015 | | o2 | 1115 | 3 | w#r | 1215 | | o2 | 1315 | 4 | a5 | 1415 | | a2 |
| 1017 | | o2 | 1117 | | wk | 1217 | | wk | 1317 | | o2 | 1417 | | o2 |
| 1019 | 3 | w#r | 1119 | | o2 | 1219 | | o2 | 1319 | 18 | se | 1419 | | o2 |
| 1021 | | wk | 1121 | | a2 | 1221 | | o2 | 1321 | | wl | 1421 | | a2 |
| 1023 | | a2 | 1123 | 3 | m2 | 1223 | 19 | se | 1323 | | o2 | 1423 | 3 | a2 |
| 1025 | | a2 | 1125 | | o2 | 1225 | | o2 | 1325 | | o2 | 1425 | | o2 |
| 1027 | | o2 | 1127 | | a2 | 1227 | | o2 | 1327 | 3 | w#p | 1427 | 3 | m2 |
| 1029 | | o2 | 1129 | | wk | 1229 | | wk | 1329 | | c1 | 1429 | | c1 |
| 1031 | 6 | a2 | 1131 | | a2 | 1231 | | y2 | 1331 | | a2 | 1431 | | o2 |
| 1033 | | wk | 1133 | 4 | a2 | 1233 | | a2 | 1333 | | o2 | 1433 | | m1 |
| 1035 | | a2 | 1135 | | s3 | 1235 | | o2 | 1335 | | o2 | 1435 | | o2 |
| 1037 | | o2 | 1137 | | a2 | 1237 | | c1 | 1337 | | a2 | 1437 | 4 | o1 |
| 1039 | 3 | a2 | 1139 | | o2 | 1239 | | o2 | 1339 | | s3 | 1439 | 19 | se |
| 1041 | | c1 | 1141 | | c1 | 1241 | | o2 | 1341 | | o2 | 1441 | 3 | a2 |
| 1043 | | o2 | 1143 | | o2 | 1243 | | o2 | 1343 | | o2 | 1443 | | o2 |
| 1045 | | o2 | 1145 | | o2 | 1245 | | o2 | 1345 | | c1 | 1445 | | a2 |
| 1047 | | o2 | 1147 | | o2 | 1247 | | a2 | 1347 | | a2 | 1447 | 19 | se |
| 1049 | | p3 | 1149 | | c1 | 1249 | | wl | 1349 | 4 | o2 | 1449 | | o2 |
| 1051 | 3 | w#q | 1151 | | a2 | 1251 | | a2 | 1351 | | o2 | 1451 | 6 | a2 |
| 1053 | | a2 | 1153 | | wk | 1253 | | a2 | 1353 | | o2 | 1453 | | wl |
| 1055 | | a2 | 1155 | | o2 | 1255 | 3 | a2 | 1355 | | a2 | 1455 | | o2 |
| 1057 | | c1 | 1157 | | o2 | 1257 | 5 | c6 | 1357 | | o2 | 1457 | | a2 |
| 1059 | | o2 | 1159 | | o2 | 1259 | 4 | a2 | 1359 | 3 | c6 | 1459 | | s3 |
| 1061 | | a2 | 1161 | | a2 | 1261 | | o2 | 1361 | | a2 | 1461 | | a2 |
| 1063 | 3 | w#q | 1163 | | a2 | 1263 | | a2 | 1363 | | o2 | 1463 | | a2 |
| 1065 | | a2 | 1165 | | o2 | 1265 | | a2 | 1365 | | o2 | 1465 | | o2 |
| 1067 | | o2 | 1167 | | o2 | 1267 | | o2 | 1367 | 3 | m2 | 1467 | | a2 |
| 1069 | | c1 | 1169 | 5 | o2 | 1269 | | o2 | 1369 | | o2 | 1469 | | o2 |
| 1071 | | a2 | 1171 | | s3 | 1271 | | o2 | 1371 | | a2 | 1471 | 3 | w#p |
| 1073 | | o2 | 1173 | | a2 | 1273 | | o2 | 1373 | | m1 | 1473 | 3 | a2 |
| 1075 | | o2 | 1175 | | o2 | 1275 | | a2 | 1375 | | o2 | 1475 | | o2 |
| 1077 | | c1 | 1177 | 5 | a2 | 1277 | | a2 | 1377 | | a2 | 1477 | | o2 |
| 1079 | | o2 | 1179 | | s3 | 1279 | | s3 | 1379 | | o2 | 1479 | | o2 |
| 1081 | | o2 | 1181 | | a2 | 1281 | | o2 | 1381 | | m1 | 1481 | | a2 |
| 1083 | | o2 | 1183 | | o2 | 1283 | 3 | w#q | 1383 | | a2 | 1483 | 3 | a2 |
| 1085 | | a2 | 1185 | | o2 | 1285 | | o2 | 1385 | | o2 | 1485 | | a2 |
| 1087 | 3 | w#p | 1187 | 3 | w#q | 1287 | | a2 | 1387 | | o2 | 1487 | 3 | m2 |
| 1089 | | o2 | 1189 | | o2 | 1289 | | wl | 1389 | | c1 | 1489 | | wl |
| 1091 | | a2 | 1191 | | o2 | 1291 | 3 | w#q | 1391 | | a2 | 1491 | 3 | a2 |
| 1093 | | p2 | 1193 | | wl | 1293 | | a2 | 1393 | | o2 | 1493 | | wl |
| 1095 | | o2 | 1195 | | s3 | 1295 | | a2 | 1395 | | o2 | 1495 | | o2 |
| 1097 | | wl | 1197 | | a2 | 1297 | | c1 | 1397 | 4 | o2 | 1497 | | a2 |
| 1099 | | o2 | 1199 | | o2 | 1299 | | o2 | 1399 | | s3 | 1499 | 18 | se |

Existence of Hadamard Matrices

| q | t | Method | q | t | Method | q | t | Method | q | t | Method | q | t | Method |
|------|----|--------|------|---|--------|------|----|--------|------|---|--------|------|----|--------|
| 1501 | | o2 | 1601 | | wk | 1701 | | a2 | 1801 | | wk | 1901 | | a2 |
| 1503 | | a2 | 1603 | | o2 | 1703 | 4 | o2 | 1803 | | a2 | 1903 | | o2 |
| 1505 | | o2 | 1605 | | o2 | 1705 | | o2 | 1805 | | a2 | 1905 | | o2 |
| 1507 | | o2 | 1607 | | a2 | 1707 | | a2 | 1807 | | o2 | 1907 | 3 | w#q |
| 1509 | 3 | a2 | 1609 | | c1 | 1709 | | wk | 1809 | | o2 | 1909 | | o2 |
| 1511 | | a2 | 1611 | | s3 | 1711 | | o2 | 1811 | | a2 | 1911 | | a2 |
| 1513 | | o2 | 1613 | | a2 | 1713 | 4 | a2 | 1813 | | o2 | 1913 | 4 | o1 |
| 1515 | | o2 | 1615 | | o2 | 1715 | | a2 | 1815 | | o2 | 1915 | 3 | a2 |
| 1517 | | a2 | 1617 | | o2 | 1717 | | o2 | 1817 | | o2 | 1917 | | o2 |
| 1519 | | o2 | 1619 | 3 | w#q | 1719 | 3 | a2 | 1819 | | s3 | 1919 | | o2 |
| 1521 | | o2 | 1621 | | wk | 1721 | | a2 | 1821 | | a2 | 1921 | | o2 |
| 1523 | | a2 | 1623 | | a2 | 1723 | | s8 | 1823 | | c4 | 1923 | | a2 |
| 1525 | | o2 | 1625 | | o2 | 1725 | | a2 | 1825 | | o2 | 1925 | | a2 |
| 1527 | 3 | c6 | 1627 | | s3 | 1727 | | a2 | 1827 | | a2 | 1927 | | o2 |
| 1529 | | o2 | 1629 | | o2 | 1729 | | o2 | 1829 | | o2 | 1929 | 4 | c6 |
| 1531 | | s3 | 1631 | | o2 | 1731 | | o2 | 1831 | 3 | m2 | 1931 | | a2 |
| 1533 | | a2 | 1633 | 3 | a2 | 1733 | | wl | 1833 | | a2 | 1933 | | p2 |
| 1535 | | o2 | 1635 | | o2 | 1735 | | s3 | 1835 | | o2 | 1935 | | o2 |
| 1537 | | o2 | 1637 | | a2 | 1737 | | a2 | 1837 | | c1 | 1937 | | o2 |
| 1539 | | o2 | 1639 | | o2 | 1739 | | o2 | 1839 | | o2 | 1939 | | o2 |
| 1541 | | a2 | 1641 | | a2 | 1741 | | c1 | 1841 | 4 | a5 | 1941 | | c1 |
| 1543 | 3 | a2 | 1643 | | a2 | 1743 | | a2 | 1843 | | o2 | 1943 | | o2 |
| 1545 | | o2 | 1645 | | o2 | 1745 | | o2 | 1845 | | o2 | 1945 | | o2 |
| 1547 | | o2 | 1647 | | o2 | 1747 | 3 | m2 | 1847 | 3 | m2 | 1947 | | o2 |
| 1549 | | wk | 1649 | | o2 | 1749 | | o2 | 1849 | | o2 | 1949 | 4 | a2 |
| 1551 | | a2 | 1651 | | s3 | 1751 | 4 | o2 | 1851 | | o2 | 1951 | | y2 |
| 1553 | | a2 | 1653 | | o2 | 1753 | | wl | 1853 | | a2 | 1953 | | o2 |
| 1555 | | s3 | 1655 | | a2 | 1755 | | a2 | 1855 | | o2 | 1955 | | o2 |
| 1557 | | o2 | 1657 | | c1 | 1757 | | a2 | 1857 | | o2 | 1957 | 4 | o2 |
| 1559 | 4 | a2 | 1659 | | o2 | 1759 | | s3 | 1859 | | o2 | 1959 | | s3 |
| 1561 | | c1 | 1661 | 4 | o2 | 1761 | | a2 | 1861 | | s4 | 1961 | | o2 |
| 1563 | | o2 | 1663 | 3 | m2 | 1763 | | o2 | 1863 | | a2 | 1963 | | s7 |
| 1565 | | o2 | 1665 | | a2 | 1765 | | o2 | 1865 | | a2 | 1965 | | c1 |
| 1567 | 19 | se | 1667 | 3 | m2 | 1767 | | o2 | 1867 | | s3 | 1967 | | a2 |
| 1569 | | c1 | 1669 | | p2 | 1769 | | o2 | 1869 | | o2 | 1969 | 10 | a5 |
| 1571 | 18 | se | 1671 | | o2 | 1771 | | o2 | 1871 | 3 | m2 | 1971 | | a2 |
| 1573 | | o2 | 1673 | | a2 | 1773 | | o2 | 1873 | | wk | 1973 | | m1 |
| 1575 | | a2 | 1675 | | o2 | 1775 | | o2 | 1875 | | a2 | 1975 | | o2 |
| 1577 | | o2 | 1677 | | o2 | 1777 | | wl | 1877 | | a2 | 1977 | | a2 |
| 1579 | 5 | a2 | 1679 | | o2 | 1779 | | o2 | 1879 | 3 | a2 | 1979 | 4 | a2 |
| 1581 | | a2 | 1681 | | o2 | 1781 | | o2 | 1881 | | a2 | 1981 | 5 | a2 |
| 1583 | 3 | m2 | 1683 | | o2 | 1783 | 18 | se | 1883 | 3 | w#r | 1983 | | o2 |
| 1585 | | o2 | 1685 | | o2 | 1785 | | o2 | 1885 | | o2 | 1985 | | o2 |
| 1587 | | o2 | 1687 | | o2 | 1787 | 3 | m2 | 1887 | | a2 | 1987 | 16 | se |
| 1589 | 4 | a2 | 1689 | 3 | c6 | 1789 | | p2 | 1889 | | m1 | 1989 | | o2 |
| 1591 | | o2 | 1691 | | a2 | 1791 | | o2 | 1891 | | o2 | 1991 | | a2 |
| 1593 | | o2 | 1693 | | wl | 1793 | 4 | a2 | 1893 | 4 | c6 | 1993 | | wl |
| 1595 | | a2 | 1695 | | a2 | 1795 | 6 | a5 | 1895 | | o2 | 1995 | | o2 |
| 1597 | | wk | 1697 | | wl | 1797 | | a2 | 1897 | | o2 | 1997 | | wk |
| 1599 | | o2 | 1699 | 3 | a2 | 1799 | | o2 | 1899 | | o2 | 1999 | | y2 |

Existence of Hadamard Matrices

Toward the classification of distance-transitive graphs of affine type

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1. はじめに

近年、重要視されてきている問題に distance regular graph の完全分類問題がある。その分類問題の一つの部分問題として distance transitive graph の分類がある。distance regular graph の部分クラスである distance transitive graph は、本来 distance regular graph が持つ組合せ論的な性質の上に群論的な性質を持っており、この二つの重要な性質により、distance transitive graph の分類は distance regular graph の分類と比較して容易であると考えられる。しかも、distance transitive graph の分類は distance regular graph の分類に大きく貢献することにもなる（一つのステップとも考えられる）。すなわち、まず解き易い問題として distance transitive graph の分類を捉えることができる。現在 distance transitive graph の分類は一步一步進展しており、ここでは、distance transitive graph の分類プログラムの概要とその途中経過について報告する。

2. distance transitive graph の分類プログラム

distance transitive graph の分類問題は大きく次の部分問題（ステップ）に分けることができる。

(1) primitive distance transitive graph の分類

(2) primitive の場合の分類を用いて imprimitive の場合を分類する。

imprimitive な場合には、derived graph もしくは halved graph (bipartite half) を考えることにより primitive なグラフが現れるので、(1) \Rightarrow (2) という問題の 2 ステップ分割が考えられる。(この状況は distance regular graph でも同様であるが、現在の distance regular graph の分類においてはこの戦略はとっていないようである。)

(2) の問題は distance regular graph の場合と比べて、自己同形群の立場を使えるので、(2) の解決はかなり容易であると思われる。というのは、imprimitive なグラフ Γ より作られる primitive なグラフ Γ' が derived graph である場合には、 $\text{Aut}(\Gamma) = K \cdot \text{Aut}(\Gamma')$ 、 K はある正規部分群、となり、halved graph である場合には $\text{Aut}(\Gamma) = \text{Aut}(\Gamma')^2$ となる。この事実は群論的にかかなり強い制約であることから、『問題が解ける』ことが予想される。

さて、上により問題はまず(1)を解くことに帰着された。そこで、以下 primitive distance transitive graph の分類を考える。diameter が2以下の場合には、2重可移群の分類、ランク3の置換群の分類により完全に解決せれている。そこで、diameter は3以上であるとしてよいことになる。primitive distance transitive graph の分類は次の定理により更に細分化される。

定理 (Praeger & Saxl & Yokoyama (1987))

Γ を diameter が2以上の primitive distance transitive graph とし、その自己同形群を G とおく。この時、次のいずれかが成り立つ。

- (i) Γ は Hamming graph であるか、もしくは diameter が2である Hamming graph の complement graph である。
 - (ii) G は almost simple である。(すなわち、ある単純群 X が存在して $X \leq G \leq \text{Aut}(X)$)
 - (iii) G は Γ の頂点全体に正則に作用する elementary abelian 正規部分群を持つ。
- ((ii) の場合を almost simple の場合と呼び、(iii) の場合を affine の場合と呼ぶ。)

上の定理により分類問題は二つの部分問題へと分割されることになる。一つは (ii) almost simple の場合の分類であり、他方は (iii) affine の場合の分類である。almost simple の場合の分類は Van Bon, Cohen, Cuypers, Inglis, Ivanov, Liebeck, Praeger, Saxl 等により、有限単純群の分類定理を利用した『蝨潰し』的な方法で研究されており、置換指標の既約分解における各既約指標の重複度が1である事実を利用して解いたものと、それにグラフ的性質を利用して解いたものがある。分類リストを挙げると、

交代群 (対称群) - Saxl, Ivanov, Liebeck & Praeger & Saxl

$\text{PSL}(n, q)$ - Inglis, Liebeck & Saxl, Van Bon & Cohen

13次元以下の古典群 - Inglis

Held の群 - Van Bon & Cohen & Cuypers

almost simple の場合はほぼ終結したと言えるが、現時点では、まだ完全分類の結果は論文として出版されていない。

ここまでの話は、Brouer & Cohen & Neumaier (1989) または、Bannai & Ito (1986) を参照されたい。

3. affine の場合の分類

『affine の場合の分類をどのように解くか』については方針が固まっているわけではない。現在 affine の場合に現れるグラフは diameter に関する無限系列として 3 種類のみが知られている。

Hermitian forms graph - 隣接する 2 点を含む極大関は唯一つ

Bilinear forms graph - 隣接する 2 点を含む極大関は 2 個

Alternating forms graph - 隣接する 2 点を含む極大関は 3 個以上

(diameter が 3 の無限系列には Affine $E_6(q)$ graph がある。Quadratic forms graph は distance regular ではあるが、distance transitive ではない。)

上の状況より、隣接する 2 点を含む極大関の個数を指定してグラフを分類するアプローチが考えられる。そして、現時点では diameter が充分大きい場合に、『隣接する 2 点を含む極大関の個数を指定すればグラフは上記のものが抽出される』ことが目標になる。

このアプローチ上のひとつの結果として、隣接する 2 点を含む極大関が 2 個の場合について次がある。

定理 (Yokoyama (1989)) Γ が affine の場合でしかも正則部分群は 2-群ではないとし、更に次の 2 条件を満たすとす。

- (1) 隣接する 2 点を含む極大関は 2 個、
- (2) Γ の自己同形群 G は三角形にはならない長さ 2 のパス全体の上に可移に作用する。

この時、 Γ は Bilinear forms graph である。

残念ながら、その後の進展はまだない。現在私が挑戦している問題は

(1) 隣接する 2 点を含む極大関が唯一の場合を分類する。(Hermitian forms graph を特徴付け、抽出する。)

- (2) 上の定理を条件 (2) を取って証明する。

の 2 点である。

最後に分類を試みる上での重要と思われる点について述べておく。上記の 3 種の無限系列において、特徴的であることは、無限系列における diameter が n のグラフを $\Gamma(n)$ とおけば、次が成り立つことである。

$\Gamma(n)$ の距離 i ($2 \leq i \leq n$) の 2 頂点を x, y とする。この時、 x, y を含む部分グラフ Γ' で、 $\Gamma(i)$ に同形であるものが唯一存在する。

この性質により、与えられたグラフが $\Gamma(n)$ であることを示すためには、局所的な部分グラフとして $\Gamma(i)$ に同形なものを (帰納的に) 構成していく方法が考えられるのである。(実際、隣接する 2 点を含む極大関が 2 個の場合に Bilinear forms graph を抽出するにはこの手法を用いた。)

今回の報告の終わりにあたり、分類プログラムの概要の説明に留まり、新しい結果の発表に至らなかったことをお詫びします。

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INTERSECTION DIAGRAMS OF DISTANCE-REGULAR GRAPHS

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1. INTRODUCTION

Let G be a connected graph and let ∂ denote the usual metric on the vertex set $V = V(G)$ of G . For vertices u, v in G and non-negative integers r, s , we define

$$\Gamma_r(u) = \{x \in V \mid \partial(x, u) = r\},$$
$$D_s^r(u, v) = \Gamma_r(u) \cap \Gamma_s(v).$$

G is said to be *distance-regular* if the size of D_s^r depends only on the distance between u and v , rather than the individual vertices. In this case, we write $p_{rs}^t = |D_s^r(u, v)|$, where $t = \partial(u, v)$. Let $d = d(G)$ be the diameter of G and k be the valency of G , and let

$$\left\{ \begin{array}{cccccc} 0 & c_1 & \cdots & c_r & \cdots & c_{d-1} & c_d \\ 0 & a_1 & \cdots & a_r & \cdots & a_{d-1} & a_d \\ k & b_1 & \cdots & b_r & \cdots & b_{d-1} & 0 \end{array} \right\}$$

be the *intersection array* of G , where $c_r = p_{1\ r-1}^r$, $a_r = p_{1\ r}^r$, $b_r = p_{1\ r+1}^r$. More precise description about distance-regular graphs will be found in [1], [2].

In this note we shall prove the following result.

THEOREM 1. *Let G be a distance-regular graph with odd diameter $d = 2r + 1 \geq 3$. If G has the following intersection array*

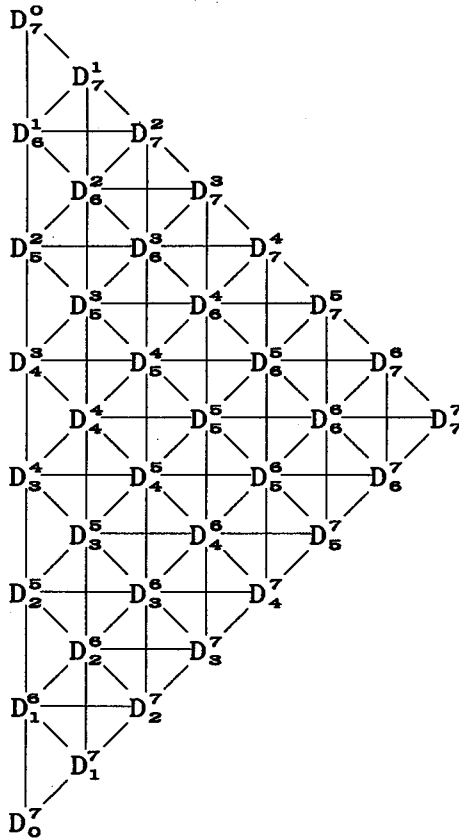
$$\left\{ \begin{array}{cccccccc} 0 & 1 & c_2 & \dots & c_r & c_{r+1} & c_{r+2} & \dots & c_{d-1} & k \\ 0 & 0 & 0 & \dots & 0 & a_{r+1} & 0 & \dots & 0 & 0 \\ k & b_1 & b_2 & \dots & b_r & b_{r+1} & b_{r+2} & \dots & b_{d-1} & 0 \end{array} \right\}$$

and $p_{j d}^d = 0$ for $j \geq 3$, then G is bipartite (i.e. $a_{r+1} = 0$).

Theorem 1 gives a partial answer to a question of E. Bannai and T. Ito about multiple P-polynomial structure of association schemes (see [1], III.4, pp.259). The notion of the intersection diagrams plays a very important role in the proof of the above theorem. We have already obtained several applications of the intersection diagram with respect to adjacent vertices. In this note we shall use the intersection diagram with respect to non-adjacent vertices.

2. THE INTERSECTION DIAGRAM

Let G be a distance-regular graph which satisfies the assumptions of Theorem 1. Fix two vertices u, v in G with $\partial(u, v) = d = d(G)$, and put $D_j^i = D_j^i(u, v)$. We draw the family $\{D_j^i\}_{ij}$ as follows. Throughout this paper we shall give figures for the case $d = 7 = 2r + 1$, $r = 3$.



In the above diagram, a line between two entries indicates possibility of existence of edges between them. We call the above diagram the *intersection diagram* of G with respect to u, v .

LEMMA 2. $D_j^d = D_d^j = \emptyset$ for $3 \leq j \leq d$.

Proof. Since $\partial(u, v) = d$, $D_s^i = p_{i s}^d$ holds for every i, s . By the assumption of Theorem 1 we have $p_{j d}^d = 0$, therefore $D_d^j = \emptyset$ for $j \geq 3$. By similar arguments we have also $D_j^d = \emptyset$ since $p_{d j}^d = p_{j d}^d = 0$ for $j \geq 3$. ■

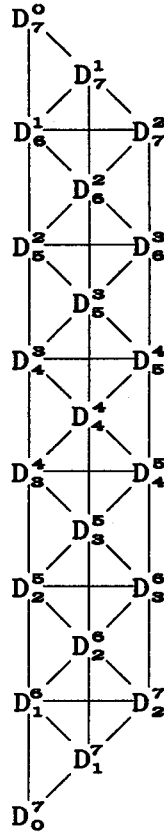
LEMMA 3. $D_j^i = \emptyset$ for all i, j with $i + j \geq d + 3$.

Proof. Assume $D_j^i \neq \emptyset$ for some i, j with $i + j \geq d + 3$, where we take i to be maximal. Then $i < d$ by Lemma 2. Let x be a vertex in D_j^i . Since $x \in \Gamma_i(u)$, there are b_i edges from x to $\Gamma_{i+1}(u)$, where $b_i > 0$ by $i < d$. This implies

$$D_{j-1}^{i+1} \cup D_j^{i+1} \cup D_{j+1}^{i+1} \neq \emptyset.$$

This contradicts to the maximality of i . ■

By Lemma 3, the intersection diagram becomes as follows.

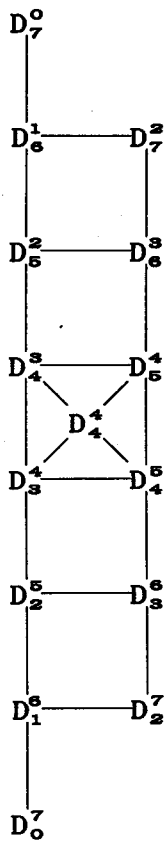


LEMMA 4. $D_{d+1-i}^i = \emptyset$ for $i \neq r+1$.

Proof. Assume $D_{d+1-i}^i \neq \emptyset$ for some i , where we take i to be minimal. It is clear $i > 0$, and we may assume $i \leq r$. Put $j = d+1-i$, and take a vertex x in D_j^i . We have $D_{j+1}^{i-1} = \emptyset$ by the minimality of i . Since $x \in \Gamma_i(u)$, there are c_i edges from x to $\Gamma_{i-1}(u)$. So we get $e(x, D_j^{i-1}) = c_i > 0$. Therefore there is an edge xy with $y \in D_j^{i-1}$. Thus xy is an edge in $\Gamma_j(v)$. But we

have $a_j = 0$ since $r+2 \leq j \leq d$. a contradiction. ■

By Lemma 4, the diagram becomes as the following.



3. EDGE PATTERNS

We consider the intersection diagram given in the previous section. For a vertex x in a entry D_j^i of the diagram, we shall determine the number of edges from x to another entry $D_j^{i'}$. For a subset A in G , the number of edges from a vertex x to A will be denoted by $e(x, A)$.

LEMMA 5. *If $i + j = d$ and $x \in D_j^i$. Then $e(x, D_{j+1}^{i-1}) = c_i$ and $e(x, D_{j-1}^{i+1}) = c_j$ hold.*

Proof. Since $x \in \Gamma_i(u)$, there are c_i edges from x to $\Gamma_{i-1}(u)$. Then we get $e(x, D_{j+1}^{i-1}) = c_i$. We get also $e(x, D_{j-1}^{i+1}) = e(x, \Gamma_{j-1}(v)) = c_j$.

LEMMA 6. *If $i + j = d + 2$ and $x \in D_j^i$. Then $e(x, D_{j+1}^{i-1}) = b_j$ and $e(x, D_{j-1}^{i+1}) = b_i$ hold.*

LEMMA 7. *If $i + j = d$, $1 \leq i \leq r$ and $x \in D_j^i$, then $e(x, D_{j+1}^{i+1}) = b_j - c_i$.*

Proof. Since $x \in \Gamma_j(v)$ and D_{j+1}^i is empty, we have

$$b_j = e(x, \Gamma_{j+1}(v)) = e(x, D_{j+1}^{i-1}) + e(x, D_{j+1}^{i+1}).$$

We have also, by Lemma 5, $e(x, D_{j+1}^{i-1}) = c_i$. Thus we get $e(x, D_{j+1}^{i+1}) = b_j - c_i$. ■

LEMMA 8. If $i + j = d$, $1 \leq i \leq r$ and $x \in D_{j+1}^{i+1}$, then $e(x, D_j^i) = c_{i+1} - b_{j+1}$.

LEMMA 9. If $x \in D_{r+1}^{r+1}$ then $e(x, D_{r+1}^r) = c_{r+1}$ and $e(x, D_r^{r+1}) = b_{r+1}$.

LEMMA 10. If $x \in D_{r+1}^r$ then $e(x, D_{r+1}^{r+1}) = a_{r+1}$.

Proof. Since $x \in \Gamma_{r+1}(v)$, there are a_{r+1} edges from x to $\Gamma_{r+1}(v)$. Here, there is no edge from x to D_{r+1}^r since $D_{r+1}^r \subset \Gamma_r(u)$ and $a_r = 0$. This implies $e(x, D_{r+1}^{r+1}) = a_{r+1}$. ■

LEMMA 11. If $x \in D_{r+2}^{r+1}$ then $e(x, D_{r+1}^{r+1}) = a_{r+1}$.

4. PROOF OF THEOREM 1

Let G be a distance-regular graph which satisfies the assumption of Theorem 1. We assume $a_{r+1} > 0$. Let m denote the size of D_{r+1}^{r+1} . Remark $m > 0$ by Lemma 10. By counting the number of edges between D_{r+1}^{r+1} and D_{r+1}^r using Lemma 9 and Lemma 10, we get $|D_{r+1}^r| = \frac{mc_{r+1}}{a_{r+1}}$. Similarly we get $|D_{r+2}^{r+1}| = \frac{mb_{r+1}}{a_{r+1}}$ by Lemma 9, 11. Then by counting the number of edges between D_{r+1}^r and D_{r+2}^{r+1} using Lemma 7, 8, we get,

$$\frac{mc_{r+1}}{a_{r+1}}(b_{r+1} - c_r) = \frac{mb_{r+1}}{a_{r+1}}(c_{r+1} - b_{r+2})$$

This implies the following relation by $m > 0$.

$$c_r c_{r+1} = b_{r+1} b_{r+2}. \quad (1)$$

By counting the number of edges between D_{r+2}^{r-1} and D_{r+3}^r , we get

$$\frac{mc_{r+1}c_r}{a_{r+1}c_{r+2}}(b_{r+2} - c_{r-1}) = \frac{mb_{r+1}b_{r+2}}{a_{r+1}b_r}(c_r - b_{r+3}).$$

Here we have $b_{r+2} - c_{r-1} = b_{r-1} - c_{r+2}$, $c_r - b_{r+3} = c_{r+3} - b_r$. So the above equality implies, by using (1),

$$b_{r-1}b_r = c_{r+2}c_{r+3} \quad (2)$$

By counting the number of edges between D_{r+3}^{r-2} and D_{r+4}^{r-1} , we get

$$\frac{mc_{r+1}c_r c_{r-1}}{a_{r+1}c_{r+2}c_{r+3}}(b_{r+3} - c_{r-2}) = \frac{mb_{r+1}b_{r+2}b_{r+3}}{a_{r+1}b_r b_{r-1}}(c_{r-1} - b_{r+4}).$$

This implies, by (1) and (2),

$$c_{r-2}c_{r-1} = b_{r+3}b_{r+4} \quad (3)$$

We repeat the above arguments. In the case r is odd, we get

$$c_1 c_2 = b_{2r} b_{2r+1}.$$

This is impossible since $b_{2r+1} = b_d = 0$. In the case r is even, we get

$$b_1 b_2 = c_{2r} c_{2r+1}.$$

Here we have $b_1 = k - 1$ and $c_{2r+1} = k$. Thus we get

$$(k - 1) b_2 = c_{2r} k.$$

Then k divides b_2 and hence $k \leq b_2$, a contradiction. ■

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Factor sets associated with regular collineation groups

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(Yutaka Hiramine)

§1. Introduction

正則な自己同型群 A をもつ対称テザイン \mathbb{D} は A が含む差集合を用いて記述できることが知られているが、点上正則な自己同型群をもつアフィン平面について類似の方法でとらえることを考えてみたい。以下ではこれが群の特殊な拡大の問題と関連していることを示し、それをもとに得られるいくつかの結果を紹介する。

§2 位数 n^2 のアベル群と位数 n のアフィン平面

A を位数 n^2 のアベル群とし H を位数 n の部分群, $G = A/H = \{1, \sigma, \dots\}$, $\{c(\sigma, \tau)\}$ を拡大 $A > H$ の正規化された因子団の一つとする。すなわち, $\{c(\sigma, \tau)\}$ は A/H の適当な代表系 $D = \{t(\sigma) \mid \sigma \in G = A/H\} \ni 1$ (ただし各 $t(\sigma) \in A$ で $\sigma = t(\sigma)H$) をとれば次をみたす。

$$t(\sigma)t(\tau) = t(\sigma\tau)C(\sigma,\tau), \quad C(\sigma,\tau) \in H$$

$$C(\sigma,\tau)C(\sigma\tau,\rho) = C(\sigma,\tau\rho)C(\tau,\rho)$$

$$C(1,\sigma) = C(\sigma,1) = 1 \quad \forall \sigma, \tau, \rho \in G$$

この因子団 $\{C(\sigma,\tau)\}$ に対して次のような結合構造

$D = \mathcal{D}(C)$ を考える:

D の点集合 $P = A$ の元全体の

D のフローの集合 $B = \{Dx \mid x \in A\} \cup A/H$

D の結合関係 = "e" により自然に定まるもの。

この時、次の(*) が成り立っている。

$$(*) \quad |P| = n^2, \quad |B| = n^2 + n, \quad |B| = n, \quad \forall B \in B$$

このことから $\mathcal{D}(C)$ は位数 n のアフィン平面

($2-(n^2, n, 1)$ デザイン) に近い形をしていることが

分かる。これが実際にアフィン平面であるための

条件が因子団の言葉を用いて次のように言える。

定理 1. $\mathcal{D}(C)$: アフィン平面

$$\Leftrightarrow \{C(\sigma,\tau)\}: \text{bijective } \forall \tau \in G^* = G - \{1\}.$$

(ここで $C(\sigma,\tau)$ を $\sigma \mapsto C(\sigma,\tau)$ により G から H

への写像とみている。)

(証明) 上の(*) より $\mathcal{D}(C)$: アフィン平面 $\Leftrightarrow |B_1 \cap B_2| \leq 1$

Dx ($x \in A$) も A/H の代表系であるから、
($\forall B_1, \forall B_2 \in B$)
(*)

$$\begin{aligned}
& B_1, B_2 \in \{Dx \mid x \in A\} \text{の時以外は常に } |B_1 \cap B_2| \leq 1. \\
\therefore D(c): \text{アフィン平面} & \Leftrightarrow |Dx \cap Dy| \leq 1 \quad \forall x, y \in A \\
& \Leftrightarrow |D \cap Dz| \leq 1 \quad \forall z \in A - \{1\} \\
& \Leftrightarrow |\{(\sigma, \rho) \mid t(\sigma)t(\rho)h = t(\rho) \exists \sigma, \exists \rho \in G\}| \leq 1 \\
& \quad (\forall z = t(\tau)h \quad \tau \in G, h \in H \quad z \neq 1) \\
& \Leftrightarrow |\{(\sigma, \tau) \mid c(\sigma, \tau) = h^{-1} \exists \sigma \in G\}| \leq 1 \\
& \Leftrightarrow \{c(\sigma, \tau)\}: \text{bijective}
\end{aligned}$$

bijective な因子団の例

(1) G, H を位数 n の群とし $f \in G$ から H の中への planar function ([1] §5.1 参照) とするとき $c(\sigma, \tau) = f(\sigma\tau)f(\tau)^{-1}f(\sigma)^{-1}$ は bijective な因子団となる。

(2) $F = F(+, 0)$ を可換な semifield (i.e. “+” に関して群, “ \cdot ” に関して loop, 左右の分配律をみた可代数系) とし $H = G = F^+$ とみる。 $c(\sigma, \tau) = \sigma \cdot \tau$ ($G \times G \rightarrow H$) とおくと bijective な因子団 $\{c(\sigma, \tau)\}$ が得られる。これが定める拡大は $F \times F$ に次のように積を定めたものに同型である: (**)
 $(\sigma_1, \tau_1)(\sigma_2, \tau_2) = (\sigma_1 + \sigma_2, \tau_1 + \tau_2 + \sigma_1 \cdot \sigma_2)$
このとき H は $\{(0, \tau) \mid \tau \in F\}$ と同一視される。また

有限な semifield の加法群 F^+ はある素数 p に対して
基本可換 p 群であることが知られている。 ([2])

$(\sigma_1, \tau_1)(\sigma_2, \tau_2) = (\sigma_2, \tau_2)(\sigma_1, \tau_1)$ が $(**)$ より
確かめられる。従って A は、アベル群であるが。

$(\sigma, \tau)^m = (m\sigma, m\tau + (1+2+\dots+m-1)\sigma \cdot \sigma)$ である
から F の性質 $p\sigma = 0$ ($\forall \sigma \in F$) を用いて

$$A \cong \begin{cases} \text{基本可換 } p \text{ 群} & (p > 2) \\ \text{exponent } 4 \text{ の homocyclic } 2 \text{ 群} & (p = 2) \end{cases}$$

結合構造 $D(C)$ では群 A が点集合 P 上に
可移作用するよう定義されていたが、アフィン平面
の場合にこの逆について考える。

$D = (P, B)$ を位数 n のアフィン平面とし $\text{Aut } D$ が
位数 n^2 の可換部分群 A をもつと仮定する。

D が定める射影平面 \tilde{D} に Orbit Theorem ([2] §13)
を用いることにより A は P または B 上に長さ n^2 の orbit
をもつ。このことから、必要ならば \tilde{D} の中の部分アフィン
平面と取り換えることにより 次のいずれかが起るといえる。

$$(1) \exists B_1, \dots, \exists B_{n+1} \in B, B = B_1^A \cup \dots \cup B_{n+1}^A$$

ここで各 $B_i^A = \{B_i x \mid x \in A\}$ は平行類

$$(2) \exists B_1, \exists B_2 \in B, B = B_1^A \cup B_2^A, B_1^A \text{ は平行類}$$

(1) が起きる時はよく知られているように A は translation group となるのである素数 p に対して基本可換 p 群となりこれはすべての p^{2^m} に対して多くの例が存在する。

従って以下では (2) を仮定しこの場合について考える。

$B_1 \cap B_2 \ni Q$ とする。 A は P 上に正則に作用するから P の任意の点 $Q \in A$ と A の元 x を同一視して B の各元は A の部分集合とみる。 とくに $1 \in B_1 \cap B_2$

$\forall x \in B_1 \rightarrow B_1 \cap B_1 x \ni x \quad \therefore B_1^A$ が平行類であることより $B_1 x = B_1 \quad \therefore B_1$ は A の部分群。

B_1^A は剰余類 A/B_1 に一致。故に B_1^A のどの元も B_2 と一点で交わるということは B_2 が A/B_1 の代表系であることを意味する。 $\therefore H = B_1, D = B_2$ とおくと

$A > H$ (部分群), $A \supset D$ (部分集合)

$|A| = n^2, |H| = |D| = n, A = HD$

さらに次が成り立つ

定理 2 D を拡大 $A > H$ の代表系として選ぶとき対応する因子団は bijective である。

(証明) $G = A/H = \{1, \sigma, \dots\}, D = \{t(\sigma) \mid \sigma \in G\}$ とおく。ここで $t(\sigma) \in A, t(\sigma)H = \sigma$ とする。

$B \in B_2^A$ とおると B を含む平行類は B^H であるから

$G \ni z \neq 1$ を fix すると $|D \cap Dt(z)H| = 1 \quad \forall z \in H.$

$\therefore |\{(\sigma, \rho) \mid t(\sigma) = t(\rho)t(z)h\}| = 1 \quad \forall z \in H.$

因子団 $\{C\}$ は $t(\rho)t(\tau) = C(\rho, \tau)t(\rho\tau)$ に依り
定義されるから

$|\{(\rho\tau, \rho) \mid C(\rho, \tau) = h^{-1}, \rho \in G\}| = 1 \quad \forall h \in H$

$\therefore \{C\}$ は bijective.

以上のことにより 次の対応が示された.

{アーベル群から得られる bijective な因子団の全体}

$\downarrow \uparrow$

{点上可移な可換自己同型群(除 translation gps)}

をもつアフィン平面の全体}

§3 bijective な因子団の性質といくつかの応用.

定理3. A を群 H の群 G による中心拡大とし.

$\{t(\sigma)\}$ を A/H の代表系 ($t(1)=1$), $\{C(\sigma, \tau)\}$ を対応する因子団とする. G のある元 ρ が次の条件を満たせば.

1. $\langle H, t(\rho) \rangle$ は H 上の分裂拡大である:

$$\prod_{m=0}^{o(\rho)-1} C(\rho, \rho^m) = 1 \quad (o(\rho) = \rho \text{ の位数})$$

(証明) $\psi(j) = C(\rho, \rho^j) \quad j=0, \pm 1, \dots$ とおく.

$$C(\rho^i, \rho^j)C(\rho^i\rho^j, \rho^k) = C(\rho^i, \rho^j\rho^k)C(\rho^j, \rho^k) \quad \forall i, j, k$$

$$\begin{aligned} \therefore \text{帰納法により } C(\rho^p, \rho^q) &= \left(\prod_{s=0}^{p-1} \psi(\rho+s) \right) \left(\prod_{t=0}^{q-1} \psi(t) \right)^{-1} \\ &= \left(\prod_{s=0}^{p+q-1} \psi(s) \right) \left(\prod_{s=0}^{q-1} \psi(s) \right)^{-1} \left(\prod_{t=0}^{p-1} \psi(t) \right)^{-1} \text{ が成り立つ.} \end{aligned}$$

$$\begin{aligned} \therefore C(\rho^p, \rho^q) &= \Gamma(p+q) \Gamma(p)^{-1} \Gamma(q)^{-1} \\ &\quad \text{ここで } \Gamma(m) = \prod_{s=0}^{m-1} \psi(s) \end{aligned}$$

$\langle \rho \rangle$ から H への写像を $\Delta(\rho^m) = \Gamma(m)^{-1}$ と

定義するとこれは well-defined. なぜなら

$$\begin{aligned} r=0(\rho) \text{ とおくと } \Gamma(\rho^{m+r}) &= \prod_{s=0}^{m-1} \psi(s) \prod_{s=m}^{m+r-1} \psi(s) \\ &= \Gamma(m) \prod_{t=0}^{r-1} C(\rho, \rho^t) = \Gamma(m) \end{aligned}$$

$$\therefore C(\rho^p, \rho^q) = \Delta(\rho^p \rho^q)^{-1} \Delta(\rho^p) \Delta(\rho^q)$$

つまり $C(*, *)$ は coboundary, および $\langle H, t(\rho) \rangle$ は H 上 split する.

定義 A, H, G を定理 3 と同じ条件にとる. 因子団

$C(\sigma, \tau)$ が $\prod_{\tau \in G} C(\sigma, \tau) = 1$ ($\forall \sigma \in G$) をみたすとき

$C(\sigma, \tau)$ を homogeneous な因子団という.

補題 4 $\{C(\sigma, \tau)\}$ が bijective な因子団ならば

$\{C(\sigma, \tau)\}$ は homogeneous か又は $n=2$ ($A \cong \mathbb{Z}_4$)

が成り立つ.

(証明) A は中心拡大であるから H は p -ベリ群であるが

p -ベリ群の元すべての積が H の 2-Sylow 群 S が唯一つの

位数 2 の元を含む場合を除き単位元であるという事を

用いて示すことができる。 $S \neq 1$, かつ S が cyclic の時は
定理1を用いてアフィン平面を考えて、アフィン平面の中で
考えて $n=2$ を示す。

定理5 $\{C(\sigma, \tau)\}$ を homogeneous, p を奇素数,
 P を A の p -Sylow 群とする。もし $P/P \cap H$ が
cyclic ならば P は $P \cap H$ 上分裂する。とくに
 A の各 p -Sylow 群 ($p \neq 2$) は cyclic でない。

定理6 $\{C(\sigma, \tau)\}$ を bijective, P を A の 2-Sylow
群とする。このとき次が成り立つ

(1) A の位数 2 の元はすべて P に含まれる。

(2) $P \cap H$ が cyclic ($\neq 1$) ならば $n=2$ ($A \cong \mathbb{Z}_4$).

(定理5, 6 は Hoffman の定理 ([1] p210) の
一般化になっている。)

以上の結果は A が p -ベル群でなくても成り立つもの
が多い (定理1, 2 など)。また、 $2 \nmid |A|$ の時は、 $\{C(\sigma, \tau)\}$
は分裂因子団であることが予想される。 $2 \parallel |A|$ の時は
 A が 2 群になると予想されるがこの場合は前者よりやさしい
ような気がする。一般にアフィン平面 (射影平面) の
位数は素数中に等しいと予想されているが、その根拠

は今かりにこのように思う。可移な自己同型群をもつアフィン平面の場合は納得させるに足る理由があるように思われる。([3] に詳しい) A が P -ベル群の場合は証明できてもよいような気がする。

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Splitting Fields of Association Schemes

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定義 d -class association scheme $(X, \{R_i\}_{i=0}^d)$

とは、有限集合 X と $X \times X$ の分割 $\{R_i\}_{i=0}^d$ で、

$$(i) \quad R_0 = \{(x, x) \mid x \in X\}$$

$$(ii) \quad \forall i, \exists i' \text{ st } \{(x, y) \mid (y, x) \in R_i\} = R_{i'}$$

$$(iii) \quad \forall i, j, k, \forall (x, y) \in R_k,$$

$$P_{ij}^k = \#\{z \in X \mid (x, z) \in R_i, (z, y) \in R_j\} \text{ は } (x, y)$$

によらず一定

$$(iv) \quad P_{ij}^k = P_{ji}^k \quad \forall i, j, k.$$

adjacency matrix A_i を次のように定義する

$$(A_i)_{xy} = \begin{cases} 1 & (x, y) \in R_i \\ 0 & (x, y) \notin R_i \end{cases}$$

\mathcal{A} を A_0, A_1, \dots, A_d で \mathbb{C} 上生成される線形空間と

すると、 \mathcal{A} は可換な \mathbb{C} 代数になる。 \mathcal{A} を adjacency algebra と呼ぶ。 \mathcal{A} は $M_n(\mathbb{C})$ ($n=|X|$) の subalgebra と考えると自然に \mathbb{C}^n に作用する。 \mathcal{A} の極大共通固有空間は丁度 $d+1$ 個あることがわかり、そのうちのひとつは $(1,1,\dots,1)$ で生成される 1 次元空間である： $\mathbb{C}^n = V_0 \oplus V_1 \oplus \dots \oplus V_d$, $V_0 = (1,1,\dots,1)$, \mathbb{C}^n から V_i への射影を E_i とおく。特に $E_0 = \frac{1}{n}J$, J は all 1 matrix である。すると

$$A_j = \sum_{i=0}^d P_j(i) E_i \quad (1)$$

と書ける。ただし $P_j(i)$ は A_j の V_i 上での固有値である。

$P = (P_j(i))$ を character table と呼び、 $K = \mathbb{Q}(P_j(i),$

$0 \leq i, j \leq d) \subset \mathbb{C}$ を splitting field と呼ぶ。

問題 K は \mathbb{R} 分体に含まれるか。

つまり、association scheme の character table が 1 の中根で書けるかということだが、この問題は未解決である。association scheme の代表的な例として、有限群の等質空間があるが、この場合には、その置換表現に現れる指標の値を使って $P_j(i)$ を表せるので、確かに \mathbb{R} 分体に含まれている。

さて、(1)式は $\{E_i\}_{i=0}^d$ について解くことができ、

$$\mathcal{A} = \langle A_0, A_1, \dots, A_d \rangle = \langle E_0, E_1, \dots, E_d \rangle.$$

となる。従って、 \mathcal{A} は行列の乗法だけでなく、行列の成分ごとの乗法 (Hadamard 積) についても閉じていることがわかる。

$$E_i \circ E_j = \frac{1}{n} \sum_{k=0}^d q_{ij}^{(k)} E_k$$

と書いたとき (\circ は Hadamard 積) $q_{ij}^{(k)}$ を Krein parameter と呼ぶ。 $L = Q(q_{ij}^{(k)}, 0 \leq i, j, k \leq d)$ とおくと $L \subset K$ である。さらに次のことがわかる。

定理 $\text{Gal}(K/L) \subset Z(\text{Gal}(K/\mathbb{Q}))$

証明の概略を述べよう。 $\sigma \in \text{Gal}(K/\mathbb{Q})$ とすると、 σ は \mathcal{A} の自己同型を引き起こし、primitive idempotents $\{E_i\}_{i=0}^d$ を置換する。 $\tau \in \text{Gal}(K/L)$ とすると τ は Krein parameter を fix するので、Hadamard 積に関する primitive idempotents $\{A_j\}_{j=0}^d$ を置換する。これより、

$$\begin{aligned} P_j(i)^{\sigma\tau} &= P_j(i^{\psi(\sigma)})^{\tau} = P_j(\psi(\tau))(i^{\psi(\sigma)}) \\ &= P_j(\psi(\tau))(i)^{\sigma} = P_j(i)^{\tau\sigma} \end{aligned}$$

(ただし $\psi(\sigma), \psi(\tau)$ は $\{0, 1, \dots, d\}$ の置換) となり、 $\sigma\tau = \tau\sigma$

が成り立つ。すなわち $\tau \in Z(\text{Gal}(K/\mathbb{Q}))$ 。

系 Krein parameter が有理数ならば、 K は円分体に含まれる。

実際、定理から、 $L = \mathbb{Q}$ なら $\text{Gal}(K/\mathbb{Q})$ が abelian であることがわかり、Kronecker-Weber の定理により、 K が円分体に入ることがわかる。

$L = \mathbb{Q}$ でない association scheme は、等質空間の場合にもあり得るが、筆者の知る限り、 L は 2 次体が 4 次体である。Krein parameter は、character table の成分から計算する公式があるので、character table が与えられれば Krein parameter が有理数かどうか判定できる。定理の証明に使われている議論を応用すると、この判定が容易にできる。この方法を示す例をあげよう。 $\text{PGL}(3, 2)$ は、 $X = \{(x, \ell) \mid x \text{ は射影空間 } \text{PG}(2, 2) \text{ の点, } \ell \text{ は直線, } x \notin \ell\}$ の上に可移に作用し、Coxeter graph と呼ばれる distance-transitive graph (従って association scheme) ができる。その character table は

$$P = \begin{pmatrix} 1 & 3 & 6 & 12 & 6 \\ 1 & -1+\sqrt{2} & -2\sqrt{2} & -2 & 2+\sqrt{2} \\ 1 & -1-\sqrt{2} & 2\sqrt{2} & -2 & 2-\sqrt{2} \\ 1 & 2 & 1 & -2 & -2 \\ 1 & -1 & -2 & 4 & -2 \end{pmatrix}$$

従って $K = \mathbb{Q}(\sqrt{2})$ である。 $\sigma: \sqrt{2} \mapsto -\sqrt{2} \in \text{Gal}(K/\mathbb{Q})$ は、 P の行の置換を引き起こす ($\{e_i\}_{i=0}^d$ の置換) が、列の置換を引き起こさない。つまり Hadamard 積を保存しない、Krein parameter を fix しない、ということがわかる。このようにして、 P の形を見ただけで Krein parameter が有理数がどうか判定できるのである。

一方、もし K が \mathbb{Q} 分体に含まれないような association scheme があるとしたら、どのような形をしているだろうか。 $\text{Gal}(K/\mathbb{Q})$ は非可換でなければならぬ。手始めとして、 $\text{Gal}(K/\mathbb{Q}) = S_3$ とする例はあるだろうか。 A_j の最小多項式は $\prod_{i=0}^d (x - P_j(i))$ で $P_j(i)$ は整数だから、3次の既約成分を持つためには $d \geq 3$ でなくてはならない。 $d=3$ の時を考えると $\dim V_1 = \dim V_2 = \dim V_3$

となり P_{ij}^k は次の形になる

$$B_1 = \begin{pmatrix} 0 & 1 & 0 & 0 \\ a+b+c & a-1 & b & c \\ 0 & b & c & a \\ 0 & c & a & b \end{pmatrix} \quad B_2 = \begin{pmatrix} 0 & 0 & 1 & 0 \\ 0 & b & c & a \\ a+b+c & c & a-1 & b \\ 0 & a & b & c \end{pmatrix}$$

$$B_3 = \begin{pmatrix} 0 & 0 & 0 & 1 \\ 0 & c & a & b \\ 0 & a & b & c \\ a+b+c & b & c & a-1 \end{pmatrix}$$

$\forall i \in \mathbb{L} \quad P_{ij}^k = (B_i)_{jk}$, a, b, c は $a^2+b^2+c^2-ab-bc-ca = a$ を満たす整数, character table P は

$$P = \begin{pmatrix} 1 & m & m & m \\ 1 & \theta_1 & \theta_2 & \theta_3 \\ 1 & \theta_2 & \theta_3 & \theta_1 \\ 1 & \theta_3 & \theta_1 & \theta_2 \end{pmatrix}$$

$\forall i \in \mathbb{L} \quad m = a+b+c$, $\{\theta_1, \theta_2, \theta_3\}$ は 3次方程式 $x^3 + x^2 - mx + bc - a^2 = 0$ の3根である。(cf.

Brouwer-Cohen-Neumaier, "Distance-regular graphs", Lemma 12.7.4 on p.389). 実は上の3次方程式の判別式は平方であることがわかって、Galois群は S_3 でないことがわかる。

ところで、上にあげた本には、この3次方程式が可約

となる必要十分条件は $b=c$ である、と書かれているが、これは間違いである。実際、

$$a = 343t^2 + 464t + 157$$

$$b = 343t^2 + 463t + 156$$

$$c = 343t^2 + 445t + 144$$

とおくと

$$(x+7t+5)(x+28t+19)(x-35t-23)$$

と因数分解される。 $t=-1$ のとき association scheme は実在し、それは $GF(7^3)$ の 3 乗剰余からつくられる cyclotomic scheme と呼ばれるものである。

一般に $GF(q)$ 上の e -class cyclotomic scheme とは association scheme $(X, \{R_i\}_{i=0}^e)$ で、 $X=GF(q)$, $GF(q)^* = \langle \theta \rangle$, $C_i = \langle \theta^e \rangle \theta_i$ (e 乗剰余), $R_i = \{(x, y) \mid x-y \in C_i\}$ ($i=1, 2, \dots, e$) で定義される。 cyclotomic scheme の splitting field が \mathbb{Q} に存するのはどのような時かという質問を提示したところ、時村良雄氏(神戸大)、山本幸一氏(東京女子大)から回答があり、それらを一般化して次の結果が得られた。

定理 $GF(p^r)$ 上の e -class cyclotomic scheme (p は素数, $p^r - 1 = ef$, $e > 1, f > 1$) の splitting field は $\mathbb{Q}(\zeta_p)/\mathbb{Q}$ の次数 $(p-1)/(p-1, f)$ の中間体である。

特に $GF(7^3)$ 上 3-class のときは $7^3 - 1 = 3 \cdot 114$ で $(p-1)/(p-1, f) = 6 / (6, 114) = 1$ であり splitting field は \mathbb{Q} である。一般に cyclotomic scheme の character table は次のように書ける

$$P = \begin{pmatrix} 1 & f & f & \dots & f \\ 1 & \chi(\underline{C}_1) & \chi(\underline{C}_2) & \dots & \chi(\underline{C}_e) \\ 1 & \chi(\underline{C}_e) & \chi(\underline{C}_1) & \dots & \chi(\underline{C}_{e-1}) \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ 1 & \chi(\underline{C}_2) & \chi(\underline{C}_3) & \dots & \chi(\underline{C}_1) \end{pmatrix}$$

ここで、 χ は nontrivial な $GF(p)$ の加法的指標、

$$\chi(\underline{C}_i) = \sum_{a \in \underline{C}_i} \chi(a)$$

で、これは Gauss の周期と呼ばれるものである。

On spherical t -designs (a survey)

坂内 英一 (九大・理)

§1. Introduction

この講演では spherical t -design に関する最近の
仕事についての survey を与える。この方面の研究は
日本ではあまりなじみがないと思われたので、証明など
の技術的なことの解説よりも、どのような問題が考えら
れてきたか、またどのような文献があるかということを
主に解説する。(この報告集では講演で述べられなかつ
たことも少し補足してある。) この方面に興味を持ち、
研究を始められる方の役に立てば幸いである。

$$S^d = \{(x_1, \dots, x_{d+1}) \in \mathbb{R}^{d+1} \mid x_1^2 + \dots + x_{d+1}^2 = 1\} \subset \mathbb{R}^{d+1}$$

を単位球とする。 S^d の有限部分集合 X についての
"Combinatorics" を研究したいわけである。 S^d の有
限部分集合 X の研究には色々の方向があるであろう。
コーディング理論の立場からは (大雑把に言えば) X の 2
点間の距離 $d(x, y)$ 達に注目し、例えば $d(x, y)$ が

特別な値を取り (あるいは $\text{Min } d(x, y)$ が出来ただけ大きく) $|X|$ が出来ただけ大きくなったものに面白がある。デザイニ理論の立場からは、単位球 S^d を有限個の点で "良く近似するもの" に面白がある。次の S^d における t -design の概念は Delsarte-Goethals-Seidel (1977) により導入された非常に自然なかつ非常に役立つ概念である。

定義 (Delsarte-Goethals-Seidel (1977)) 単位球 S^d の有限部分集合 X が spherical t -design (または t -design in S^d) であるとは、

$$\frac{1}{|S^d|} \int_{S^d} f(x) dx = \frac{1}{|X|} \sum_{x \in X} f(x)$$

for $\forall f(x) = f(x_1, \dots, x_{d+1})$: polynomials of degree $\leq t$ が成り立つことを定義する。

(この条件は次の条件

$$\sum_{x \in X} f(x) = 0 \quad \text{for } \forall f \in \text{Harm}(i), \quad 1 \leq i \leq t,$$

と同値である。ここで $\text{Harm}(i)$ は x_1, \dots, x_{d+1} の i -次の homogeneous harmonic polynomials の \mathbb{R} 線形空間を表わす。 $L^2(S^d) = \bigoplus_{i \geq 0} \text{Harm}(i)$ であり、

$$\dim \text{Harm}(i) = \binom{d+i}{i} - \binom{d+i-2}{i-2} \quad \text{となる。}$$

従って X が spherical t -design であるか否かは有限個の多項式 $f(x)$ に関して check すればよい条件である。

Spherical t -design の基本的な性質についてはこのまゝと詳しいことは、Delsarte-Goethals-Seidel (1977) を、あるいは、survey paper: Bannai (1988) を参照したい。

S^d の有限部分集合の "Combinatorics" については重要な結果を要約する。

定理 (Delsarte-Goethals-Seidel (1977)) X が S^d の t -design ならば、

$$|X| \geq \begin{cases} \binom{d + \lfloor \frac{t}{2} \rfloor}{\lfloor \frac{t}{2} \rfloor} + \binom{d + \lfloor \frac{t}{2} \rfloor - 1}{\lfloor \frac{t}{2} \rfloor - 1} & \text{for } t = \text{even} \\ 2 \cdot \binom{d + \lfloor \frac{t}{2} \rfloor}{\lfloor \frac{t}{2} \rfloor} & \text{for } t = \text{odd} \end{cases}$$

(Fisher 型不等式)

(上で等号の成り立 $\Rightarrow X$ は tight t -design と呼ぶ。)

定理 (Bannai-Damerell (1979, 80)) Tight t -design in S^d ($d \geq 2$) が存在するならば、 $t = 2, 3, 4, 5, 7$ or 11 である。

注意 $t = 4, 5, 7$ に対する tight t -design in S^d の分類は依然未解決である。

S^d における t -design X の色々の例は $O(d+1)$ の有限部分群 G の orbits : $X = \bar{x}^G$ with $\bar{x} \in S^d$ と

なるものを中心に色々と調べられた。一方、 $d \geq 2$ の時、
大きい t に対して t -design in S^d が存在するが否
かは、次の Seymour-Zaslavsky の定理が得られた。未
解決であった。

定理 (Seymour-Zaslavsky (1984)) $\forall t, \forall d$ に対
して t -design X in S^d が存在する。

この Seymour-Zaslavsky の結果は非常に一般的な結
果であり、上の定理はその極く特別な場合である。この
論文で3種類の証明が与えられている (例えば一番易し
い証明は implicit function theorem (陰函数定理)
を用いる) が、いずれも完全に non-constructive な証
明である。例えば t と d を与えたとき、 X の存在は
わかるが、 $|X|$ がどの位小さく出来るか (小さい程
良い) は一般的には全く計算不可能であった。

従って、残された問題として、与えられた t と d に対
して、どの位小さい $|X|$ を持つ t -design X in S^d
が存在するかが決まるか、また与えられた X をどの
ように explicit に構成することが出来るか、といったこと
が重要な問題となる。次の § でこの問題に関しての
最近の進展を述べる。

§2. Spherical t -design に関する最近の結果

この節では survey paper: Bannai (1988) 以後の spherical t -design についての極く最近の研究について述べる。

$S^1 = \text{unit circle } (\subset \mathbb{R}^2)$ に内接する正 $(t+1)$ -角形の頂点は t -design in S^1 である。Fisher型不等式は $|X| \geq t+1$ なる t -design in S^1 に対して、正 m -角形 ($m \geq t+1$) を考えることにより、 $|X| \geq t+1$ とする任意の $|X|$ に対して t -design in S^1 が explicit に構成出来る。また $t=1$ の時は問題は trivial とする。以下 $t \geq 2$ の $d \geq 2$ の場合のみを考える。

$t=2$ の時、Fisher型不等式は $|X| \geq d+2$ とする。 $(S^d$ に内接する regular simplex の頂点が $|X|=d+2$ (tight) の例を与える。) $t=2$ に対する完全な結果は、1987/8 に Ohio に滞在していた味村良雄氏 (神戸大) により得られた。

定理 (Mimura (to appear))

- (i) $d = \text{even}$, $|X| = d+3$ の 2-design X in S^d は存在しない。
- (ii) 上の例外を除き、任意の $|X| \geq d+2$ に対して、

2-design X in S^d が存在する (かつ explicit に構成される)。

注意 $t \geq 3$ の時はこの種の完全な結果は未解決である。他の小さな t の値に対しての味村氏の仕事の拡張の試みは、Bela Bajnok により成された。

定理 (Bajnok (1989, to appear)) (結果の要約)
 $|X| \geq (d+1) 2^{d+2}$ とする任意の $|X|$ に対して、5-design X in S^d が存在する (かつ explicit に構成される)。

(注意: Fisher 型 bound は $|X| \geq 2 \binom{d+2}{2} \approx d^2$)

定理 (Bajnok (1989)) 上と同様の結果が $t=7$ に対して得られる。[ただし $|X|$ は大きく、 $|X|$ がある合同条件を満たさなければならず、また X が multi-set を許す (i.e., X が同じ要素 2 個以上含む可能性 — ブロウワー・テグイシの理論における repeated blocks の概念に対応する) 可能性が完全に排除されていないので $t=5$ の時と比べてとまあ少し弱い結果である。]

注意 この方向を更に進めて、一般の t に対して同様の結果が得られることが望まれた (がまだ未解決である)。

Spherical t -design と似た (あるいは意味が易しく、また群の表現論が使えないという) 意味では難かしくかつか) 概念である次の interval design の概念は興味深い。

定義 (実区間 $[-1, 1]$ で考えよ — 他の区間でもよいか) 有限集合 $X \subset [-1, 1]$ が t -design in $[-1, 1]$ (interval t -design) であるとは、

$$\frac{1}{|X|} \sum_{x \in X} f(x) = \frac{1}{2} \int_{-1}^1 f(x) dx$$

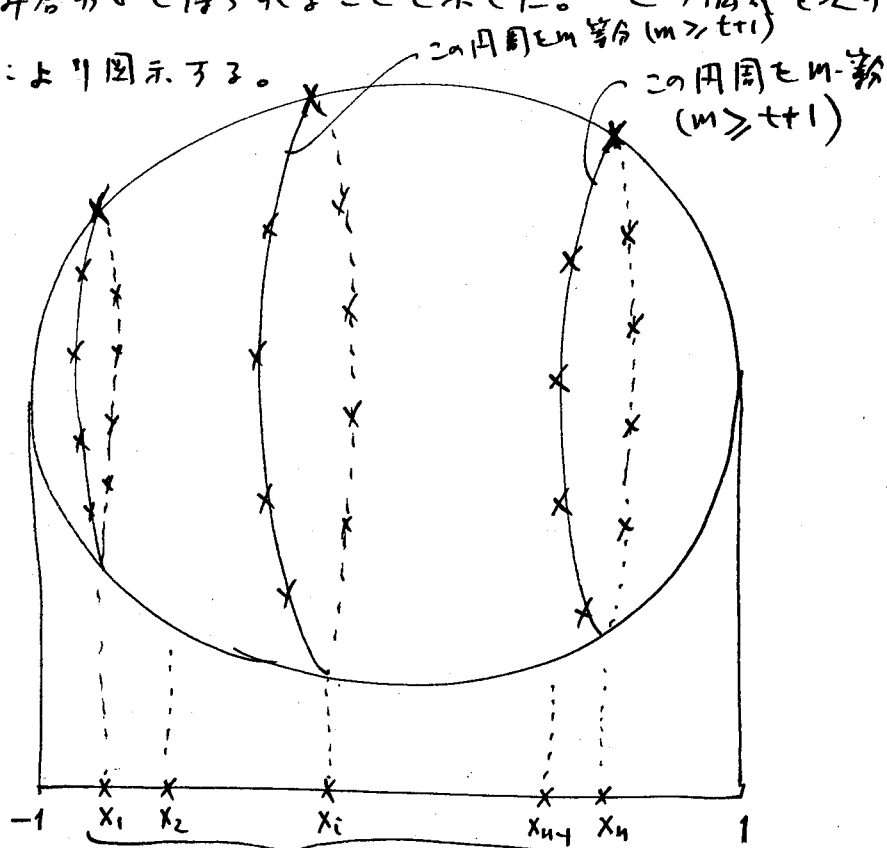
for \forall polynomial $f(x)$ of degree $\leq t$ と定義する。

注意 この概念は "quadrature formula with equal weight" の存在と同値であり、解析 (主に approximation theory) で古くから研究されて来た。 X が interval t -design の時 $|X| \geq t+1$ が成り立ち、tight (i.e. $|X| = t+1$) となるのは $t \leq 7$ かつ $t = 9$ の時に限ることから S. Bernstein に於て 1940 年 4 月に証明された。 (Krylov (1962), Natanson (1965) 参照。) 更に Bernstein は、ある absolute constant c で X が interval t -design $\Rightarrow |X| \geq c \cdot t^2$ とするものが存在することを示した。 (Krylov (1962), Natanson (1965) 参照。) 従って $t \rightarrow \infty$ の時 $|X|$ は

tight の bound より はるかに大きくなる。

注意 Seymour-Zaslavsky の存在定理は、一番初めに (私が Ohio の seminar で提出した) interval t -design の存在問題を解いたのが始まりであり、それが spherical t -design を含む一般の存在定理に直ちに拡張されたわけである。

Bajnok (1989) は S^2 の spherical t -design を S^1 の spherical t -design と interval t -design を組み合わせさせて得られることを示した。この構成を次の図により図示する。



interval t -design
(t -design X は nm 個の点からなる)

注意 (何らかの形で) 類似なことが S^d ($d \geq 3$) の場合に言えよと望ましいが未解決である。

de Reyna (1988) は interval t -design の存在 (Seymour-Zaslavsky の定理の特別な場合) の別証明を与えた。これは 3 ページの短い論文であり、トポロジー (ホモトピー論) の極く基本的なことのみに用いる。私はこの証明を用いて $|X|$ を具体的に計算出来よ可能性に気付く。Bajnok はこれを実行して、(ある具体的な定数 C_i ($i=1,2$) が存在して) 任意の $n \geq C_1 \cdot t^{9/2}$ に対して $|X|=n$ とする interval t -design X が常に存在する事を証明した。従って Bajnok の前に述べた結果から、任意の $n \geq C_2 \cdot t^{1/2}$ に対して、 $|X|=n$ とする t -design X in S^2 が常に存在する事がわかる、たしかである (Bajnok (in preparation))。

極く最近 (1989年7月) Wagner (to appear) の preprint を見た。そこでは interval t -design の explicit な存在についての別証明を与えた。(この結果は de Reyna-Bajnok の bound よりも悪くは思われる)

子。) その論文の中で、任意の $n \geq C_d \cdot t^{12d^4}$ に
 対して、 $|X|=n$ とする t -design X in S^d が常に存
 在することを announce された (ただし詳細は
 今の所不明である。) 以下にせよ、存在かわか、この
 size $|X|$ は tight の bound よりず、と大きい。
 この間をどの位埋めたいのか、また explicit な t -design
 がどのように構成されたのか、今後の興味深い問題
 であろう。

§3. 補足

1. Spherical t -design は n が大きいとき存在するわけ
 であるが、これは特別に興味深いのはどのような
 だろうか。私自身は rigid t -design の概念が
 好き。その分類問題は非常に重要になり
 考えたい。

(詳しくは、Bannai (1988, 1988bis) 等を参照。)

2. sphere の中の t -design の概念は compact
 rank 1 symmetric space の t -design に拡張され、
 色々な類似の結果がそこで証明されている。
 (Bannai (1988), Bannai-Hoggar (1989) 等を参照。)

3. \mathbb{R}^d (または non-compact rank 1 symmetric

space) の有限部分集合 X の "combinatorics" は S^d の場合にくらべて難しくなりました。 \mathbb{R}^d における t -design X もどう定義するかは、Neumaier-Seidel (1988), DelSarte-Seidel (1989) に一つの試みがあります。これはおもしろい論文ですが、まだ最終的な結果ではないと思います。さらに、彼らの定義は、 X が " t -design in \mathbb{R}^d " であるとは、

$$\sum_{x \in X} (x, x)^{\ell} h_i(x) = 0 \quad \text{for } \forall h_i(x) \in \text{Harm}(i)$$

with $2\ell + i \leq t$ ($i \geq 1$) となることである。(ただし、 $(x, x) = x_1^2 + \dots + x_d^2$.)

X が \mathbb{R}^d の s -distance set の時、 $|X| \leq \binom{d+s}{s}$ が知られている (Blokhuis (1984), Bannai-Bannai-Stanton (1983)). t -design in \mathbb{R}^d の最上の定義としては、 $|X| = \binom{d+s}{s}$ を満たす s -distance set X in \mathbb{R}^d を $2s$ -design になるか一番望ましいが、先の定義ではそれは多分真ではないであろう。野田隆三郎氏は最近、 $X \subset \mathbb{R}^d$ が $|X| = \binom{d+2}{2}$ なる 2 -distance set ならば

$$\sum_{x \in X} (x, x)^{\ell} h_i(x) = 0 \quad \text{for } \forall h_i(x) \in \text{Harm}(i)$$

with $\lambda + i \leq 2$ ($i \geq 1$) を示された。(これが何を意味するか、またこの結果が \mathbb{R}^d における t -design の最終的定義を見つけたのに役立ったか否かは今の所不明であるが、何かの役に立つ可能性はありと思う。

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TRIANGULATIONS WITH WEIGHTED VERTICES

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Abstract. If three points A, B, C, are given in the plane, then the right bisectors of the sides of triangle ABC divide the plane up into regions that "belong" to the three vertices. If weights are attached to the vertices, then the regions obtained are bounded by circular arcs, and the conditions for intersection can be found by an application of the Heron formula for the area of a triangle. A series of diagrams illustrate various configurations obtained in some practical applications in the field of forestry.

1. Introduction. In the biological sciences, the occurrence of competition models is frequent (cf. Daniels, Burkhardt, and Clason [1]); for a description of various different competition designs, we refer to Street and Street [2]). As one instance, we may take three points A, B, C, in the plane and imagine them to represent three trees that are competing for nutrients and water from the soil. If the trees are of equal size, we can reasonably postulate that they have equal attractive strengths, and it is natural to use the right bisectors of AB, AC, and BC to divide the plane up into three regions that "belong" to the three trees A, B, C; the circumcentre of triangle ABC then becomes the point at which the attractive forces of the three trees are equal. We shall look at the case when the trees are of different sizes, and consequently their attractive strengths are different; we model this situation by attaching positive weights u , v , and w , to points A, B, and C. There is no loss in generality in assuming that the points are labelled in such a way that u, v , and w are in decreasing magnitude, that is, $u \geq v \geq w > 0$.

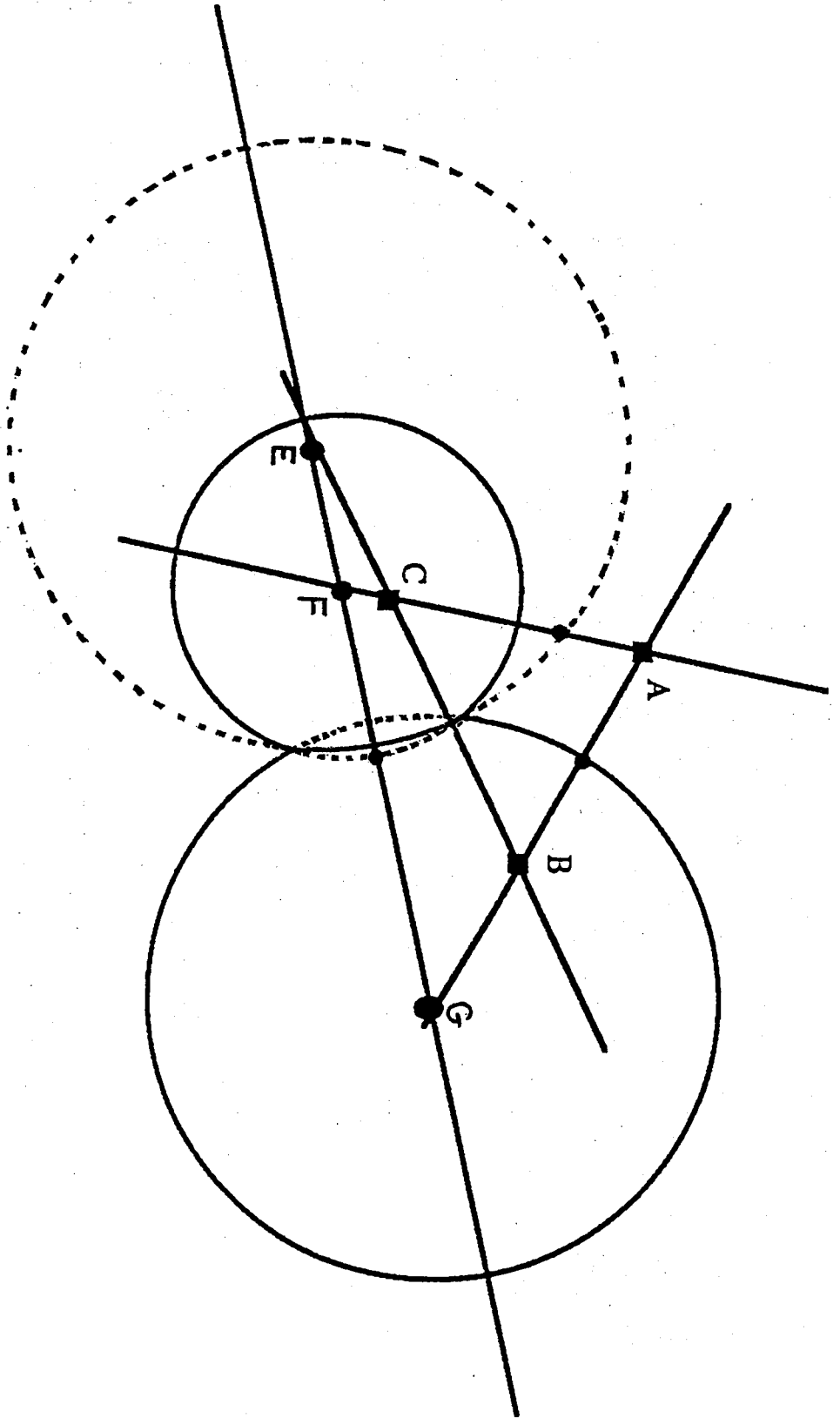
2. Discussion of the Model. In Figure 1, the triangle ABC is arbitrary (the three vertices can be considered as representing the positions of three trees). Weights u , v , and w , are attached to the vertices A, B, and C, respectively. It is assumed that all the weights are positive and that $u \geq v \geq w$. If the co-ordinates of A, B, and C are represented by (x_1, y_1) , (x_2, y_2) , and (x_3, y_3) , respectively, then the locus of points P_1 such that $AP_1/BP_1 = u/v$ is given by

$$v^2 [(x - x_1)^2 + (y - y_1)^2] = u^2 [(x - x_2)^2 + (y - y_2)^2].$$

This locus has the general form

$$K_1(x^2 + y^2) + K_2x + K_3y + K_4 = 0,$$

Figure 1



and hence is a circle whose centre lies on the line AB. Similarly, we find that the locus of points P_2 such that $BP_2/CP_2 = v/w$ is a circle whose centre is on the line BC and whose equation is

$$w^2 [(x - x_2)^2 + (y - y_2)^2] = v^2 [(x - x_3)^2 + (y - y_3)^2].$$

Finally, the locus of points P_3 such that $CP_3/AP_3 = w/u$ is a circle whose centre is on the line AC and whose equation is

$$u^2 [(x - x_3)^2 + (y - y_3)^2] = w^2 [(x - x_1)^2 + (y - y_1)^2].$$

These three circles are just the circles of Apollonius associated with each of the pairs of points. So we may state our first result as

Lemma 1. The loci partitioning the plane into areas belonging to A, B, and C (in the sense that the attractive forces from A, B, and C are dominant in the areas) are simply arcs of circles.

Figure 1 illustrates such a situation; the partitioning arcs are shown as solid lines, whereas the remainder of each circle is shown as a dotted arc. In this diagram, the circles intersect in two real points, that is, there are two points K such that $AK : BK : CK = u : v : w$ (zero points of intersection and one point of contact are also possibilities)

We next note that the equations of any two of the three circle can be combined to give the other equation; this means that any points common to two of the circles also lie on the third circle. Consequently, the three circles all possess the same common chord. Since the centres of the three circles lie on the common chord, we have

Lemma 2. The centres of the three circles are all collinear.

It is interesting to note that this result can be derived in another way. If we denote the centres of the three circles by G, E, and F, respectively, then it is easy to calculate that

$$AG/GB = u^2/v^2, \quad BE/EC = v^2/w^2, \quad CF/FA = w^2/u^2.$$

It thus follows that the points G, E, and F divide the sides of the triangle in such a way that the product of AG/GB , BE/EC , and CF/FA is unity. Consequently, the Theorem of Menelaus guarantees that G, E, and F all lie upon a straight line.

It is clear that, as the ratio $u : v : w$ approaches the ratio $1 : 1 : 1$, then the circles approach circles of infinite radius (straight lines). One of the points of intersection recedes to infinity, as does the line GEF. The other point of intersection becomes the intersection of the right bisectors of the three sides, that is, the circumcentre of triangle ABC.

3. Condition for Real Intersections. Figure 1 illustrates the situation when the circles have two real points of intersection. These two points may coincide or they may be non-real. The case when the two points coincide is clearly the dividing line between the real and non-real situations. We shall now investigate the algebraic conditions for real intersections; the conditions take a particularly symmetric form if we employ the Heron formula for the area of a triangle.

First, it is useful to obtain a somewhat different form for the Heron formula; normally, this formula states that the area Δ is given by

$$\Delta^2 = s(s - a)(s - b)(s - c),$$

where a, b, c , are the lengths of the sides opposite A, B, C , respectively, and where $2s = a + b + c$. If we substitute for s , we have

$$16\Delta^2 = (a+b+c)(a+b-c)(a-b+c)(b+c-a);$$

from this form of the equation, we note that failure of the triangle inequality (that is, a negative value for one term such as $a+b-c$) shows up as a negative value for Δ^2 (or, equivalently, as a non-real value for Δ). Further algebra produces the alternative formula

$$16\Delta^2 = 2(a^2b^2 + a^2c^2 + b^2c^2) - a^4 - b^4 - c^4.$$

It will also be useful to introduce an auxiliary triangle whose sides are au, bv , and cw . The area of this auxiliary triangle is a quantity T where

$$16T^2 = 2[(au)^2(bv)^2 + (au)^2(cw)^2 + (bv)^2(cw)^2] - (au)^4 - (bv)^4 - (cw)^4.$$

We now consider Figure 2, where K represents a point of intersection of the circles; then the distances AK, BK , and CK are given by ku, kv , and kw , respectively, where k is a constant of proportionality that remains to be determined. The cosine formula immediately gives us the results that

$$\cos AKB = [k^2(u^2 + v^2) - c^2]/2k^2uv,$$

$$\cos BKC = [k^2(v^2 + w^2) - a^2]/2k^2vw,$$

$$\cos CKA = [k^2(u^2 + w^2) - b^2]/2k^2uw.$$

However, the angles AKB, BKC , and CKA add to 360° , and it is well known that, for any angles P, Q , and R that add to 360° , we have the identity

$$\cos^2P + \cos^2Q + \cos^2R - 2\cos P \cos Q \cos R = 1.$$

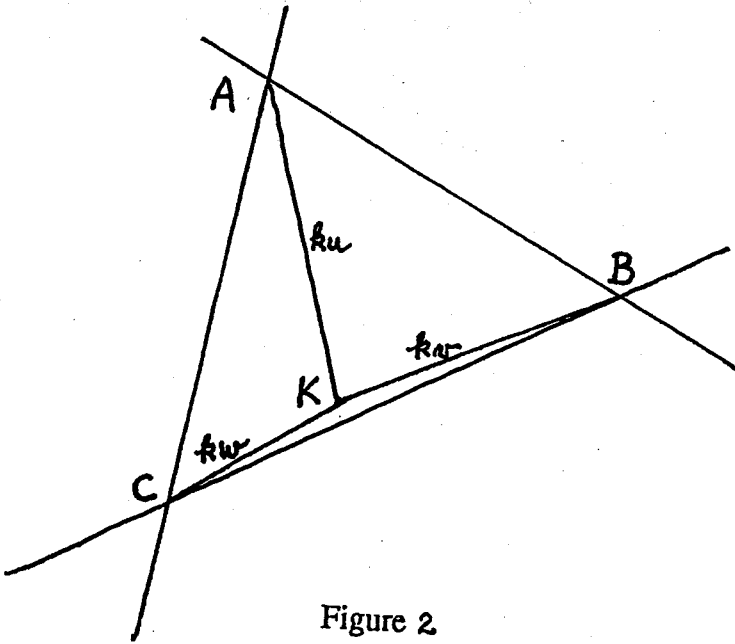


Figure 2

Substitution in this identity yields the following equation for k (all terms in k^6 cancel out):

$$k^4[a^2(u^2-v^2)(u^2-w^2)+ b^2(v^2-u^2)(v^2-w^2)+ c^2(w^2-v^2)(w^2-u^2)] + k^2[(au)^2(a^2-b^2-c^2)+ (bv)^2(b^2-a^2-c^2)+ (cw)^2(c^2-b^2-c^2)] + (abc)^2 = 0.$$

We note that, if $u = v = w = 1$, then k simply becomes R , the circumradius of triangle ABC . The equation then reduces to

$$R^2(-16\Delta^2) + (abc)^2 = 0,$$

and this gives the familiar formula for the circumradius of the triangle, namely,

$$R = abc/4\Delta.$$

We further note that the equation is quadratic in k^2 ; consequently, the condition for k^2 to be real can immediately be written down as

$$[(au)^2(a^2-b^2-c^2)+ (bv)^2(b^2-a^2-c^2)+ (cw)^2(c^2-b^2-c^2)]^2 - 4(abc)^2 [a^2(u^2-v^2)(u^2-w^2)+ b^2(v^2-u^2)(v^2-w^2)+ c^2(w^2-v^2)(w^2-u^2)] \geq 0.$$

This condition simplifies to the requirement that the product of

$$\{2(a^2b^2 + a^2c^2 + b^2c^2) - a^4 - b^4 - c^4\}$$

$$\text{and } \{2[(au)^2(bv)^2 + (au)^2(cw)^2 + (bv)^2(cw)^2] - (au)^4 - (bv)^4 - (cw)^4\}$$

be non-negative; thus

$$(16\Delta^2)(16T^2) \geq 0.$$

Since ABC is a given triangle, we certainly have $\Delta^2 > 0$. However, the auxiliary triangle may not be real. If the auxiliary triangle is real, then $T^2 > 0$, and the two values for k^2 correspond to the two points of intersection of the circles. If $T^2 = 0$, then there is only one point of intersection of the circles (the common chord becomes a common tangent). Finally, if $T^2 < 0$, then the auxiliary triangle ceases to be real, and the circles do not possess real points of intersection.

We thus have

Lemma 4. There are two real points of intersection of the circles if and only if the conditions

$$au + bv > cw, \quad bv + cw > au, \quad cw + au > bv,$$

are satisfied. If one of these equalities becomes an equality, then $T = 0$, and there is a single point of intersection. If any inequality is reversed, then the circles do not have real points of intersection.

We note that, if $u = v = w = 1$, then the conditions just become the ordinary triangle inequalities for the triangle ABC; these conditions are automatically satisfied and so we have the comforting result that the right bisectors of the sides of the triangle do meet in a real point.

It is also possible to deduce the result of Lemma 4 in a somewhat less symmetrical fashion. Let R_1 and R_2 be the radii of the circles with centres G and F, respectively. Then there will be real points of intersection if and only if we have the two conditions

$$(R_1 + R_2)^2 \geq (FG)^2,$$

$$(R_1 - R_2)^2 \leq (FG)^2.$$

It is easy to calculate that

$$R_1 = uvc/(u^2 - v^2), \quad R_2 = uwb/(u^2 - w^2).$$

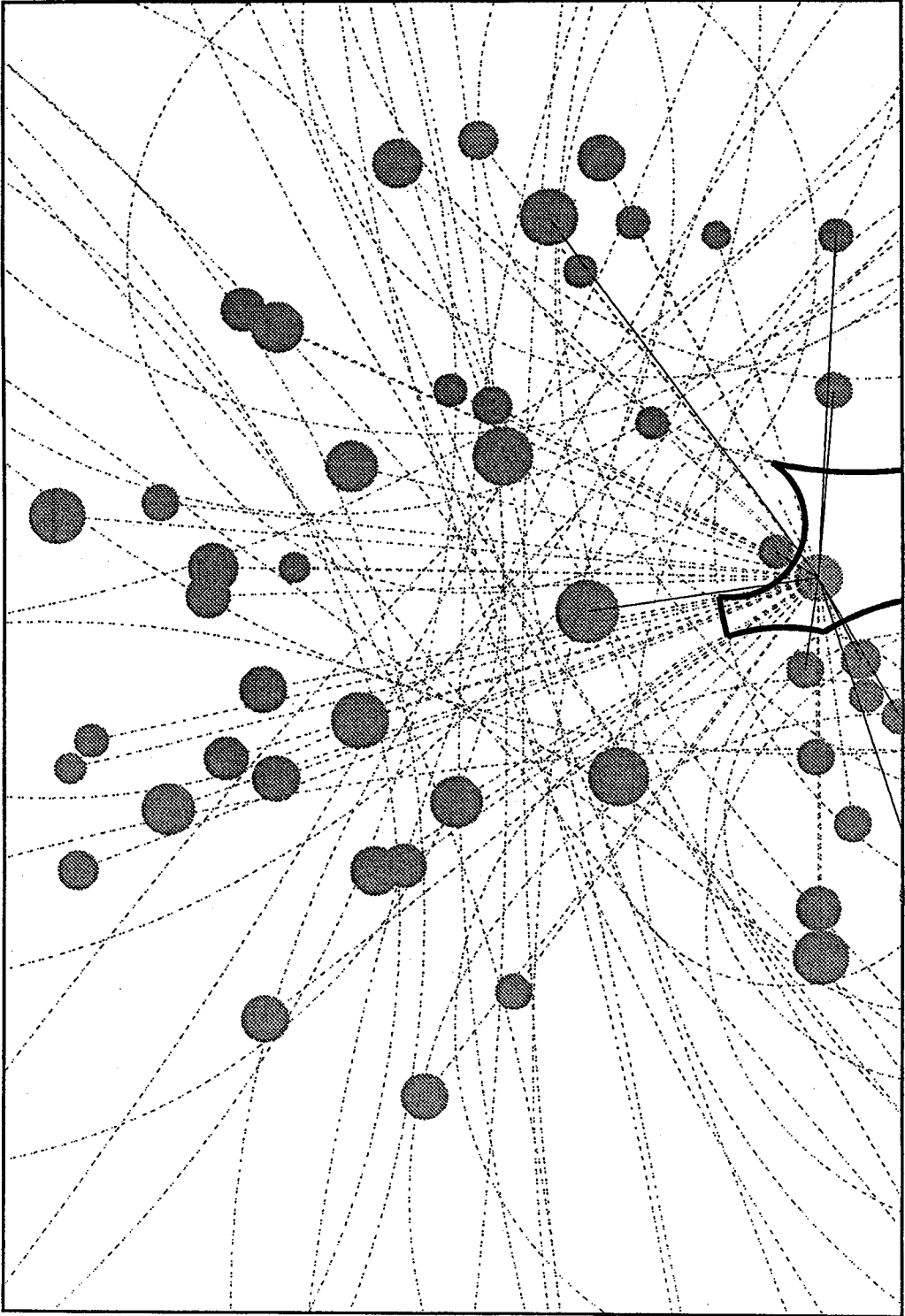
Also, FG can be calculated from the cosine law for triangle AFG. After considerable algebra, these two inequalities reduce to the result already given in Lemma 4.

4. Conclusion. The results described in this paper have been implemented in algorithmic form, and Figures 3 to 10 show some actual results of the algorithm as applied to forest data. The values of u , v , w , are proportional to the attractive powers of the trees (these attractive powers are indicated by the size of the representing points). Figure 3 shows all the construction lines for various combinations of points, and the resulting region dominated by a single point. Figure 4 is identical with Figure 3, except that the construction lines have been eliminated and the regions associated with each point have been drawn in. Figures 5 through 10 indicate some of the different possibilities can arise when there are trees of vastly differing attractive powers (in practice, this involves trees of different sizes or ages).

5. Acknowledgment. We take this opportunity to express our gratitude to Mr Robert Chan who implemented the computer algorithm that produced Figures 3 to 10.

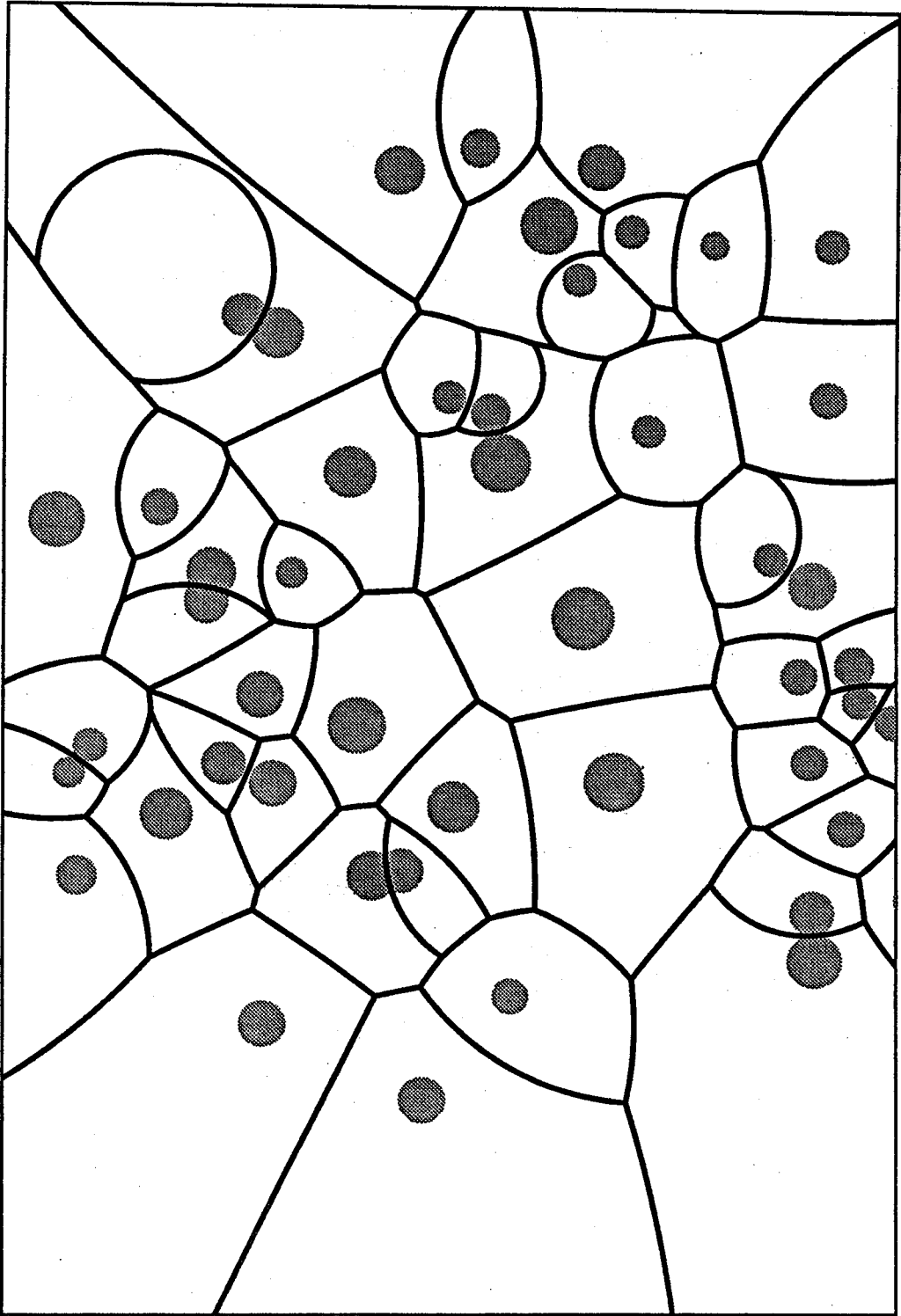
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Internal:RChan:WDT Folder:Vertices Folder:a.a.78.ps.one.tile

Figure 3



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Figure 4

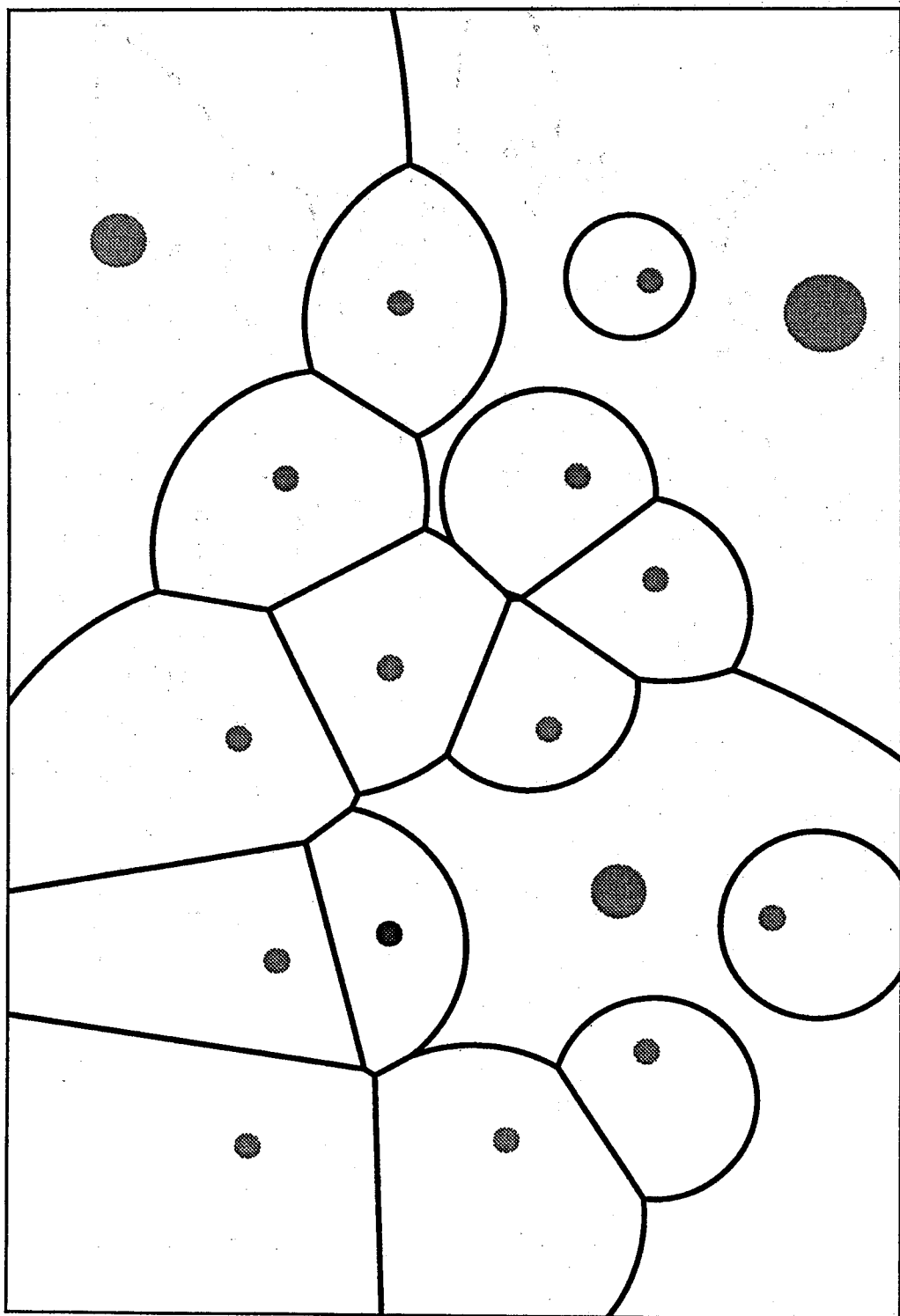
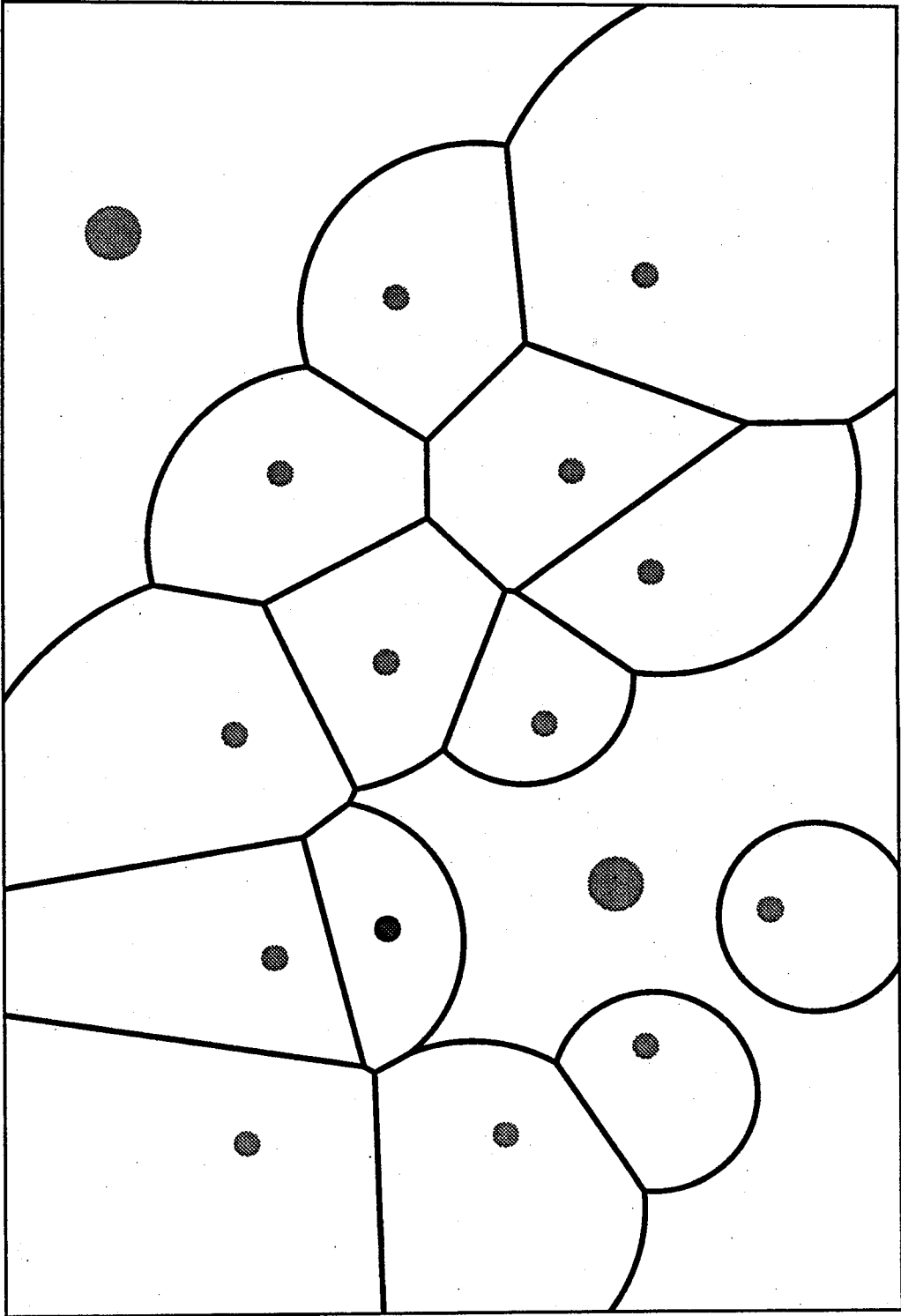


Figure 5

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Figure 6

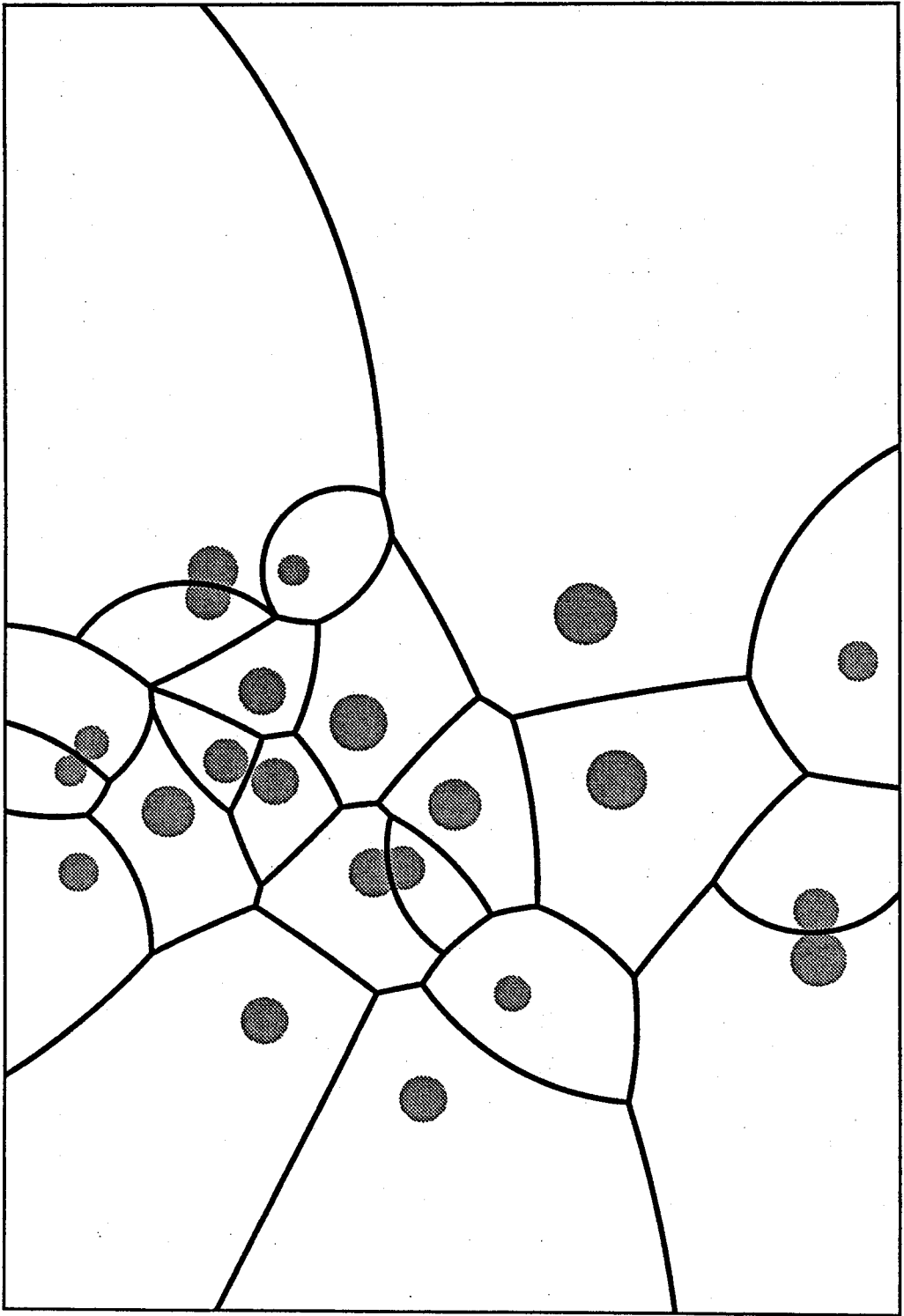
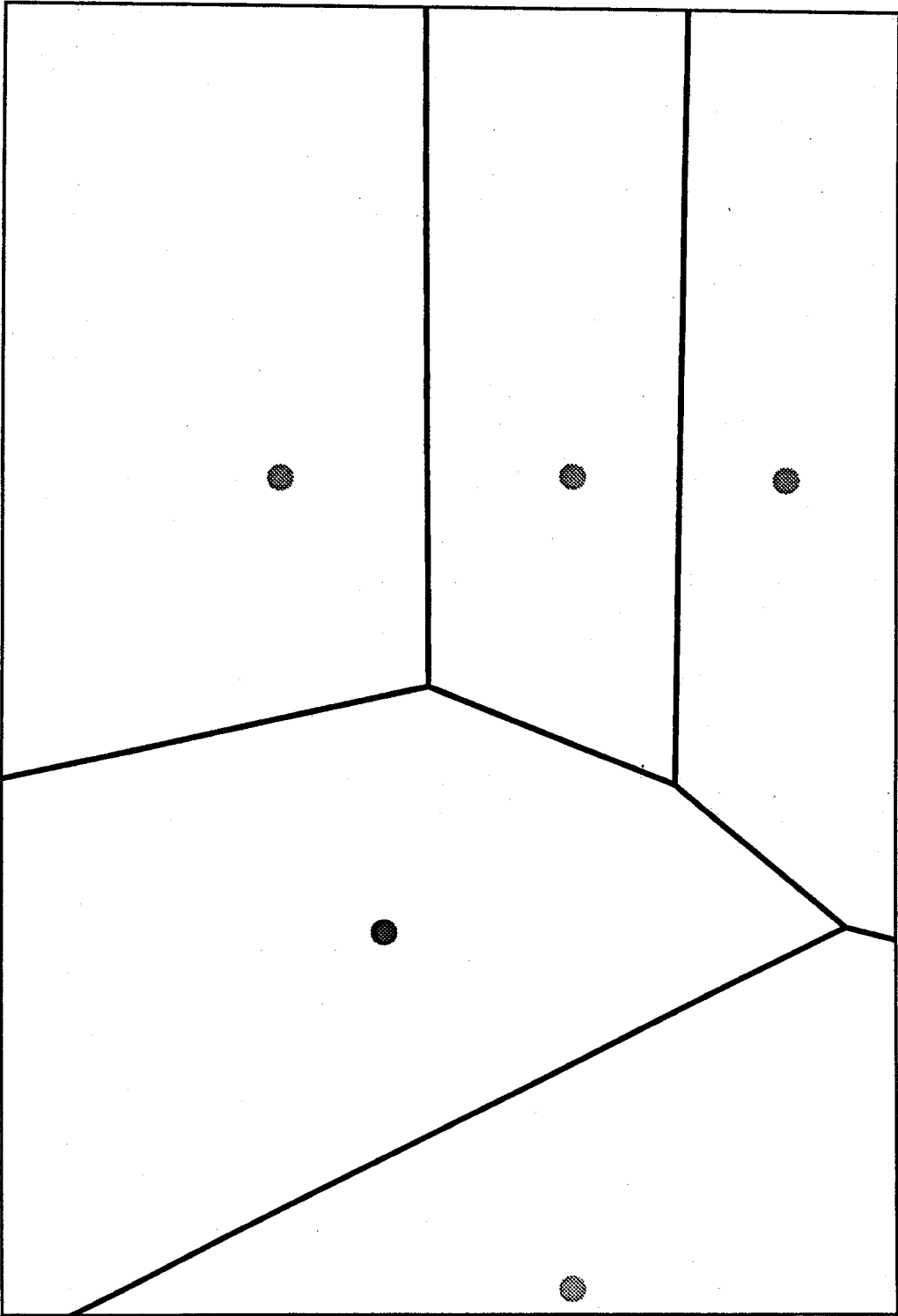


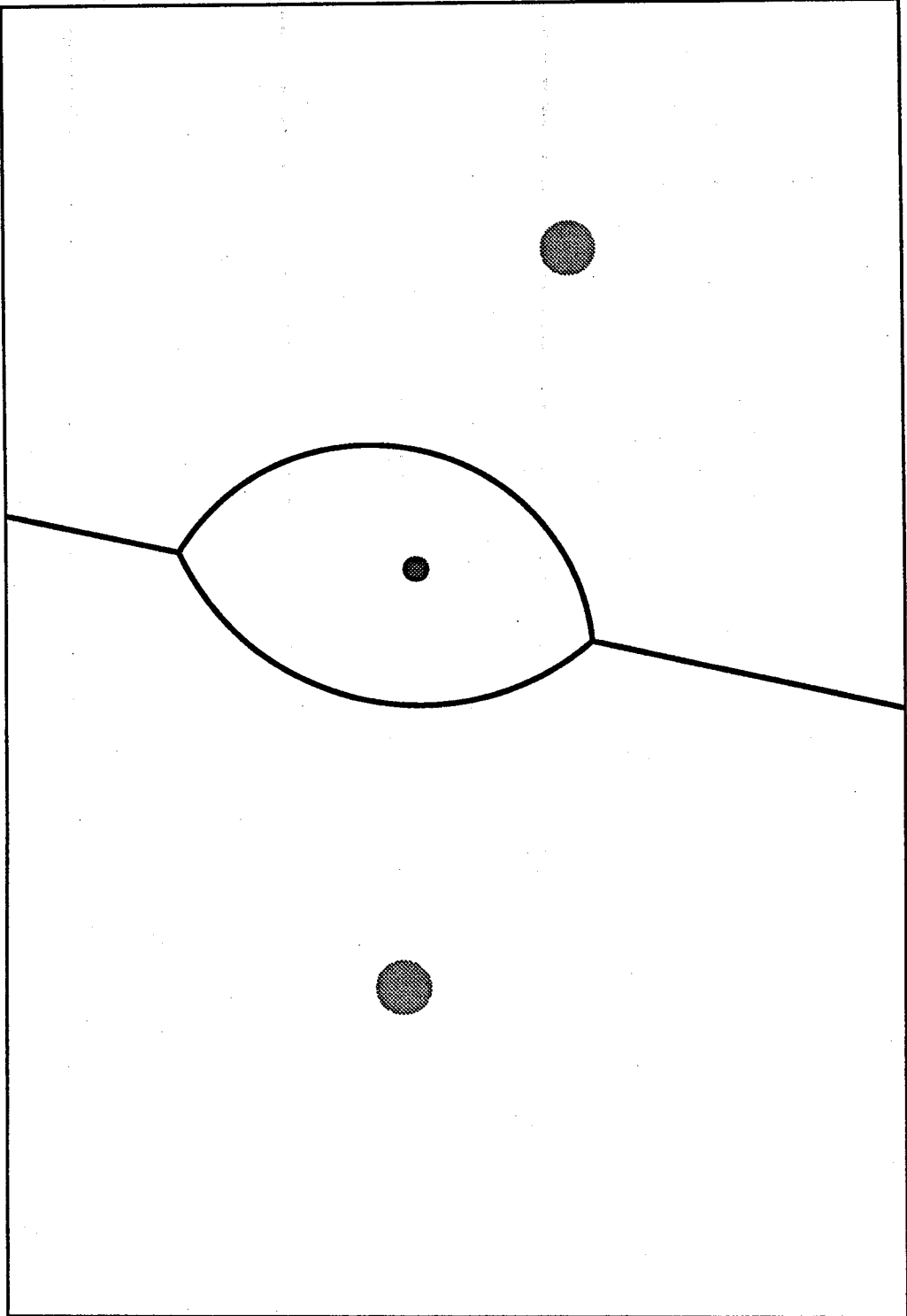
Figure 7

80meg:3



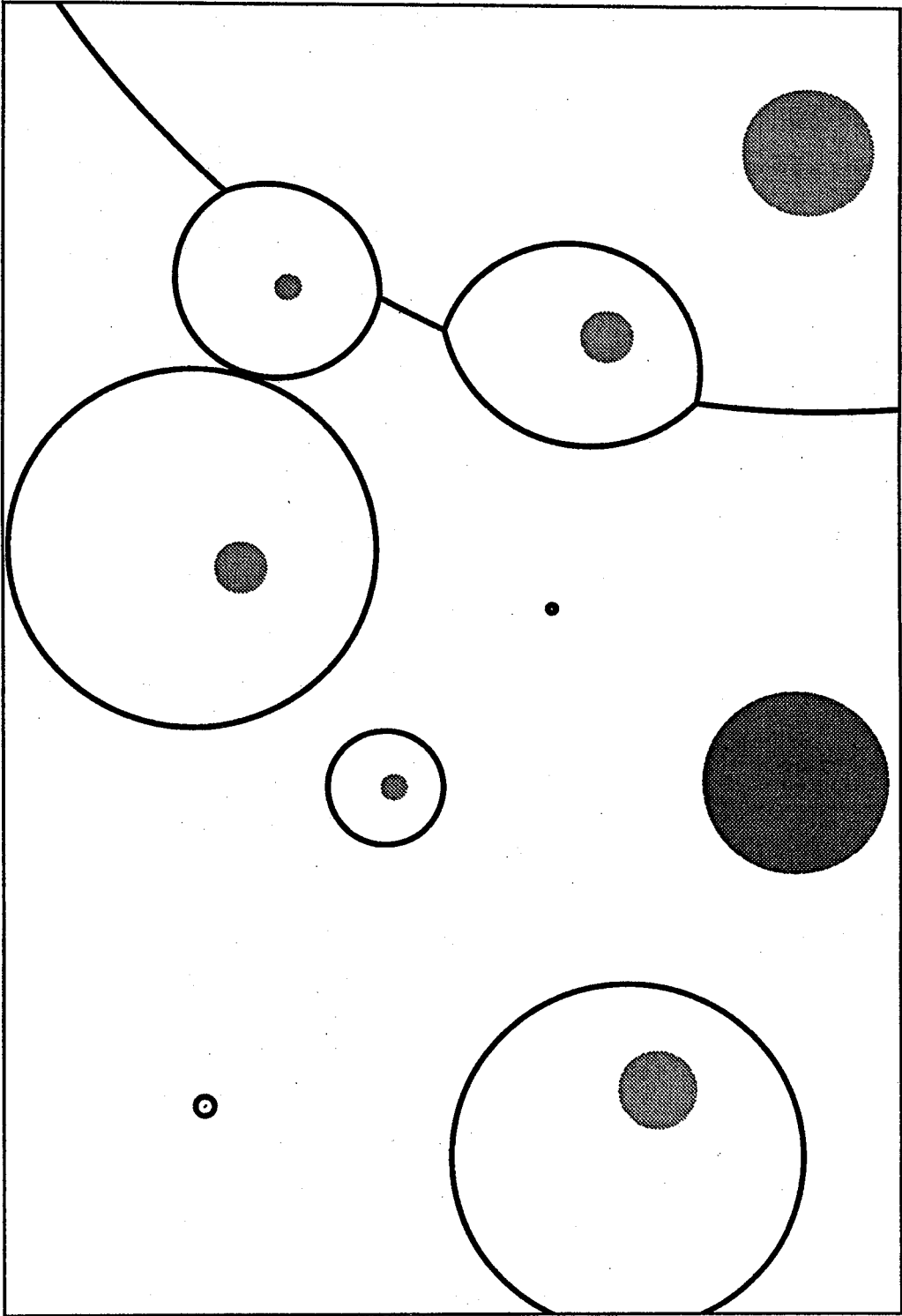
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Figure 8



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Figure 9



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Figure 10

Algebraic-geometric code on a curve

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1. 緒言

V. D. Goppa は代数曲線論と符号理論の間の重要な関係を見い出して代数幾何符号を導入し、それ以前のすべての線形符号は射影直線を基底とする代数幾何符号とみなせることを示した^{2),3)}。M. Tsfasman, S. G. Vladut, Th. Zink は Y. Ihara の結果¹⁾を用いて、従来の線形符号では到達できなかった極めて性能の良い代数幾何符号の無限列を構成した¹⁰⁾。また、J. H. van Lint and T. A. Springer⁶⁾、J. P. Hansen⁵⁾、J. Yustesen et al.¹²⁾等は、具体的な平面代数曲線の上に代数幾何符号を構成し、復号法に関しても考察した。Hansenは $GF(2^3)$ 上定義された Klein の4次曲線

$$X: x^3y + y^3z + z^3x = 0$$

を基底とする符号長21の代数幾何符号を構成し、その性能がBCH符号とほぼ同等であることを調べた。

ここでは、 $l \equiv 3 \pmod{6}$ 、 $s=2^l+1$ とするとき、 $GF(2^{3l})$ 上で定義された平面代数曲線

$$X_s: x^s y + y^s z + z^s x = 0$$

を基底とする代数幾何符号について報告する。この符号の構成に用いる X_s の F 有理点全体に正則に作用する非可換群 G が存在し、構成された符号は群環 $F[G]$ のイデアルとみなせる。R. M. Tannar は同じ非可換群が作用する符号（ただし代数幾何符号ではない）を構成している。このTannarの結果やJ. Yustesen et al.

の結果についても一部紹介する。

2. 代数幾何符号

代数幾何符号について、以下で必要となる事柄を復習しておく。代数幾何符号一般についてはGoppaの文献⁴⁾を見られたい。また水野⁶⁾、山西¹¹⁾にその概要が紹介されている。

$F = GF(q)$ を位数 q の有限体とし、 n 次元ベクトル空間 F^n の部分ベクトル空間 C を符号長 n の線形符号と呼ぶ。

$$x = (x_1, x_2, \dots, x_n), \quad y = (y_1, y_2, \dots, y_n) \in F^n$$

に対し、

$$d(x, y) = \# \{ i \mid x_i \neq y_i \}$$

とにおいて、 F におけるHamming 距離 $d(x, y)$ を定義する。

$$d(C) = \min \{ d(x, y) \mid x, y \in C, x \neq y \}$$

とにおいて $d(C)$ を符号 C の最小距離と呼ぶ。

$$k = \dim_F C$$

を符号 C の次元と呼ぶ。符号は雑音のある通信路における誤りを検出し、訂正することを目標とするので、与えられた符号長および次元に対して大きな最小距離を有する符号を構成することが重要である。

X を有限体 $F = GF(q)$ 上定義された特異点を持たない m 次代数曲線とする。 X は F の代数的閉包 \bar{F} 上で既約、すなわち絶対既約であるとする。 X の種数を $g(X)$ であらわす。このとき

$$g(X) = \frac{(m-1)(m-2)}{2}$$

体 F 上定義される X 上の有理的関数全体を $F(X)$ であらわす。

F 上有理的な因子

$$E = m_1 Q_1 + m_2 Q_2 + \dots + m_j Q_j$$

に対して

$$L(E) = \{ f \in F(X) \mid (f) \geq -E \}$$

とおく。すなわち $L(E)$ は Q_j で高々 m_j 位の極をもち他では正則な X 上の有理的関数全体である。このとき代数曲線論で次の定理が知られている。

定理 (Riemann-Roch) $L(E)$ は F 上有限次元ベクトル空間となり、その次元は

$$\dim_F L(E) = \deg E - g(X) + 1$$

で与えられる。もし $\deg(E) > 2g(X) - 2$ ならば $i(E) = 0$ である。

$$E = \sum_{i=1}^j m_i Q_i, \quad D = \sum_{i=1}^n P_i$$

をともに体 F 上有理的な因子とし、 E と D のどの成分も互いに異なるものとする。また D の各成分 P_i は X 上の F -有理点でさらに

$$\deg E < \deg D = n$$

とする。ベクトル空間 $L(E)$ からベクトル空間 F^n への写像 ψ を

$$\psi(f) = (f(P_1), f(P_2), \dots, f(P_n))$$

で定義する。このとき写像 ψ の像を曲線 X を底曲線とする代数幾何符号と呼び、 $C(D, E; X)$ とあらわす。符号 $C(D, E; X)$ の次元を k 、最小距離を d とすると Riemann-Roch の定理より

(i) $k = \dim_F L(E)$

(ii) $k \geq \deg E - g(X) + 1$

(iii) $d \geq n - \deg E$

が成り立つ。したがって、良いパラメータをもつ符号を構成するためにはその種数に比較して有理点の数の多い曲線を見出す必要がある。

3. X_s 上の F -有理点

l を $l \equiv 3 \pmod{6}$ を満たす整数、 $s = 2^l + 1$ 、 $F = \text{GF}(2^{3l})$ とおく。体 F 上で

$$X_s : x^s y + y^s z + z^s x = 0$$

定義される平面代数曲線

$$g(X_s) = \frac{s(s-1)}{2} = 2^{l-1}(2^l + 1)$$

を考えよう。 X_s は種数

をもち、絶対既約で特異点をもたない曲線である。

$$Q_0 = [1:0:0], \quad Q_1 = [0:1:0], \quad Q_2 = [0:0:1]$$

は任意の s について X_s の F -有理点になるので、これらを自明な有理点と呼ぶ。

X_s の自明でない F -有理点全体を $R(X_s, F)$ であらわす。

2次射影線形群を $\text{PGL}(2; F)$ であらわす。 β を $\beta^3 + \beta + 1 = 0$ を満たす $\text{GF}(2^3)$ の原始元とし

$$t = \frac{2^{3l} - 1}{2^3 - 1}$$

とおく。さらに $\alpha^t = \beta$ を満たす F の原始元 α を1つとり固定する。 $u = 2^{2l}$, $v = 2^l$ とおき、 $\text{PGL}(2; F)$ の2元 a, b を

$$a = \begin{pmatrix} \alpha & 0 & 0 \\ 0 & \alpha^u & 0 \\ 0 & 0 & \alpha^v \end{pmatrix}, \quad b = \begin{pmatrix} 0 & 0 & 1 \\ 1 & 0 & 0 \\ 0 & 1 & 0 \end{pmatrix}$$

と定義する。このとき

$$K = \langle a, b \rangle \subset \text{PGL}(2, F)$$

とおけば K は X_s に作用し、その位数は $3(2^{2l} + 2^l + 1)$ である。

$r = (2^{2l} + 2^l + 1)$ とおく。 K は基本関係式 $a^r = 1, b^3 = 1, ab = b^v a$ をもつ有限非可換群である。 F の 0 以外の元をつくる乗法群の位数 r の部分群を U とする。

このとき、次の定理が成り立つ。すなわち

定理 $x^s + x + 1 = 0$ が U の生成元をその根にもてば、

$$\#R(X_s, F) \geq r(1+2).$$

定理の証明には先に構成した群 K を用いる。その詳細は水野, 安藤⁸⁾ を見られたい。

4. X_s 上の代数幾何符号

前節の定理により得られる X_s の F 有理点を

$$\{P_i \mid i = 0, 1, \dots, 3(1+2)(2^{2l} + 2^l + 1)\}$$

とおく。

m を $\frac{2^{l-1}(2^l - 1)}{3} < m < \frac{n}{3}$ を満たす整数とする。

X_s の正因子として

$$D = \sum_{i=1}^{3(1+2)(2^{2l} + 2^l + 1)} P_i, \quad E = m(Q_0 + Q_1 + Q_2)$$

をとり、 X_s 上の代数幾何符号 $C = C(D, E; X_s)$ を考える。 C の符号長は

$$\deg D = n$$

である。 C の次元を k 、最小距離を d とおくと、

$$(i) \quad n = 3(1+2)(2^{2l} + 2^l + 1),$$

$$(ii) \quad k \geq 3m - 2^{l-1}(2^l + 1) + 1,$$

$$(iii) \quad d \geq n - 3m$$

が成立する。

例 $l = 3$ の場合。

このとき、 $s = 9$ 、 $F = \text{GF}(2^9)$ 、 $g(X_s) = 36$ 、 $r = 73$ 、 $|K| = 219$ となる。 α を $\alpha^9 + \alpha^4 + 1 = 0$ を満たす $\text{GF}(2^9)$ の原始元とする。このとき、 $\xi = \alpha^{119}$ は $\xi^9 + \xi + 1 = 0$ を満たす U の生成元である。定理より X_9 上には 1095 点の F -有理点が存在する。

これらの有理点を用いて $36 \leq m \leq 394$ なる整数 m について、

- (i) $n = 1095$,
- (ii) $k \geq 3m - 35$,
- (iii) $d \geq n - 3m$

なる $\text{GF}(2^9)$ 上の符号が構成される。

5. Yustesen 達の代数幾何符号と Tanner の符号

この節では Yustesen 達の代数幾何符号と Tanner の符号について簡単に紹介する。

まず、Yustesen 達の代数幾何符号の定義を述べよう。

$F = \text{GF}(q)$ を有限体、 X を m 次代数曲線、 $\{P_1, \dots, P_n\}$ を X 上の F -有理点の集合とする。 j を $j < q$ なる整数とすると、 j 次同次多項式全体と $\{0\}$ からなる F 上のベクトル空間を V_j であらわす。曲線 X を基底とする符号長 n の 2 つの線形符号 G および H を

$$G_X(j) = \{f(P_1), f(P_2), \dots, f(P_n) : f \in V_j\},$$

$$H_X(j) = G_X(j)^\perp$$

で定義する。これらが Yustesen 達の代数幾何符号である。

$G_X(j)$ の次元を $k(G)$ 、最小距離を $d(G)$ とおくと、代数曲線論の Bezout の定理より、 $n > mj$ なる m について

$$k(G) = \begin{cases} \binom{j+2}{2} & j < m \\ \binom{j+2}{2} - \binom{j-m+2}{2} & j \geq m \end{cases}$$

$$d(G) \geq n - mj,$$

が成り立つ。

その定義より $H_X(j)$ の次元 $k(H)$ は $k(H) = n - k(G)$ 。よって X が非特異でかつ $j \geq m$ ならば

$$k(H) = mj - g(X) + 1$$

$H_X(j)$ の最小距離 $d(H)$ については、 X が非特異、 $n > mj$ 、 $j \geq m - 2$ の条件もとに Riemann-Roch の定理から

$$d(H) \geq mj - 2g(X) + 1$$

が成り立つ。

また彼等は 1 変数多項式から出発する方法で素数 p と正整数 r にたいして、 $(p^{3r} - 1)(p^{2r} - 1)$ 点の $\text{GF}(p^{6r})$ -有理点をもつ曲線の族

$$x^{p^r - 1} z^{p^{2r} - 1} - z^{p^r - 1} y^{p^{2r} - 1} + x^{p^{2r} - 1} y^{p^r - 1} = 0$$

を得ることに成功した。

一方 Tanner は群の作用する線形符号について研究し、新しい線形符号を構成した(代数曲線上の符号ではない)。 n を正整数、 j, q を n と互いに素な数とする。 Z_n の置換 $i \rightarrow i + j, i \rightarrow qi$ をそれぞれ A, M であらわし、 A, M で生成される群

$$\langle A, M \rangle$$

を考える。 Tanner はこの群の作用する符号長 n の $\text{GF}(q)$ 上の符号を調べ、多くの線形符号を構成した。彼の符号の中には、例えば、符号長 315、次元 271、最小距離 11 なるパラメータを有するものもある。

Yustesen 達や Tanner はともに復号に関しても考察している。詳細については文献 9), 12) を参照されたい。

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Construction of Even Unimodular Lattices from Self-Dual Codes over Finite Fields

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Dec.1,1989

1 Self-dual Code over a Finite Field

Let p be a prime and $\mathbb{F}_p = GF(p)$ the field of p elements. Let $V = \mathbb{F}_p^n$ be the vector space of dimension n over \mathbb{F}_p . A linear $[n, k]$ code C is a vector subspace V of dimension k . In V , the inner product, which is denoted by (x, y) for x, y in V , is defined as usual. The dual code C^\perp of C is defined by

$$C^\perp = \{u \in V \mid (u, v) = 0 \forall v \in C\}.$$

The code C is called self-orthogonal $\iff C \subseteq C^\perp$.

The code C is called self-dual $\iff C = C^\perp$.

2 Complete Weight Enumerator

Let X_0, X_1, \dots, X_{p-1} be independent variables, then the complete weight enumerator $W_C(X_0, X_1, \dots, X_{p-1})$ of C is defined by

$$\begin{aligned} W_C(X_0, X_1, \dots, X_{p-1}) \\ = \sum_{v \in C} X_0^{n_0(v)} X_1^{n_1(v)} \dots X_{p-1}^{n_{p-1}(v)}, \end{aligned}$$

where $n_i(v)$ is the number of the coordinates v_j of the codeword $v = (v_1, v_2, \dots, v_n)$ such that $v_j \equiv i \pmod{p}$. Obviously we have

$$n_0(v) + n_1(v) + \dots + n_{p-1}(v) = n$$

and

$$n_0(\mathbf{v}) + 2^2 n_1(\mathbf{v}) + \cdots + (p-1)^2 n_{p-1}(\mathbf{v}) \equiv 0 \pmod{p}.$$

The last congruence comes from the fact $(\mathbf{v}, \mathbf{v}) = 0$ for $\mathbf{v} \in \mathbf{C} = \mathbf{C}^\perp$. The polynomial $W_{\mathbf{C}}(X_0, X_1, \dots, X_{p-1})$ can be rewritten as

$$W_{\mathbf{C}}(X_0, X_1, \dots, X_{p-1}) = \sum_{n_0+n_1+\dots+n_{p-1}=n} A(n_0, n_1, \dots, n_{p-1}) X_0^{n_0} X_1^{n_1} \cdots X_{p-1}^{n_{p-1}},$$

where $A(n_0, n_1, \dots, n_{p-1})$ is the number of codewords \mathbf{v} in \mathbf{C} such that

$$n_0(\mathbf{v}) = n_0, n_1(\mathbf{v}) = n_1, \dots, n_{p-1}(\mathbf{v}) = n_{p-1}.$$

The cases where we have clear description of $W_{\mathbf{C}}(X_0, X_1, \dots, X_{p-1})$, are $p=2$ and $p=3$ (Conf.[4],[5], and [6]). A rather unsatisfactory description of the complete weight enumerator for $p=5$ is given in [7].

3 Construction of Even Unimodular Lattices from Codes over $GF(p)$

Let \mathbf{C} be a self-dual $[2m, m]$ code over \mathbf{F}_p . We take vectors $\omega_1, \omega_2, \dots, \omega_{2m} \in \mathbf{R}^{2m}$ so that

$$(\omega_i, \omega_j) = p\delta_{ij}.$$

holds.

We form a lattice M by

$$M = [\pm\omega_i \pm \omega_j]_{\mathbf{Z}}$$

Note that M is even integral and $d(M) = 2^2 p^{2m}$. Put $M^\dagger = \{\mathbf{z} \in \mathbf{R}^{2m} \mid (\mathbf{x}, \mathbf{z}) \in \mathbf{Z} \ \forall \mathbf{x} \in M\}$. M^\dagger is the dual lattice of M and it satisfies

$$\text{index}[M^\dagger : M] = 2^2 p^{2m}.$$

Let $\mathbf{v}_i = (v_{i1}, v_{i2}, \dots, v_{i2m})$ ($1 \leq i \leq m$) be a basis of \mathbf{C} . From \mathbf{v}_i we define vectors \mathbf{y}_i in a following manner :

$$\mathbf{y}_i = \frac{1}{p} \sum_{j=1}^{2m} b_{ij} \omega_j \quad (1 \leq i \leq m),$$

where b_{ij} are integers satisfying

$$b_{ij} \bmod p = v_{ij}.$$

We put

$$J = M + \mathbb{Z}y_1 + \mathbb{Z}y_2 + \cdots + \mathbb{Z}y_m.$$

It can be verified that J is an integral lattice. By taking b_{ij} 's suitably, we can make J to be even (by adjusting $(y_i, y_i) \equiv 0 \pmod{2}$). A general element y of J is expressed as

$$y = u + \sum_{i=1}^m a_i y_i, \quad a_i \in \mathbb{Z}, u \in M.$$

$\text{supp}(y)$ is a vector in \mathbb{F}_p^{2m} defined by

$$\text{supp}(y) = \sum_{i=1}^m \bar{a}_i v_i,$$

where

$$\bar{a}_i = a_i \bmod p.$$

We can prove that the mapping

$$\varphi : y \mapsto \text{supp}(y)$$

defines a linear mapping from J to \mathbb{C} and the kernel of φ is M . Therefore

$$J/M \simeq \mathbb{C}.$$

In particular, we have

$$[J : M] = |\mathbb{C}| = p^m.$$

At this stage we know that

$$M^{\perp} \supset J^{\perp} \supset J \supset M$$

and

$$[M^{\perp} : M] = [M^{\perp} : J^{\perp}][J^{\perp} : J][J : M] = 2^2 p^{2m}.$$

Since $[M^{\perp} : J^{\perp}] = [J : M] = p^m$, we have

$$[J^{\perp} : J] = 2^2.$$

If we choose $y_0 \in J^{\perp} - J$ with $(y_0, y_0) \equiv 0 \pmod{2}$ and form a lattice L :

$$L = J + \mathbb{Z}y_0,$$

then L is an even unimodular lattice.

4 Special Topics

The Gosset lattice E_8 has many representations. It is a well-known fact that E_8 is constructed from the Hamming binary $[8, 4, 4]$ code. Here we give several constructions of E_8 from other codes over $GF(p)$ ($p \geq 3$).

*(T) Ternary code construction of E_8 .

$$C = \begin{pmatrix} 1 & 1 & 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 2 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 1 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 & 2 & 1 \end{pmatrix}.$$

C is a generator matrix of a self-dual $[8, 4]$ code over $GF(3)$. Let ω_i ($1 \leq i \leq 8$) be vectors in \mathbb{R}^8 s.t. $(\omega_i, \omega_j) = 3\delta_{ij}$ and $M = [\pm\omega_i \pm \omega_j]_{\mathbb{Z}}$. Representatives y_1, y_2, y_3 and y_4 of J/M are

$$\begin{aligned} y_1 &= \frac{1}{3}(\omega_1 + \omega_2 - 2\omega_3), y_2 = \frac{1}{3}(\omega_2 - \omega_3 - 2\omega_4) \\ y_3 &= \frac{1}{3}(\omega_5 + \omega_6 - 2\omega_7), y_4 = \frac{1}{3}(\omega_6 - \omega_7 - 2\omega_8) \\ y_0 &= \frac{1}{6}(\omega_1 + \omega_2 + \omega_3 - 3\omega_4 + \omega_5 + \omega_6 + \omega_7 - 3\omega_8) \end{aligned}$$

$$E_8 = J + \mathbb{Z}y_0, 2y_0 \in J, y_0 \notin J.$$

*(Q) Quinary code construction of E_8 .

$$C = \begin{pmatrix} 1 & 0 & 0 & 0 & 1 & 1 & 1 & 1 \\ 0 & 1 & 0 & 0 & 1 & 1 & 4 & 4 \\ 0 & 0 & 1 & 0 & 1 & 4 & 1 & 4 \\ 0 & 0 & 0 & 1 & 1 & 4 & 4 & 1 \end{pmatrix}.$$

C is a generator matrix of a self-dual $[8, 4]$ code over $GF(5)$. Let ω_i ($1 \leq i \leq 8$) be vectors in \mathbb{R}^8 s.t. $(\omega_i, \omega_j) = 5\delta_{ij}$ and $M = [\pm\omega_i \pm \omega_j]_{\mathbb{Z}}$. Representatives y_1, y_2, y_3 and y_4 of J/M are

$$\begin{aligned} y_1 &= \frac{1}{5}(\omega_1 + \omega_5 + \omega_6 + \omega_7 - 4\omega_8), y_2 = \frac{1}{5}(\omega_2 + \omega_5 + \omega_6 + 4\omega_7 - \omega_8) \\ y_3 &= \frac{1}{5}(\omega_3 + \omega_5 + 4\omega_6 + \omega_7 - \omega_8), y_4 = \frac{1}{5}(\omega_4 - 4\omega_5 - \omega_6 - \omega_7 + \omega_8) \end{aligned}$$

$$y_0 = \frac{1}{10}(\omega_1 + \omega_2 - \omega_3 + \omega_4 - 3\omega_5 - 3\omega_6 + 3\omega_7 - 3\omega_8)$$

$$E_8 = J + \mathbf{Z}y_0, 2y_0 \in J, y_0 \notin J$$

*(S) Septenary code construction of E_8 .

$$\mathbf{C} = \begin{pmatrix} 1 & 0 & 0 & 0 & 2 & 3 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 & 0 & 2 & 3 \\ 1 & 0 & 1 & 1 & 0 & 2 & 0 & 0 \\ 1 & 0 & 1 & 6 & 3 & 0 & 1 & 6 \end{pmatrix}.$$

\mathbf{C} is a generator matrix of a self-dual $[8, 4]$ code over $GF(7)$. Let ω_i ($1 \leq i \leq 8$) be vectors in \mathbf{R}^8 s.t. $(\omega_i, \omega_j) = 7\delta_{ij}$ and $M = [\pm\omega_i \pm \omega_j]_{\mathbf{Z}}$. Representatives y_1, y_2, y_3 and y_4 of J/M are

$$y_1 = \frac{1}{7}(\omega_1 + 2\omega_5 + 3\omega_6), y_2 = \frac{1}{7}(\omega_2 + 2\omega_7 + 3\omega_8)$$

$$y_3 = \frac{1}{7}(\omega_1 + \omega_3 + \omega_4 - 5\omega_6), y_4 = \frac{1}{7}(\omega_1 + \omega_2 - \omega_4 + 3\omega_5 + \omega_7 - \omega_8)$$

$$y_0 = \frac{1}{14}(-2\omega_1 + 2\omega_2 - 4\omega_3 + 2\omega_4 - 2\omega_5 + 2\omega_6 - 4\omega_7 + 2\omega_8)$$

$$E_8 = J + \mathbf{Z}y_0, 2y_0 \in J, y_0 \notin J$$

*(11) $GF(11)$ code construction of E_8 .

$$\mathbf{C} = \begin{pmatrix} 1 & 1 & 3 & 0 & 0 & 0 & 0 & 0 \\ 3 & 8 & 0 & 2 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 1 & 3 & 0 \\ 0 & 0 & 0 & 0 & 3 & 8 & 0 & 2 \end{pmatrix}.$$

\mathbf{C} is a generator matrix of a self-dual $[8, 4]$ code over $GF(11)$. Let ω_i ($1 \leq i \leq 8$) be vectors in \mathbf{R}^8 s.t. $(\omega_i, \omega_j) = 7\delta_{ij}$ and $M = [\pm\omega_i \pm \omega_j]_{\mathbf{Z}}$. Representatives y_1, y_2, y_3 and y_4 of J/M are

$$y_1 = \frac{1}{11}(\omega_1 + \omega_2 - 8\omega_3), y_2 = \frac{1}{11}(3\omega_1 - 3\omega_2 + 2\omega_4)$$

$$y_3 = \frac{1}{11}(\omega_5 + \omega_3 + \omega_6 - 8\omega_7), y_4 = \frac{1}{11}(3\omega_5 - 3\omega_6 + 2\omega_8)$$

$$y_0 = \frac{1}{11}(2\omega_1 - \omega_2 - 4\omega_3 + \omega_4 + 2\omega_5 - \omega_6 - 4\omega_7 + \omega_8)$$

$$E_8 = J + Zy_0, 2y_0 \in J, y_0 \notin J$$

It may be possible that E_8 can be constructed from a self-dual [8, 4]

code over arbitrary finite field.

*(T) Ternary code construction of the Leech lattice

Let $C = [I_{12}H_{12}]$ be a generator matrix of the Pless code P_{24} , where $H_{12} = (h_{ij})$ ($1 \leq i, j \leq 12$) is a Paley type matrix of order 12 viewed as the matrix over $GH(3)$. Let ω_i ($1 \leq i \leq 24$) be vectors in \mathbb{R}^{24} s.t. $(\omega_i, \omega_j) = 3\delta_{ij}$ and $M = [\pm\omega_i \pm \omega_j]_{\mathbb{Z}}$. Representatives x_i ($1 \leq i \leq 24$) of J/M are given by

$$x_i = \frac{1}{3}(\omega_i + \sum_{j=1}^{12} a_{ij}\omega_{j+12}),$$

where a_{ij} are 0,1,-1 according as h_{ij} are 0,1,2 respectively. Put

$$x_0 = \frac{1}{6}(\sum_{j=1}^{23} \omega_j - 5\omega_{24}),$$

$$Leech = J + \mathbb{Z}x_0, 2x_0 \in J, x_0 \notin J$$

There also is a quinary code construction of the Leech lattice (c.f.[10]).

The code construction on the general finite fields in our sense may play an important role in the theory of quadratic forms. To glimpse this we show some illustrating examples.

(I) An even unimodular lattice of rank n is called extremal if the minimal vectors of L is a $2k$ -vector (i.e. a vector x satisfying $(x, x) = 2k$) with $2k = 2[\frac{n}{24}] + 2$. Here we give a small table of results.

| rank n extremal lattice | 8 | 16 | 24 | 32 | 40 | 48 | 56 | 64 |
|----------------------------|---|----|----|----|----|----|----|----|
| | B | B | B | B | B | | | |
| | T | T | T* | T† | T† | T* | T† | T† |
| | Q | Q | Q | ? | ? | ? | ? | |
| | S | S | ? | ? | ? | ? | ? | |

B: it admits binary code construction
 T: it admits ternary code construction
 Q: it admits quinary code construction
 S: it admits septenary code construction
 * : it is found by Leech-Sloane ([2]).
 † : it is found by Ozeki ([9]).

Remark : The places which are marked "??" are not examined.

(II) 24-dimensional Niemeier lattices

| | | | | | | | |
|-------------|-----|----------------|----|------------|-----|-------------|-----|
| D_{24} | B T | $A_{15}D_9$ | T | E_8^3 | B T | $A_9^2D_6$ | T |
| E_8D_{16} | B T | $A_7^2D_5^2$ | T | D_{12}^2 | B T | $A_5^4D_4$ | T |
| D_8^3 | B T | A_4^6 | T? | A_8^3 | T? | A_1^{24} | B T |
| D_4^6 | B T | A_6^4 | T? | A_2^{12} | T | A_3^8 | T |
| D_6^4 | B T | $A_{11}D_7E_6$ | T | E_6^4 | T? | $A_{17}E_7$ | T |
| A_{12}^2 | T? | $D_{10}E_7^2$ | T | A_{24} | T? | Leech | B T |

B: it admits binary code construction
 T: it admits ternary code construction

Remark : The places marked by "T?" are not yet determined whether it admits a ternary code construction. The reason for this is that the classification of self-dual ternary codes of length 24 is not complete (c.f. [3]).

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